

HIGHLY CONNECTED SUBGRAPHS WITH LARGE CHROMATIC NUMBER

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ABSTRACT. For integers $k \geq 1$ and $m \geq 2$, let $g(k, m)$ be the least integer $n \geq 1$ such that every graph with chromatic number at least n contains a $(k+1)$ -connected subgraph with chromatic number at least m . Refining the recent result Girão and Narayanan that $g(k-1, k) \leq 7k+1$ for all $k \geq 2$, we prove that $g(k, m) \leq \max(m+2k-2, \lceil(3+\frac{1}{16})k\rceil)$ for all $k \geq 1$ and $m \geq 2$. This sharpens earlier results of Alon, Kleitman, Saks, Seymour, and Thomassen, of Chudnovsky, Penev, Scott, and Trotignon, and of Penev, Thomassé, and Trotignon.

Our result implies that $g(k, k+1) \leq \lceil(3+\frac{1}{16})k\rceil$ for all $k \geq 1$, making a step closer towards a conjecture of Thomassen from 1983 that $g(k, k+1) \leq 3k+1$, which was originally a result with a false proof and was the starting point of this research area.

1. INTRODUCTION

All graphs in this paper are finite and with no loops or parallel edges. Given a graph G with vertex set $V(G)$, a *cutset* in G is a (possibly empty) set $X \subseteq V(G)$ whose removal from G results in a disconnected graph. For an integer $k \geq 1$, G is said to be *k -connected* if it has more than k vertices and has no cutset of cardinality less than k . A *stable set* of G is a vertex set with pairwise nonadjacent vertices in G . The *chromatic number* of G , denoted by $\chi(G)$, is the least integer $m \geq 0$ such that $V(G)$ can be partitioned into m stable sets.

The starting point of this paper is the following conjecture of Thomassen [13, Theorem 11] from 1983 which was originally a result with a false proof.

Conjecture 1.1. *For every integer $k \geq 1$, every graph with chromatic number more than $3k$ contains a $(k+1)$ -connected graph with chromatic number more than k .*

For integers $k \geq 1$ and $m \geq 2$, let $g(k, m)$ be the smallest integer $n \geq 1$ such that every graph with chromatic number at least n contains a $(k+1)$ -connected subgraph with chromatic number at least m . Thus $g(k, 2)$ is simply the least integer $n \geq 1$ such that every graph with chromatic number at least n contains a $(k+1)$ -connected subgraph, $g(1, m) = \max(m, 3)$ for all $m \geq 2$, and Conjecture 1.1 says $g(k, k+1) \leq 3k+1$ for all $k \geq 1$. Here are the known estimates on $g(k, m)$ for $k, m \geq 2$.

- Alon, Kleitman, Saks, Seymour, and Thomassen [1], seeking for a remedy for the incorrect proof of Conjecture 1.1 in [13], initiated the study of $g(k, m)$ and proved that¹

$$\max(m+k-1, 2k+1) \leq g(k, m) \leq \max(m+10k^2, 100k^3+1).$$

- Partly motivated by the study of χ -boundedness (see [12] for a survey on this topic), Chudnovsky, Penev, Scott, and Trotignon [5] proved (among other things) that

$$g(k, m) \leq \max(m+2k^2, 2k^2+k+1)$$

and Penev, Thomassé, and Trotignon [11] proved that

$$g(k, m) \leq \max(m+2k-2, 2k^2+1).$$

- Motivated by recent progress on Hadwiger's conjecture [6], Girão and Narayanan [8] proved that $g(k-1, k) \leq 7k+1$; in fact their argument can be slightly modified to get $g(k-1, k) \leq 4k$.

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¹The authors of [1] only stated that $g(k, m) \geq m+k-1$, but their lower bound construction in fact shows that $g(k, m) \geq \max(m+k-1, 2k+1)$.

A classical result of Mader [9] (see also [7, Theorem 1.4.3]) says that for every $k \geq 1$, every graph with average degree at least $4k$ contains a $(k+1)$ -connected subgraph with more than $2k$ vertices. This leads to two natural questions:

- What is the smallest constant $C > 0$ such that every graph with average degree at least Ck contains a $(k+1)$ -connected subgraph with more than $2k$ vertices? Carmesin [4] recently showed that $C = 3 + \frac{1}{3}$ is the correct answer.
- What if we just ask for the $(k+1)$ -connectivity without any demands on the number of vertices of the subgraph? Mader [10] conjectured that every graph with average degree at least $3k-1$ contains a $(k+1)$ -connected subgraph, and the current record on this problem is held by Bernshteyn and Kostochka [2] who proved that $(3 + \frac{1}{6})k$ suffices as long as the host graph has at least $\frac{5}{2}k$ vertices.

Our main result, which can be considered as an analogue of Carmesin's theorem for the chromatic number in some sense, improves all of the aforementioned upper bounds on $g(k, m)$. It shows that

$$g(k, m) \leq \max \left(m + 2k - 2, \left\lceil \left(3 + \frac{1}{16} \right) k \right\rceil \right)$$

for all $k \geq 1$ and $m \geq 2$, approaching the lower bound $g(k, m) \geq \max(m + k - 1, 2k + 1)$; and consequently, Conjecture 1.1 is true with 3 replaced by $3 + \frac{1}{16}$.

Theorem 1.2. *For every integer $k \geq 1$, every graph G with $\chi(G) \geq (3 + \frac{1}{16})k$ contains a $(k+1)$ -connected subgraph with more than $\chi(G) - k$ vertices and chromatic number at least $\chi(G) - 2k + 2$.*

We suspect that the constant factor $3 + \frac{1}{16}$ in Theorem 1.2 can be replaced by 3, which would verify Conjecture 1.1.

2. TEMPLATES AND INEXTENSIBILITY

The proof of Theorem 1.2 employs the “template-inextensibility” method first appeared in [5] (under the name “coloring constraints”) and developed further in [8], with a number of modifications. Given sets $A \subseteq B$ and a map f with domain B , let $f|_A$ be the restriction of f to A , and let $f(A) := \{f(a) : a \in A\}$. For a graph G , let $|G|$ denote the number of vertices of G . For $S \subseteq V(G)$, let $G[S]$ denote the subgraph of G with vertex set S and edges whose endpoints are in S ; a graph H is an *induced subgraph* of G if there is $S \subseteq V(G)$ with $H = G[S]$, and H is a *proper induced subgraph* of G if $|H| < |G|$.

For a finite set \mathcal{C} of colors, a *proper \mathcal{C} -coloring* of G is a function $f : V(G) \rightarrow \mathcal{C}$ satisfying $f(u) \neq f(v)$ for all u, v adjacent vertices of G ; thus $\chi(G)$ is the least integer $m \geq 0$ such that there is a proper \mathcal{C} -coloring of G with $|\mathcal{C}| = m$. A *\mathcal{C} -template on G* is a triple $T = (S, c, F)$ where

- S is a subset of $V(G)$;
- $c : S \rightarrow \mathcal{C}$ is a proper \mathcal{C} -coloring of $G[S]$; and
- F is a map from $V(G) \setminus S$ to the family of all subsets of \mathcal{C} .

It might be helpful to think of the vertices of S as the *precolored* vertices, and each $v \in V(G) \setminus S$ as an *uncolored* vertex with a list of *forbidden colors* specified by $F(v)$.

For an integer $k \geq 1$, given a \mathcal{C} -template $T = (S, c, F)$ on G , define the *k -cost* of T by

$$\text{cost}_k(T) := k|S| + \sum_{v \in V(G) \setminus S} |F(v)|.$$

Every $A \subseteq V(G)$ naturally gives rise to the \mathcal{C} -template $T_A := (S \cap A, c|_{S \cap A}, F|_{A \setminus S})$ on $G[A]$. Note that cost_k is additive under disjoint unions, that is, $\text{cost}_k(T_{A \cup B}) = \text{cost}_k(T_A) + \text{cost}_k(T_B)$ for all disjoint $A, B \subseteq V(G)$.

A proper \mathcal{C} -coloring f of G is said to *respect* T if $f|_S = c$ and $f(v) \in \mathcal{C} \setminus F(v)$ for all $v \in V(G) \setminus S$. Say that G is *\mathcal{C} -inextensible* if there exists a \mathcal{C} -template $T = (S, c, F)$ on G such that

- $\text{cost}_k(T) < 2k^2$;
- $|F(v)| \leq k$ for all $v \in V(G) \setminus S$; and
- there is no proper \mathcal{C} -coloring of G respecting T .

In this case, say that T *witnesses* the \mathcal{C} -inextensibility of G ; note that $|S| < 2k$. Observe that if $\chi(G) > 0$, then there exists \mathcal{C} so that G is \mathcal{C} -inextensible: indeed, such a \mathcal{C} can be chosen to be any set of $\chi(G) - 1$ colors, then the \mathcal{C} -inextensibility of G is witnessed by the empty \mathcal{C} -template with no precolored vertices and no forbidden colors at each uncolored vertex. Say that G is \mathcal{C} -*extensible* if it is not \mathcal{C} -inextensible.

In what follows, when there is no danger of ambiguity, we drop the prefix \mathcal{C} from the notions of proper \mathcal{C} -colorings, \mathcal{C} -templates, and \mathcal{C} -inextensibility. We also drop the prefix k from the notion of k -costs, and drop the subscript k from the notation cost_k .

3. CONNECTIVITY

We wish to show that, given a set \mathcal{C} of colors, if $|\mathcal{C}|$ is sufficiently large, then every inextensible graph contains a $(k + 1)$ -connected subgraph. Here is an example showing that $|\mathcal{C}| \geq 3k - 1$ is necessary: if $|\mathcal{C}| = 3k - 2$, let G be a star² with $2k$ vertices, and consider a template on G where the leaves are precolored by different colors and the center has $k - 1$ forbidden colors which are not used for the leaves; then this template witnesses the inextensibility of G while G certainly has no 2-connected subgraphs. It turns out that $|\mathcal{C}| \geq 3k - 1$ is also sufficient. To see this, let us say that a template $T = (S, c, F)$ on an inextensible graph G is *good* if T witnesses the inextensibility of G and $|F(v)| \leq k - 1$ for all $v \in V(G) \setminus S$, that is, each uncolored vertex has fewer than k forbidden colors. The following lemma says that every inextensible graph has a good template as long as $|\mathcal{C}| \geq 3k - 1$.

Lemma 3.1. *If $|\mathcal{C}| \geq 3k - 1$ and G is inextensible, then there is a good template on G .*

Proof. Since G is inextensible, there exists a template $T = (S, c, F)$ witnessing its inextensibility; choose T with $|S|$ maximal. Suppose there is $v \in V(G) \setminus S$ with $|F(v)| = k$; let $S' := S \cup \{v\}$. We have that

$$2k^2 - k|S| > \text{cost}(T) - k|S| \geq |F(v)| = k,$$

and so $|S| < 2k - 1$. Because $|\mathcal{C}| \geq 3k - 1$, it follows that (recall that $c(S) = \{c(u) : u \in S\}$)

$$|\mathcal{C} \setminus (c(S) \cup F(v))| \geq |\mathcal{C}| - |S| - |F(v)| > (3k - 1) - (2k - 1) - k = 0,$$

so $\mathcal{C} \setminus (c(S) \cup F(v)) \neq \emptyset$. Let $T' := (S', c', F|_{V(G) \setminus S'})$ be a template on G with c' satisfying $c'|_S = c$ and $c'(v) \in \mathcal{C} \setminus (c(S) \cup F(v))$. Then

$$\text{cost}(T') = \text{cost}(T) + k - |F(v)| = \text{cost}(T) < 2k^2$$

so T' would witness the inextensibility of G with $|S'| > |S|$, contradicting the maximality of $|S|$. Therefore $|F(v)| \leq k - 1$ for all $v \in V(G) \setminus S$, that is, T is good on G . This proves Lemma 3.1. \blacksquare

A graph is said to be *minimally inextensible* if it is inextensible while its proper induced subgraphs are extensible. It is immediate that every inextensible graph contains a minimally inextensible induced subgraph; and as the following lemma shows, every minimally inextensible graph is $(k + 1)$ -connected as long as $|\mathcal{C}| \geq 3k - 1$. This constitutes the connectivity part of Theorem 1.2.

Lemma 3.2. *If $|\mathcal{C}| \geq 3k - 1$, then every minimally inextensible graph has more than $|\mathcal{C}| - k + 1$ vertices and is $(k + 1)$ -connected.*

Proof. Let G be a minimally inextensible graph. By Lemma 3.2, there is a good template $T = (S, c, F)$ on G .

Claim 3.1. *Every vertex in S has more than k neighbors in $V(G) \setminus S$.*

Proof. Let $u \in S$, and let M be the set of neighbors of u in $V(G) \setminus S$. Let $T' := (S \setminus \{u\}, c|_{S \setminus \{u\}}, F')$ be the template on $G \setminus u$ with F' defined by

- $F'(v) := F(v)$ for all $v \in V(G) \setminus (S \cup M)$; and
- $F'(v) := F(v) \cup \{c(u)\}$ for all $v \in M$.

²A *star* is a complete bipartite graph with one part having only one vertex (called the *center*); and the vertices in the other part are called the *leaves*. If both parts have size one each then the center can be either one of the two vertices.

As $|F(v)| \leq k - 1$ for all $v \in V(G) \setminus S$, we have that $|F'(v)| \leq k$ for all $v \in V(G) \setminus S$. Observe that

$$\begin{aligned} \text{cost}(T') &\leq k|S \setminus \{u\}| + \sum_{v \in V(G) \setminus (S \cup M)} |F(v)| + \sum_{v \in M} (|F(v)| + 1) \\ &= k|S| - k + \sum_{v \in V(G) \setminus S} |F(v)| + |M| = \text{cost}(T) + |M| - k. \end{aligned}$$

If $|M| \leq k$, then $\text{cost}(T') \leq \text{cost}(T) < 2k^2$; so the extensibility of $G \setminus u$ would give a proper coloring of $G \setminus u$ respecting T' and so a proper coloring of G respecting T , a contradiction. This shows that $|M| > k$, as required. This proves Claim 3.1. \square

Claim 3.2. *Every vertex in $V(G) \setminus S$ has more than $|\mathcal{C}| - k$ neighbors in G .*

Proof. Let $v \in V(G) \setminus S$, and let N be the set of neighbors of v in G . The extensibility of $G \setminus v$ yields a proper coloring c' of $G \setminus v$ respecting $T_{V(G) \setminus \{v\}}$. If $|N| \leq |\mathcal{C}| - k$, then by the goodness of T

$$|\mathcal{C} \setminus (c'(N) \cup F(v))| \geq |\mathcal{C}| - |N| - |F(v)| > |\mathcal{C}| - (|\mathcal{C}| - k) - k = 0$$

so there would be a proper coloring of G respecting T , a contradiction. This proves Claim 3.2. \square

Now, Claims 3.1 and 3.2 together imply that $|G| > |\mathcal{C}| - k + 1 \geq 2k \geq k + 1$. Next, suppose that G is not $(k + 1)$ -connected, then there would be a cutset X of G with $|X| \leq k$ and with disjoint nonempty sets of vertices A, B with $A \cup B = V(G) \setminus X$ and no edges between them. Since cost is additive under disjoint unions, we have $\text{cost}(T_A) + \text{cost}(T_B) \leq \text{cost}(T) < 2k^2$; and so we may assume $\text{cost}(T_B) < k^2$. Let $D := A \cup X \cup S$; then $|D| < |G|$ since $B \not\subseteq S$ by Claim 3.1. Thus $G[D]$ is extensible, and so has a proper coloring c' respecting T_D . Let $T' := (S', c'', F|_{B \setminus S})$ be the template on $G[B \cup X]$ with $S' := X \cup (B \cap S)$ and c'' defined by $c''|_X := c'|_X$ and $c''|_{B \cap S} := c|_{B \cap S}$; note that c'' is a proper coloring of $G[S']$. Since

$$\text{cost}(T') = k|X \cup (B \cap S)| + \sum_{v \in B \setminus S} |F(v)| = k|X| + \text{cost}(T_B) < k^2 + k^2 = 2k^2,$$

and since $G[B \cup X]$ is extensible, it has a proper coloring c''' respecting T' . As $c'|_X = c''|_X = c'''|_X$, gluing c' and c''' would give a proper coloring of G respecting T , a contradiction. This proves Lemma 3.2. \blacksquare

4. CHROMATIC NUMBER

This section deals with the chromatic part of Theorem 1.2. We aim to prove that if $|\mathcal{C}|$ is sufficiently large then every inextensible graph has chromatic number as large as desired. To do so, we prove that if an inextensible graph G has small chromatic number, then we can find a proper coloring of G respecting a good template on G (this is similar to the approach in the proof of Lemma 3.2). For an integer $n \geq 0$, let $[n]$ be $\{1, 2, \dots, n\}$ if $n \geq 1$ and \emptyset if $n = 0$. Here is the main result of this section.

Lemma 4.1. *If $|\mathcal{C}| \geq (3 + \frac{1}{16})k - 1$, then every inextensible graph G satisfies $\chi(G) \geq |\mathcal{C}| - 2k + 3$.*

It is worth noting that the bound $|\mathcal{C}| - 2k + 3$ is optimal in Lemma 4.1, since it is false if we ask for $|\mathcal{C}| - 2k + 4$ under the weaker condition $|\mathcal{C}| \geq 2k - 1$: indeed, let $m := |\mathcal{C}| - 2k + 4 \geq 3$, let S be a stable set of cardinality $2k - 1$, let K be a complete graph on $m - 2$ vertices, and let $H_{k,m}$ be the graph obtained by joining every vertex in S to every vertex in K ; then $\chi(H_{k,m}) = m - 1$, but the template on $H_{k,m}$ with the vertices in S precolored differently and no forbidden colors at the vertices in K witnesses the inextensibility of $H_{k,m}$. However, we do not know whether Lemma 4.1 still holds when $|\mathcal{C}| \geq 3k - 1$, which was the condition on $|\mathcal{C}|$ needed to guarantee the $(k + 1)$ -connectivity of minimally inextensible graphs back in Section 3.

For the rest of this section we make use of the following setup. Let $|\mathcal{C}| \geq 3k - 1$, let G be an inextensible graph, let $\chi := \chi(G)$, and let $S_1 \cup \dots \cup S_\chi$ be a partition of $V(G)$ into χ stable sets. Let $T = (S, c, F)$ be a good template of G given by Lemma 3.1, that is, $|F(v)| < k$ for all $v \in V(G) \setminus S$. For every $P \subseteq V(G)$, let $w(P) := \sum_{v \in P} |F(v)|$ be the *weight* of P ; note that w is additive under disjoint unions. For every

$i \in [\chi]$, let $P_i := S_i \setminus S$ and $p_i := \lfloor w(P_i)/k \rfloor$; we may assume $P_i \neq \emptyset$, since we can add an isolated vertex to P_i if $P_i = \emptyset$. Let $p := p_1 + \dots + p_\chi$ and $t := 2k - |S|$, then

$$kt = 2k^2 - k|S| > \text{cost}(T) - k|S| \geq w(P_1) + \dots + w(P_\chi). \quad (1)$$

Thus, since $w(P_i) \geq kp_i$ for all $i \in [\chi]$, we obtain $kt > k(p_1 + \dots + p_\chi) = kp$ which yields

$$p \leq t - 1. \quad (2)$$

4.1. $4k - 1$ with a loss of $2k - 1$ suffices. To give a better explanation of the main argument, we give a quick proof of the bound $g(k, m) \leq \max(m + 2k - 1, 4k - 1)$ for all $k \geq 1$ and $m \geq 2$, whose chromatic part follows from the following weaker version of Lemma 4.1.

Lemma 4.2. *If $|\mathcal{C}| \geq 4k - 2$ then $\chi \geq |\mathcal{C}| - 2k + 2$.*

The proof of this lemma resembles the argument in [8], utilizing the following simple fact.

Lemma 4.3. *Let $k, q \geq 1$ be integers, and let Q be a set of integers (repeated members allowed) with $0 \leq a \leq k$ for all $a \in Q$ and $qk \leq \sum_{a \in A} a < (q + 1)k$. Then there exists a partition $Q = Q_1 \cup \dots \cup Q_q$ such that $\sum_{a \in Q_i} a < 2k$ for all $i \in [q]$.*

Proof. The lemma is trivial for $q = 1$. Let $q > 1$ and assume that it holds for $q - 1$. Let R be a maximal subset of Q with $\sum_{a \in R} a \geq (q - 1)k$, then $\sum_{a \in R} a < qk$ as $0 \leq a \leq k$ for all $a \in Q$; so by induction there is a partition $R = Q_1 \cup \dots \cup Q_{q-1}$ with $\sum_{a \in Q_i} a < 2k$ for all $i \in [q - 1]$. Then let $Q_q := Q \setminus R$ and note that $\sum_{a \in Q_q} a < (q + 1)k - (q - 1)k = 2k$. This proves Lemma 4.3. \blacksquare

Proof of Lemma 4.2. Let $I_0 := \{i \in [\chi] : p_i = 0\}$, let $I_1 := \{i \in [\chi] : p_i \geq 1\}$, and let $P := \bigcup_{i \in I_1} P_i$; note that $p = \sum_{i \in I_1} p_i$. By Lemma 4.3, for each $i \in I_1$ there is a partition $P_i = P_{i1} \cup \dots \cup P_{ip_i}$ such that $w(P_{ij}) < 2k$ for all $j \in [p_i]$. For every $i \in I_1$ and $j \in [p_i]$, let $L_{ij} := \mathcal{C} \setminus (c(S) \cup \bigcup_{v \in P_{ij}} F(v))$, then since $p \leq t - 1 = 2k - |S| - 1$ by (2) and since $|\mathcal{C}| \geq 4k - 2$, we see that

$$|L_{ij}| \geq |\mathcal{C}| - |S| - w(P_{ij}) \geq |\mathcal{C}| - (2k - p - 1) - (2k - 1) = |\mathcal{C}| - (4k - 2) + p \geq p.$$

Hence we can assign $p = \sum_{i \in I_1} p_i$ different colors to the stable sets in $\bigcup_{i \in I_1} \{P_{ij} : j \in [p_i]\}$, obtaining a proper coloring c' of $G[S \cup P]$ respecting $T_{S \cup P}$ with $|c'(P)| \leq p$. Put $S' := S \cup P$ and $t' := t - p \geq 1$; then

$$|c'(S')| \leq |S| + p = 2k - t + p = 2k - t'.$$

Now, suppose for a contradiction that $\chi \leq |\mathcal{C}| - 2k + 1$, and let $L_i := \mathcal{C} \setminus (c'(S') \cup \bigcup_{v \in P_i} F(v))$ for all $i \in I_0$. Let $I' := \{i \in I_0 : w(P_i) \geq t'\}$. Note that for every $i \in I' \setminus I_0$, we have that

$$|L_i| \geq |\mathcal{C}| - |c'(S')| - w(P_i) \geq |\mathcal{C}| - (2k - t') - (t' - 1) = |\mathcal{C}| - 2k + 1 \geq \chi \geq |I_0|.$$

Thus, if we can assign $|I'|$ different colors to the stable sets in $\{P_i : i \in I'\}$ such that each P_i gets a color in L_i , then we can assign $|I_0|$ different colors to the stable sets in $\{P_i : i \in I_0\}$ such that each P_i gets a color in L_i . This would give a proper coloring of G respecting T , a contradiction.

We now assign colors to $\{P_i : i \in I'\}$. We may assume that $I' \neq \emptyset$. Let $x_i := w(P_i)$ for all $i \in I'$, and assume that $I' = [n]$ and $x_1 \geq \dots \geq x_n$ for some $n \geq 1$. Since $p = \sum_{i \in I_1} p_i$, by (1) we see that for every $i \in [n]$,

$$kt > \sum_{i \in I_0} x_i + \sum_{i \in I_1} w(P_i) \geq ix_i + k \sum_{i \in I_1} p_i = ix_i + kp$$

so $i < kt'/x_i$. It follows that, since $x_i = w(P_i) \leq k$,

$$x_i + i < x_i + \frac{kt'}{x_i} = k + t' - \frac{(k - x_i)(x_i - t')}{x_i} \leq k + t'$$

so $x_i + i \leq k + t' - 1$, which yields (note that $|\mathcal{C}| \geq 4k - 2 \geq 3k - 1$)

$$|L_i| - i \geq |\mathcal{C}| - |c'(S')| - x_i - i \geq |\mathcal{C}| - (2k - t') - (k + t' - 1) = |\mathcal{C}| - (3k - 1) \geq 0.$$

Thus, $|L_i| \geq i$ for all $i \in [n]$, and so we can greedily assign $n = |I'|$ different colors to the stable sets P_1, \dots, P_n such that each P_i gets a color in L_i . This proves Lemma 4.2. \blacksquare

We can now give a proof that $g(k, m) \leq \max(m + 2k - 1, 4k - 1)$ for all $k \geq 1$ and $m \geq 2$.

Proposition 4.1. *For every integer $k \geq 1$, every graph G with $\chi(G) \geq 4k - 1$ contains a $(k + 1)$ -connected subgraph with more than $\chi(G) - k$ vertices and chromatic number at least $\chi(G) - 2k + 1$.*

Proof. Let \mathcal{C} be a set of $\chi(G) - 1$ colors, then G is \mathcal{C} -inextensible; so it has a minimally \mathcal{C} -inextensible subgraph H . Then H is $(k + 1)$ -connected and has more than $\chi(G) - k$ vertices by Lemma 3.2, and satisfies $\chi(H) \geq |\mathcal{C}| - 2k + 2 = \chi(G) - 2k + 1$ by Lemma 4.2. This proves Proposition 4.1. ■

4.2. Reduction step. The proof of Lemma 4.2 consists of two steps: coloring the stable sets P_i with $p_i \geq 1$ by p colors, then coloring those P_i with $p_i = 0$. As shown in the proof, the second step can be done smoothly under the condition $|\mathcal{C}| \geq 3k - 1$ as long as the first step has been carried out, which is possible when $|\mathcal{C}| \geq 4k - 2$ thanks to Lemma 4.3. In order to go below $4k - 2$ significantly, it might be tempting to improve the constant factor 2 in Lemma 4.3 to a (much) smaller constant independent of $q \geq 1$. This, however, is not possible even if one asks for a partition of Q into $q + r$ sets each with sum of elements less than $(2 - \varepsilon)k$, for any given integer $r \geq 0$ and any given small $\varepsilon > 0$.³ Another potential shortcoming of the proof of Lemma 4.2 is that the goodness of the template T was never really used, since the calculations involving every $|F(v)|$ only require them to be at most k . On the other hand, given the goodness of T , every partition of P_i into stable sets of weight less than k must have at least $p_i + 1$ components as $w(P_i) \geq p_i k$. In fact, we shall prove that under the condition $|\mathcal{C}| \geq 3k - 1$, it is possible to get around Lemma 4.3 and reduce Lemma 4.1 essentially to the case when each P_i has a partition into exactly $p_i + 1$ sets of weight less than k . More precisely, for each $i \in [\chi]$, there is a nonnegative integer $q_i \leq p_i$ and a subset of P_i of weight at least $q_i k$ whose vertices can be colored by at most q_i colors and whose complement in P_i can be partitioned into $p_i - q_i + 1$ stable sets each of weight less than k . This is the content of the following lemma.

Lemma 4.4. *If $|\mathcal{C}| \geq 3k - 1$, then there exist integers q_1, \dots, q_χ with $0 \leq q_i \leq p_i$ for all $i \in [\chi]$ and subsets P'_1, \dots, P'_χ of P_1, \dots, P_χ respectively, such that*

- for every $i \in [\chi]$, $w(P'_i) \geq q_i k$, and there is a partition $P_i \setminus P'_i = \bigcup_{j \in [t_i + 1]} P_{ij}$ with $w(P_{ij}) < k$ for all $j \in [t_i + 1]$ where $t_i := p_i - q_i$; and
- for $P^1 := \bigcup_{i \in [\chi]} P'_i$, there is a proper coloring c_1 of $G[S \cup P^1]$ respecting $T_{S \cup P^1}$ with $|c_1(P'_i)| \leq q_i$ for all $i \in [\chi]$, in particular $|c_1(P^1)| \leq q_1 + \dots + q_\chi$.

To prove Lemma 4.4, we need several notions. Given a subset P of $V(G)$, a sequence $\mathcal{Q} = (Q_1, \dots, Q_n)$ of nonempty disjoint subsets of P is called *fit* if $0 \leq w(Q_j) < k$ for all $j \in [n]$ and $P = Q_1 \cup \dots \cup Q_n$. Then every sequence $(v_1, \dots, v_{|P|})$ of vertices in P is a *fit sequence* of P since $w(\{v_j\}) = |F(v_j)| < k$ for all $j \in [n]$ by the goodness of T ; also, note that permuting the components of \mathcal{Q} preserves the fitness of the resulting sequence. A fit sequence $\mathcal{Q} = (Q_1, \dots, Q_n)$ is said to be *increasing* if $w(Q_1) \leq \dots \leq w(Q_n)$; observe that a fit sequence of \mathcal{Q} can be made increasing by sorting its components, and so P_i always has an increasing fit sequence.

Given a fit sequence $\mathcal{Q} = (Q_1, \dots, Q_n)$ of P , let $w_j(\mathcal{Q}) := w(Q_1 \cup \dots \cup Q_j)$ for all $j \in [n]$; then $w_n(\mathcal{Q}) = w(P)$. For $j \in [n]$, say that Q_j is a *jump* of \mathcal{Q} if $\lfloor w_j(\mathcal{Q})/k \rfloor = \lfloor w_{j-1}(\mathcal{Q})/k \rfloor + 1$, that is, there is an integer $q \geq 1$ such that

$$(q - 1)k \leq w_{j-1}(\mathcal{Q}) < qk \leq w_j(\mathcal{Q}) < (q + 1)k$$

and say that Q_j is a *non-jump* of \mathcal{Q} if $\lfloor w_j(\mathcal{Q})/k \rfloor = \lfloor w_{j-1}(\mathcal{Q})/k \rfloor$, that is, there is an integer $q \geq 0$ with

$$qk \leq w_{j-1}(\mathcal{Q}) \leq w_j(\mathcal{Q}) < (q + 1)k,$$

and in this case say that q is a *landmark* of \mathcal{Q} . Observe that, as $0 \leq w(Q_j) < k$ for all $j \in [n]$,

- Q_1 is always a non-jump, in particular 0 is always a landmark of \mathcal{Q} ;
- each Q_j is either a jump or a non-jump;
- there are $\lfloor w(P)/k \rfloor$ jumps in total; and

³To see this, let $\varepsilon > 0$ be sufficiently small, let k be much larger than $1/\varepsilon$, let q be such that $r/\varepsilon < (q + r)/2 < (r + 1)/\varepsilon$, and let Q be some set of $q + r + 1$ integers at least $(1 - \varepsilon/2)k$ and smaller than k .

- every landmark of \mathcal{Q} is at most $\lfloor w(P)/k \rfloor$.

A fit sequence of P is called *critical* if it has no two consecutive non-jumps. By choosing an increasing fit sequence of P then combining two consecutive non-jumps (if exist) and sorting the new fit sequence repeatedly, we see that P always has an increasing critical sequence. Also, given a critical sequence $\mathcal{Q} = (Q_1, \dots, Q_n)$ of P with landmarks q_1, \dots, q_ℓ satisfying $0 = q_1 < \dots < q_\ell \leq \lfloor w(P)/k \rfloor$, it is not hard to see that

- $\ell = n - \lfloor w(P)/k \rfloor$;
- for every $r \in [\ell - 1]$, there are exactly $q_{r+1} - q_r$ jumps of \mathcal{Q} between $q_r k$ and $(q_{r+1} + 1)k$, in particular $q_r k \leq w_{q_r+r}(\mathcal{Q}) < (q_r + 1)k$ for all $r \in [\ell]$; and
- there are exactly $\lfloor w(P)/k \rfloor - q_\ell$ jumps of \mathcal{Q} after $q_\ell k$.

Given the above definitions, the proof of Lemma 4.4 relies on an iterative procedure which recolors and swaps stable sets within fit sequences of P_1, \dots, P_χ , as follows.

Proof of Lemma 4.4. It suffices to iterate the following claim for $i = 1, 2, \dots, \chi$ in turn.

Claim 4.1. *Let $i \in [\chi]$, and assume that there exist integers q_1, \dots, q_{i-1} with $0 \leq q_h \leq p_h$ for all $h \in [i-1]$ and subsets P'_1, \dots, P'_{i-1} of P_1, \dots, P_{i-1} respectively, such that*

- for every $h \in [i-1]$, $w(P'_h) \geq q_h k$, and there is a partition $P_h \setminus P'_h = \bigcup_{j \in [t_h+1]} P_{hj}$ such that $w(P_{hj}) < k$ for all $j \in [t_h+1]$ where $t_h := p_h - q_h$; and
- for $P' := \bigcup_{h \in [i-1]} P'_h$, there is a proper coloring c' of $G[S \cup P']$ respecting $T_{S \cup P'}$ with $|c'(P'_h)| \leq q_h$ for all $h \in [i-1]$, in particular $|c'(P')| \leq q_1 + \dots + q_{i-1}$.

Then there exists an integer q_i with $0 \leq q_i \leq p_i$ and a subset P'_i of P_i such that for $t_i := p_i - q_i$,

- $w(P'_i) \geq q_i k$, and there is a partition $P_i \setminus P'_i = \bigcup_{j \in [t_i+1]} P_{ij}$ with $w(P_{ij}) < k$ for all $j \in [t_i+1]$; and
- we can color the vertices of P'_i by at most q_i colors such that every $v \in P'_i$ gets a color not in $c'(S \cup P') \cup F(v)$.

Proof. Let $\mathcal{Q} = (Q_1, \dots, Q_n)$ be an increasing critical sequence of P_i . Let q_1^i, \dots, q_ℓ^i with $0 = q_1^i < \dots < q_\ell^i \leq \lfloor w(P_i)/k \rfloor = p_i$ be the landmarks of \mathcal{Q} , and let $n_r := q_r^i + r$ for all $r \in [\ell]$; then

$$q_r^i \leq w_{n_r-1}(\mathcal{Q}) \leq w_{n_r}(\mathcal{Q}) < (q_r^i + 1)k$$

by the criticality of \mathcal{Q} . Let $\mathcal{Q}^1 := \mathcal{Q}$; and for $r \in [\ell - 1]$, assume that there is a fit sequence $\mathcal{Q}^r = (Q_1^r, \dots, Q_n^r)$ of P_i such that

- $(Q_1^r, \dots, Q_{n_r}^r)$ is a permutation of (Q_1, \dots, Q_{n_r}) ;
- $Q_j^r = Q_j$ for all $j \in [n] \setminus [n_r]$; and
- we have assigned at most q_r^i colors to the stable sets $Q_1^r, \dots, Q_{n_r-1}^r$ such that each $v \in Q_1^r \cup \dots \cup Q_{n_r-1}^r$ gets a color not in $c'(S \cup P') \cup F(v)$.

Let $L := \mathcal{C} \setminus c'(S \cup P')$. Because $q_h \leq p_h$ for all $h \in [i-1]$ and $p \leq t-1$ by (2), we see that

$$|c'(P')| \leq q_1 + \dots + q_{i-1} \leq p_1 + \dots + p_{i-1} \leq p \leq t-1.$$

Thus, as $|\mathcal{C}| \geq 3k - 1$, we obtain

$$|L| \geq |\mathcal{C}| - |c'(S \cup P')| \geq |\mathcal{C}| - |S| - |c'(P')| \geq |\mathcal{C}| - (2k - t) - (t - 1) = |\mathcal{C}| - (2k - 1) \geq k.$$

For $j \in \{n_r, n_r + 1, \dots, n_{r+1}\}$, let $L_j := L \setminus (\bigcup_{v \in Q_j^r} F(v))$; note that $w(Q_{n_r}^r) \leq w(Q_{n_r})$ since \mathcal{Q} is increasing, and $w(Q_j^r) = w(Q_j)$ for all $j \in [n_{r+1}] \setminus [n_r]$, in particular $|L_j| \geq |L| - w(Q_{n_r}^r) > k - k = 0$ and so L_j is nonempty for all $j \in \{n_r, n_r + 1, \dots, n_{r+1}\}$. Since q_r^i and q_{r+1}^i are landmarks of \mathcal{Q} , we see that

$$\sum_{j=n_r}^{n_{r+1}} w(Q_j^r) \leq \sum_{j=n_r}^{n_{r+1}} w(Q_j) = w_{n_{r+1}}(\mathcal{Q}) - w_{n_r-1}(\mathcal{Q}) < (q_{r+1}^i + 1)k - q_r^i k = (n_{r+1} - n_r)k.$$

Therefore, as $|L| \geq k$, we deduce that

$$\begin{aligned} \sum_{j=n_r}^{n_{r+1}} |L_j| &\geq \sum_{j=n_r}^{n_{r+1}} (|L| - w(Q_j^r)) \\ &> (n_{r+1} - n_r + 1)|L| - (n_{r+1} - n_r)k = (n_{r+1} - n_r)(|L| - k) + |L| \geq |L|. \end{aligned}$$

It follows that $L_{n_r}, \dots, L_{n_{r+1}}$ are not pairwise disjoint, thus there is a color common to two among them; and so, since they are all nonempty, we can assign at most $n_{r+1} - n_r = q_{r+1}^i - q_r^i + 1$ colors to the stable sets in $\{Q_j^r : j \in \{n_r, \dots, n_{r+1}\}\}$ such that each Q_j^r gets a color in L_j . Then

- if in fact at most $q_{r+1}^i - q_r^i$ colors are needed, then obviously we have assigned at most $q_{r+1}^i - q_r^i$ colors to the stable sets $Q_{n_r}^r, \dots, Q_{n_{r+1}-1}^r$ such that Q_j^r gets a color in L_j for all $j \in \{n_r, \dots, n_{r+1} - 1\}$; we then uncolor $Q_{n_{r+1}}^r$ and let $\mathcal{Q}^{r+1} := \mathcal{Q}^r$;
- if $q_{r+1}^i - q_r^i + 1$ colors are needed and $Q_{n_{r+1}}^r$ gets a color not already used for $Q_{n_r}^r, \dots, Q_{n_{r+1}-1}^r$, then again we have assigned at most $q_{r+1}^i - q_r^i$ colors to the stable sets $Q_{n_r}^r, \dots, Q_{n_{r+1}-1}^r$ such that Q_j^r gets a color in L_j for all $j \in \{n_r, \dots, n_{r+1} - 1\}$; we then uncolor $Q_{n_{r+1}}^r$ and let $\mathcal{Q}^{r+1} := \mathcal{Q}^r$; and
- if $q_{r+1}^i - q_r^i + 1$ colors are needed and we have to color $Q_{n_{r+1}}^r$ and Q_j^r with $j \in \{n_r, \dots, n_{r+1} - 1\}$ by the same color, then pick $j' \in \{n_r, \dots, n_{r+1} - 1\} \setminus \{j\}$ (this is possible as $n_{r+1} - n_r = q_{r+1}^i - q_r^i + 1 > 1$), uncolor $Q_{j'}^r$, and swap $Q_{j'}^r$ and $Q_{n_{r+1}}^r$ in \mathcal{Q}^r to obtain a new fit sequence \mathcal{Q}^{r+1} of P_i .

In this way, we have shown that there exists a fit sequence $\mathcal{Q}^{r+1} = (Q_1^{r+1}, \dots, Q_n^{r+1})$ of P_i such that

- $(Q_1^{r+1}, \dots, Q_{n_{r+1}}^{r+1})$ is a permutation of $(Q_1, \dots, Q_{n_{r+1}})$;
- $Q_j^{r+1} = Q_j$ for all $j \in [n] \setminus [n_{r+1}]$; and
- we have assigned at most $q_r^i + (q_{r+1}^i - q_r^i) = q_{r+1}^i$ colors to the stable sets $Q_1^{r+1}, \dots, Q_{n_{r+1}-1}^{r+1}$ such that each $v \in Q_1^{r+1} \cup \dots \cup Q_{n_{r+1}-1}^{r+1}$ gets a color not in $c'(S \cup P') \cup F(v)$.

Iterating this procedure for $r = 1, 2, \dots, \ell - 1$ in turn, we have shown that there is a fit sequence $\mathcal{Q}^\ell = (Q_1^\ell, \dots, Q_n^\ell)$ of P_i such that

- $(Q_1^\ell, \dots, Q_{n_\ell}^\ell)$ is a permutation of (Q_1, \dots, Q_{n_ℓ}) ;
- $Q_j^\ell = Q_j$ for all $j \in [n] \setminus [n_\ell]$; and
- we can assign at most q_ℓ^i colors to the stable sets $Q_1^\ell, \dots, Q_{n_\ell-1}^\ell$ such that each $v \in Q_1^\ell \cup \dots \cup Q_{n_\ell-1}^\ell$ gets a color not in $c'(S \cup P') \cup F(v)$.

Now, let $q_i := q_\ell^i$, $t_i := p_i - q_i = n - n_\ell$, $P_i' := \bigcup_{j \in [n_\ell-1]} Q_j^\ell$, and $P_{ij} := Q_{j+n_\ell-1}^\ell$ for all $j \in [t_i + 1]$; then

- $w(Q_{n_\ell}^\ell) \leq w(Q_{n_\ell})$ since \mathcal{Q} is increasing, and $w_{n_\ell}(\mathcal{Q}^\ell) = w_{n_\ell}(\mathcal{Q})$ since $(Q_1^\ell, \dots, Q_{n_\ell}^\ell)$ is a permutation of (Q_1, \dots, Q_{n_ℓ}) , therefore

$$w(P_i') = w_{n_\ell-1}(\mathcal{Q}^\ell) = w_{n_\ell}(\mathcal{Q}^\ell) - w(Q_{n_\ell}^\ell) \geq w_{n_\ell}(\mathcal{Q}) - w(Q_{n_\ell}) = w_{n_\ell-1}(\mathcal{Q}) \geq q_\ell^i k = q_i k;$$

- because there are exactly $\lfloor w(P_i)/k \rfloor - q_i = t_i$ jumps of \mathcal{Q} after $q_i k$, $\bigcup_{j \in [t_i+1]} P_{ij}$ is a partition of $P_i \setminus P_i'$ with $w(P_{ij}) = w(Q_{j+n_\ell-1}^\ell) = w(Q_{j+n_\ell-1}) < k$ for all $j \in [t_i + 1]$; and
- we can color the vertices of P_i' by at most $q_\ell^i = q_i$ colors such that every $v \in P_i'$ gets a color not in $c'(S \cup P') \cup F(v)$. This proves Claim 4.1. \square

The proof of Lemma 4.4 is complete. \blacksquare

4.3. Finishing the proof. We now come to the rest of the proof of Lemma 4.1. As explained in the previous subsection, Lemma 4.4 allows us to have in mind that each P_i has a partition into $p_i + 1$ stable sets of weight less than k ; but for technical reason (there might be some i with $w(P_i \setminus P_i') < t_i k$) it might be wise to present the full argument without assuming that. We shall go through several coloring steps. To achieve the bound $\chi \geq |\mathcal{C}| - 2k + 3$ which is optimal as we have seen, we observe that each vertex in $P_i \setminus P_i'$ (which is yet to be colored) need not get a color not already used for $(S \cap S_i) \cup P_i'$. Then, careful calculations with the supposition $\chi \leq |\mathcal{C}| - 2k + 2$ (for a contradiction) will enable us to basically work with the supposition $\chi \leq |\mathcal{C}| - 2k + 1$ as in the second step in the proof of Lemma 4.2 (with more twists).

We obtain the constant factor $3 + \frac{1}{16}$ from the estimate in the last coloring step. Let us proceed with the details.

Proof of Lemma 4.1. Picking up where we left off, we let $P'_1, \dots, P'_\chi, P^1, q_1, \dots, q_\chi, t_1, \dots, t_\chi$, and c_1 be given by Lemma 4.4. Let $S^1 := S \cup P^1$, and let $T^1 := (S^1, c_1, F|_{V(G) \setminus S^1})$ be a template on G . Let $q := q_1 + \dots + q_\chi$, then for every $i \in [\chi]$,

$$|c_1(S^1 \setminus S_i)| \leq |c(S \setminus S_i)| + \sum_{h \in [\chi] \setminus \{i\}} |c_1(P'_h)| \leq |S \setminus S_i| + \sum_{h \in [\chi] \setminus \{i\}} q_h = |S| - |S \cap S_i| + q - q_i. \quad (3)$$

For every $i \in [\chi]$, let $X_i := P_i \setminus P'_i = \bigcup_{j \in [t_i+1]} P_{ij}$ and $x_i := w(X_i) - t_i k$; then $x_i < k$, and since $w(P'_i) \geq q_i k$, we have

$$w(P_i) - x_i = w(P_i) - w(X_i) + t_i k = w(P'_i) + t_i k \geq (q_i + t_i)k = p_i k.$$

Let $t' := t - p$, then $t' \geq 1$ by (2), and

$$kt > \text{cost}(T) - k|S| = \sum_{i \in [\chi]} w(P_i) \geq \sum_{i \in [\chi]} (p_i k + x_i) = pk + \sum_{i \in [\chi]} x_i$$

by (1), which implies

$$k(t - p) = kt' > \sum_{i \in [\chi]} x_i. \quad (4)$$

Now, let $I_0 := \{i \in [\chi] : t_i = 0\}$; by adding isolated vertices we may assume $X_i \neq \emptyset$ for all $i \in I_0$. Let I_2 be a maximal subset of $[\chi]$ such that we can properly color $\bigcup_{i \in I_2} X_i$ by $\sum_{i \in I_2} t_i$ colors so that for every $i \in I_2$, each $v \in X_i$ gets a color not in $c_1(S^1 \setminus S_i) \cup F(v)$. Then I_0, I_2 are disjoint. Let $I_1 := [\chi] \setminus (I_0 \cup I_2)$, and let $s_1 := \sum_{i \in I_1} t_i$ and $s_2 := \sum_{i \in I_2} t_i$. Let $P^2 := \bigcup_{i \in I_2} X_i$, and let c_2 be a proper coloring of $G[S^1 \cup P^2]$ respecting $T_{S^1 \cup P^2}^1$ with $|c_2(P^2)| \leq s_2$ whose existence is guaranteed by the definition of I_2 . Observe that

$$s_1 + s_2 = \sum_{i \in I_1} t_i + \sum_{i \in I_2} t_i = \sum_{i \in [\chi]} t_i = \sum_{i \in [\chi]} (p_i - q_i) = p - q.$$

Let $S^2 := S^1 \cup P^2$, and let $T^2 := (S^2, c_2, F|_{V(G) \setminus S^2})$ be a template on G . Let $I := I_0 \cup I_1$, then $I \cup I_2$ is a partition of $[\chi]$. For every $i \in I$, as $|c_1(S^1 \setminus S_i)| \leq |S| - |S \cap S_i| + q - q_i$ by (3) and $s_1 + s_2 = p - q$,

$$\begin{aligned} |c_2(S^2 \setminus S_i)| &\leq |c_1(S^1 \setminus S_i)| + |c_2(P^2)| \leq |S| - |S \cap S_i| + q - q_i + s_2 \\ &= (2k - t' - p) - |S \cap S_i| + q - q_i + (p - q - s_1) \\ &= 2k - t' - s_1 - |S \cap S_i| - q_i. \end{aligned} \quad (5)$$

Claim 4.2. $x_i \geq t_i(|C| - 3k + t' + s_1 + |S \cap S_i| + q_i)$ for all $i \in I$.

Proof. Suppose that there is $i \in I$ such that $x_i < t_i(|C| - 3k + t' + s_1 + |S \cap S_i| + q_i)$; then $i \in I_1$. Let $L := C \setminus c_2(S^2 \setminus S_i)$. Since $|c_2(S^2 \setminus S_i)| \leq 2k - t' - s_1 - |S \cap S_i| - q_i$ by (5), we see that

$$\begin{aligned} |L| &\geq |C| - |c_2(S^2 \setminus S_i)| \\ &\geq |C| - (2k - t' - s_1 - |S \cap S_i| - q_i) = |C| - 2k + t' + s_1 + |S \cap S_i| + q_i, \end{aligned}$$

therefore $|L| \geq |C| - 2k + 1$ since $t' \geq 1$. For every $j \in [t_i + 1]$, let $L_j := L \setminus (\bigcup_{v \in P_{ij}} F(v))$, then

$$|L_j| \geq |L| - w(P_{ij}) \geq |C| - 2k + 1 - (k - 1) = |C| - (3k - 2) > 0.$$

As $p_i = q_i + t_i$ and by our supposition on x_i , we have that

$$\begin{aligned} \sum_{j \in [t_i+1]} |L_j| &\geq (t_i + 1)|L| - \sum_{j \in [t_i+1]} w(P_{ij}) \geq |L| + t_i(|C| - 2k + t' + s_1 + |S \cap S_i| + q_i) - (x_i + t_i k) \\ &= |L| + t_i(|C| - 3k + t' + s_1 + |S \cap S_i| + q_i) - x_i > |L| \end{aligned}$$

so L_1, \dots, L_{t_i+1} are nonempty and not pairwise disjoint. Therefore we can assign at most t_i colors to the stable sets in $\{P_{ij} : j \in [t_i + 1]\}$ so that each P_{ij} gets a color in L_j , and so $I_2 \cup \{i\}$ contradicts the maximality of I_2 . This proves Claim 4.2. \square

For each $i \in I$, assume $w(P_{i1}) = \min_{j \in [t_i+1]} w(P_{ij}) =: y_i$, and let $Y_i := P_{i1}$; then $y_i = w(Y_i)$. Let $P^3 := \bigcup_{i \in I_1} (X_i \setminus Y_i)$; we next extend c_2 to P^3 by at most s_1 colors, as follows.

Claim 4.3. *There is a proper coloring c_3 of $G[S^2 \cup P^3]$ respecting $T_{S^2 \cup P^3}^2$ and satisfying $|c_3(X_i \setminus Y_i)| \leq t_i$ for all $i \in I_1$.*

Proof. Note that for every $i \in I_1$, $P^3 \cap S_i = X_i \setminus Y_i$ is the disjoint union of t_i stable sets of the form P_{ij} where $j \in [t_i+1] \setminus \{1\}$. For $i \in I_1$ and $j \in [t_i+1] \setminus \{1\}$, let $L_{ij} := \mathcal{C} \setminus (c_2(S^2 \setminus S_i) \cup \bigcup_{v \in P_{ij}} F(v))$, then as $|c_2(S^2 \setminus S_i)| \leq 2k - t' - s_1$ by (5), as $p_i = q_i + t_i$, and as $|\mathcal{C}| \geq 3k - 1$, we have

$$\begin{aligned} |L_{ij}| &\geq |\mathcal{C}| - |c_2(S^2 \setminus S_i)| - w(P_{ij}) \\ &\geq |\mathcal{C}| - (2k - t' - s_1) - (k - 1) = |\mathcal{C}| - (3k - 1) + t' + s_1 \geq s_1. \end{aligned}$$

Hence we can assign $\sum_{t \in I_1} t_i$ different colors to the stable sets in $\bigcup_{i \in I_1} \{P_{ij} : j \in [t_i+1] \setminus \{1\}\}$ such that each P_{ij} gets a color in L_{ij} . This proves Claim 4.3. \square

Let $S^3 := S^2 \cup P^3$, and let $T^3 := (S^3, c_3, F|_{V(G) \setminus S^3})$ be the template on G with c_3 given by Claim 4.3. For $i \in I$, since $|c_3(P^3 \setminus S_i)| \leq \sum_{h \in I_1 \setminus \{i\}} |c_3(X_h \setminus Y_h)| \leq \sum_{h \in I_1 \setminus \{i\}} t_h = s_1 - t_i$, since $|c_2(S^2 \setminus S_i)| \leq 2k - t - |S \cap S_i| - q_i - s_1$ by (5), and since $p_i = q_i + t_i$, we have that

$$\begin{aligned} |c_3(S^3 \setminus S_i)| &\leq |c_2(S^2 \setminus S_i)| + |c_3(P^3 \setminus S_i)| \\ &\leq (2k - t' - |S \cap S_i| - q_i - s_1) + (s_1 - t_i) = 2k - t' - |S \cap S_i| - p_i. \end{aligned} \tag{6}$$

So far our arguments only concern $|\mathcal{C}|$ instead of its relation to χ . Now, assume for a contradiction that $\chi \leq |\mathcal{C}| - 2k + 2$. Let $L_i := \mathcal{C} \setminus (c_3(S^3 \setminus S_i) \cup \bigcup_{v \in Y_i} F(v))$ for all $i \in I$. As promised, the following claim will enable us to basically work with the supposition $\chi \leq |\mathcal{C}| - 2k + 1$.

Claim 4.4. *There exists $J \subseteq I$ with $I_1 \subseteq J$, $|J| \leq |\mathcal{C}| - 2k + 1$, and $|L_i| \geq |I|$ for all $i \in I \setminus J$.*

Proof. If $|I| \leq |\mathcal{C}| - 2k + 1$ then we can just take J to be I ; so we may assume $|I| \geq |\mathcal{C}| - 2k + 2$, then $|I| = \chi = |\mathcal{C}| - 2k + 2$, $I = [\chi]$, and $I_2 = \emptyset$. It suffices to show that there exists $i \in I_0$ with $|L_i| \geq |I|$; then we can take $J := I \setminus \{i\}$. Suppose not; then $|L_i| < |I| = |\mathcal{C}| - 2k + 2$ for all $i \in I_0$, so since $x_i = y_i$ for $i \in I_0$ and $|c_3(S^3 \setminus S_i)| \leq 2k - t' - |S \cap S_i| - p_i$ by (6),

$$\begin{aligned} x_i = y_i &\geq |\mathcal{C}| - |c_3(S^3 \setminus S_i)| - |L_i| > |\mathcal{C}| - (2k - t' - |S \cap S_i| - p_i) - (|\mathcal{C}| - 2k + 2) \\ &= t' - 2 + |S \cap S_i| + p_i, \end{aligned}$$

hence $x_i \geq t' - 1 + |S \cap S_i| + p_i$ for all $i \in I_0$. For each $i \in I_1$, since $|\mathcal{C}| \geq 3k - 1$, since $s_1 + q_i \geq t_i + q_i = p_i$, and since $t_i \geq 1$, by Claim 4.2 we have that

$$x_i \geq t_i(|\mathcal{C}| - 3k + t' + s_1 + |S \cap S_i| + q_i) \geq t' - 1 + |S \cap S_i| + p_i.$$

Hence, since $|I| = \chi = |\mathcal{C}| - 2k + 2 \geq k + 1$ and $\sum_{i \in I} |S \cap S_i| = |S| = 2k - t' - p$, we obtain

$$\begin{aligned} \sum_{i \in I} x_i &\geq \sum_{i \in I} (t' - 1 + |S \cap S_i| + p_i) \\ &= |I|(t' - 1) + |S| + p \geq (k + 1)(t' - 1) + (2k - t' - p) + p = kt' + k - 1 \geq kt', \end{aligned}$$

contradicting (4). This proves Claim 4.4. \square

Let J be the subset of I obtained from Claim 4.4, and let $I' := \{i \in J : x_i \geq t'\}$; then $I_1 \subseteq I'$ as $x_i \geq t_i(|\mathcal{C}| - 3k + t' + s_1) \geq t'$ for all $i \in I_1$ by Claim 4.2 since $s_1 \geq t_i \geq 1$ and $|\mathcal{C}| \geq 3k - 1$. The following claim is sufficient to complete the proof of Lemma 4.1.

Claim 4.5. *It is possible to assign $|I'|$ different colors to the stable sets $\{Y_i : i \in I'\}$ so that each Y_i gets a color in L_i .*

To see how Claim 4.5 leads to a contradiction and finishes the proof of Lemma 4.1, observe that for every $i \in J \setminus I' \subseteq I_0$, since $y_i = x_i \leq t' - 1$ and $|c_3(S^3 \setminus S_i)| \leq 2k - t'$ by (6),

$$|L_i| \geq |\mathcal{C}| - |c_3(S^3 \setminus S_i)| - y_i \geq |\mathcal{C}| - (2k - t') - (t' - 1) = |\mathcal{C}| - 2k + 1 \geq |J|,$$

so we can assign $|J|$ different colors to the stable sets in $\{Y_i : i \in J\}$ so that each Y_i gets a color in L_i . Thus, since $|L_i| \geq |I|$ for all $i \in I \setminus J$, we can assign $|I|$ different colors to the stable sets in $\{Y_i : i \in I\}$ so that each Y_i gets a color in L_i . In this way we would obtain a proper coloring of G respecting T , a contradiction.

Hence, the rest of the proof will be devoted to Claim 4.5. We may assume that $I' \neq \emptyset$. Put $d := |\mathcal{C}| - 3k + 1 \geq \frac{1}{16}k$ and $d' := d + s_1 - 1$; then Claim 4.2 yields $x_i \geq t_i(d + t' + s_1 - 1) = t_i(t' + d')$ for all $i \in I_1$. For every $i \in I'$, as $|c_3(S^3 \setminus S_i)| \leq 2k - t'$ by (6), we see that

$$|L_i| - i \geq |\mathcal{C}| - |c_3(S^3 \setminus S_i)| - y_i - i \geq |\mathcal{C}| - (2k - t') - (y_i + i) = d + k + t' - (y_i + i) - 1. \quad (7)$$

Claim 4.6. *We may assume that $t' \geq d' + 1$.*

Proof. Suppose that $t' \leq d'$, then $x_i \geq t_i(t' + d') \geq t_i(t_i + 1)t' \geq (t_i + 1)t'$ for all $i \in I'$. Assume that $I' = [n]$ for some $n \geq 1$, and that $x_1 \geq \dots \geq x_n$. For every $i \in [n]$, since $t_i k + x_i = \sum_{j \in [t_i + 1]} w(P_{ij}) \geq (t_i + 1)y_i$, and since $i < kt'/x_i$ which follows from $kt' > x_1 + \dots + x_n \geq ix_i$ by (2), we have that

$$y_i + i < \frac{t_i k + x_i}{t_i + 1} + \frac{kt'}{x_i} = k + t' + \frac{(k - x_i)((t_i + 1)t' - x_i)}{x_i(t_i + 1)} \leq k + t'$$

and so $y_i + i \leq k + t' - 1$; thus (7) yields

$$|L_i| - i \geq d + k + t' - (k + t' - 1) - 1 = d \geq 0.$$

It follows that we can greedily assign $n = |I'|$ different colors to the stable sets Y_1, \dots, Y_n so that each Y_i gets a color in L_i , proving Claim 4.5. Thus we may assume that $t' \geq d' + 1$, as claimed. \square

Now, put $s := s_1 - 1$; then $d' = d + s$ and $t' \geq d' + 1 \geq s + 1 \geq |I_1|$ by Claim 4.6.

Claim 4.7. *We may assume that $s > d + s(t' + d')/k$.*

Proof. Suppose that $s \leq d + s(t' + d')/k$. For every $i \in I_1$, by (7) and since $t' \geq |I_1|$, we have that

$$|L_i| \geq d + k + t' - y_i - 1 \geq d + k + t' - (k - 1) - 1 = d + t' \geq |I_1|,$$

so we can assign $|I_1|$ different colors to the stable sets in $\{Y_i : i \in I_1\}$ so that each Y_i gets a color in L_i .

Now, assume that $I' \setminus I_1 = [n] \subseteq I_0$ for some $n \geq 0$, and that $x_1 \geq \dots \geq x_n$. If $n = 0$ then Claim 4.5 is proved, so we may assume that $n \geq 1$. Since $x_i \geq t_i(t' + d')$ for all $i \in I_1$, (4) implies that

$$kt' > \sum_{i \in I_1} t_i(t' + d') + \sum_{i \in [n]} x_i = s_1(t' + d') + \sum_{i \in [n]} x_i.$$

Thus, for every $i \in [n]$, $i < (kt' - s_1(t' + d'))/x_i$, and so $y_i + i = x_i + i < x_i + \frac{kt' - s_1(t' + d')}{x_i} = \phi(x_i)$ where $\phi: [t', \infty) \rightarrow \mathbb{R}$ is defined by $\phi(x) := x + \frac{kt' - s_1(t' + d')}{x}$. Since ϕ is convex and $t' \leq x_i \leq k - 1$, we have that $\phi(x_i) \leq \max(\phi(t'), \phi(k - 1))$. Observe that $\phi(t') = k + t' - s_1 - s_1 d'/t' < k + t' - s_1$, and that

$$\begin{aligned} \phi(k - 1) &= k - 1 + t' - \frac{(s_1 - 1)t' + s_1 d'}{k - 1} \\ &< k - 1 + t' - \frac{s}{k}(t' + d') \leq k - 1 + t' - s + d = k + t' - s_1 + d \end{aligned}$$

where the last inequality follows from our supposition. Thus for every $i \in [n]$, $\phi(x_i) < k + t' - s_1 + d$ which yields $y_i + i \leq k + t' - s_1 + d - 1$, and so (7) implies that

$$|L_i| - i \geq d + k + t' - (y_i + i) - 1 \geq d + k + t' - (k + t' - s_1 + d - 1) - 1 = s_1 \geq |I_1|.$$

Hence $|L_i| \geq i + |I_1|$ for all $i \in [n]$, and so we can greedily assign $n = |I' \setminus I_1|$ colors to the stable sets Y_1, \dots, Y_n such that each Y_i gets a color in L_i not already used for any stable set in $\{Y_{i'} : i' \in I_1\}$. Since this proves Claim 4.5, we may thus assume that $s > d + s(t' + d')/k$, as claimed. \square

The last part of the proof reduces to proving that

$$\frac{(k - (t' + d'))(t' - d')}{t' + d'} \leq 2d. \quad (8)$$

To see how (8) concludes the proof of Claim 4.5 and thus the proof of Lemma 4.1, assume that $I' = [n]$ for some $n \geq 1$ and $x_1 \geq \dots \geq x_n$. Then as in the proof of Claim 4.6, we have that for every $i \in [n]$,

$$y_i + i < k + t' + \frac{(k - x_i)((t_i + 1)t' - x_i)}{x_i(t_i + 1)}.$$

If $x_i \geq (t_i + 1)t'$ then $y_i + i < k + t'$. If $x_i \leq (t_i + 1)t' - 1$, then $i \in I_1$ which yields $t_i \geq 1$; thus since $k - 1 \geq x_i \geq t_i(t' + d') \geq t' + d'$ and $t' \geq d' + 1$ by Claim 4.6, by (8) we would obtain

$$\frac{(k - x_i)((t_i + 1)t' - x_i)}{x_i(t_i + 1)} \leq \frac{(k - x_i)(t' - d')}{x_i(t_i + 1)} \leq \frac{(k - (t' + d'))(t' - d')}{2(t' + d')} \leq d$$

so $y_i + i < k + t' + d$. We would have thus shown that $y_i + i \leq k + t' + d - 1$ for all $i \in [n]$, so (7) implies

$$|L_i| - i \geq d + k + t' - (k + t' + d - 1) - 1 = 0,$$

hence we could greedily assign $n = |I'|$ different colors to the stable sets Y_1, \dots, Y_n such that each Y_i gets a color in L_i , proving Claim 4.5.

Now, to prove (8), observe that

$$\begin{aligned} \frac{(k - (t' + d'))(t' - d')}{t' + d'} &= \frac{(k - (t' + d'))((t' + d') - 2d')}{t' + d'} \\ &= k + 2d' - \frac{2kd'}{t' + d'} - (t' + d') \leq k + 2d' - 2\sqrt{2kd'} = (\sqrt{k} - \sqrt{2d'})^2 \end{aligned}$$

where the inequality follows from the arithmetic mean–geometric mean inequality. If $d' \geq 4d$, then since $d \geq \frac{1}{16}k$ we would have $d' \geq \frac{1}{4}k$ so $(\sqrt{k} - \sqrt{2d'})^2 \leq (1 - \frac{1}{\sqrt{2}})^2 k < 2d$ and (8) is proved. So we may assume $d' < 4d$ which yields $s < 3d$ and so $\frac{4}{3}s < d'$ (recall that $d' = d + s$).

By Claim 4.7, we have that $s > d$ and $t' + d' < (1 - d/s)k$. We claim that $(1 - d/s)k < \sqrt{2kd'}$. Indeed, this is equivalent to $d/s + \sqrt{2d'/k} > 1$, which is true since

$$\frac{d}{s} + \sqrt{\frac{2d'}{k}} = \frac{d}{s} + \sqrt{\frac{d'}{2k}} + \sqrt{\frac{d'}{2k}} \geq 3\sqrt[3]{\frac{dd'}{2sk}} > 1$$

where the penultimate inequality follows from the arithmetic mean–geometric mean inequality, and the last inequality is true because $d \geq \frac{1}{16}k$ and $d' > \frac{4}{3}s$.

Consequently, $t' + d' < (1 - d/s)k < \sqrt{2kd'}$. Thus, because the function $x \mapsto 2kd'/x + x$ is strictly decreasing over $(0, \sqrt{2kd'})$, we deduce that

$$\frac{2kd'}{t' + d'} + t' + d' > \frac{2sd'}{s - d} + \left(1 - \frac{d}{s}\right)k = k + 2d' + \frac{2dd'}{s - d} - \frac{kd}{s} = k + 2d' - 2d + \frac{4sd}{s - d} - \frac{kd}{s},$$

and so

$$\frac{(k - (t' + d'))(t' - d')}{t' + d'} = k + 2d' - \frac{2kd'}{t' + d'} - (t' + d') < 2d - \frac{4ds}{s - d} + \frac{kd}{s} = 2d - \frac{d(4s^2 - k(s - d))}{s(s - d)}.$$

To complete the proof, it is enough to show that $4s^2 - k(s - d) \geq 0$ which is equivalent to $4(s - \frac{1}{8}k)^2 + k(d - \frac{1}{16}k) \geq 0$. This proves (8) and Lemma 4.1. \blacksquare

We are now ready to finish the proof of Theorem 1.2.

Proof of Theorem 1.2. Let $|\mathcal{C}|$ be a set of $\chi(G) - 1$ colors, then G is \mathcal{C} -inextensible; so there is a minimally \mathcal{C} -inextensible subgraph H of G . Then H is $(k + 1)$ -connected and has more than $\chi(G) - k$ vertices by Lemma 3.2, and satisfies $\chi(H) \geq |\mathcal{C}| - 2k + 3 = \chi(G) - 2k + 2$ by Lemma 4.1. This proves Theorem 1.2. \blacksquare

5. CONCLUDING REMARKS

There are some final points we would like to make. First, as remarked in the Introduction, we conjecture that the constant $\frac{1}{16}$ in Theorem 1.2 can be removed, which would yield $g(k, m) \leq \max(m + 2k - 2, 3k)$ for all $k \geq 1$ and $m \geq 2$; a possible way to approach this conjecture is to prove Lemma 4.1 when $|\mathcal{C}| \geq 3k - 1$. Note that one can simplify the proof of Lemma 4.1 to get $\chi(G) \geq |\mathcal{C}| - 3k + 4$ for every inextensible graph G whenever $|\mathcal{C}| \geq 3k - 1$, which together with Lemma 3.2 implies that $g(k, m) \leq m + 3k - 3$ for all $k \geq 1$ and $m \geq 3$.

Second, there is an analogue of $g(k, m)$ for list colorings as well (see [7, Section 5.4] for preliminaries). For integers $k \geq 1$ and $m \geq 2$, let $g_\ell(k, m)$ be the least integer $n \geq 1$ such that every graph with choosability at least n contains a $(k + 1)$ -connected subgraph with choosability at least m . It was proved in [8] that $g_\ell(k - 1, k) \leq 4k$ for all $k \geq 2$; by modifying the argument outlined there and adapting the notions of templates and inextensibility in Section 2 to the list coloring setting, one could show the following which yields $g_\ell(k, m) \leq m + 3k - 2$ for all $k \geq 1$ and $m \geq 2$.

Proposition 5.1. *For every integer $k \geq 1$, every graph G with $\chi_\ell(G) \geq 3k$ contains a $(k + 1)$ -connected subgraph with more than $\chi_\ell(G) - k$ vertices and choosability at least $\chi_\ell(G) - 3k + 2$.*

Note that the lower bound construction in [1] also shows that $g_\ell(k, m) \geq \max(m + k - 1, 2k + 1)$ for all $k \geq 1$ and $m \geq 2$. It would be interesting to narrow the gap between the lower and upper bounds on $g_\ell(k, m)$, for instance to decide whether $g_\ell(k, m) \leq \max(m + 2k, Ck)$ for some universal constant $C > 0$, similar to Theorem 1.2.

Finally, given the bounds $\max(m + k - 1, 2k + 1) \leq g(k, m) \leq \max(m + 2k - 2, \lceil (3 + \frac{1}{16})k \rceil)$ for all $k \geq 1$ and $m \geq 2$, it would be nice to prove that $g(k, m) \leq \max(m + (1 + \varepsilon)k, Ck)$ for some universal constants $\varepsilon \in (0, 1)$ and $C > 0$. New methods would be needed to accomplish this. Indeed, a loss of $2k - 2$ on the chromatic number seems to be the best that the “template-inextensibility” method could produce, because of the optimality of the bound $|\mathcal{C}| - 2k + 3$ in Lemma 4.1 as discussed in Section 4. When $k = 2$, the above bounds only yield $g(2, m) \in \{m + 1, m + 2\}$ for all $m \geq 5$, but Alex Scott and Paul Seymour (personal communication) recently showed that $g(2, m) = m + 1$ for all $m \geq 4$ by proving the following equivalent statement which was conjectured by the author.

Theorem 5.2. *For every integer $m \geq 4$ and for every graph G with chromatic number at least m , if there are nonadjacent vertices u, v in G for which every optimal coloring of G assigns different colors to u, v , then G contains a 3-connected subgraph with chromatic number at least m .*

(Here an *optimal coloring* of G is a proper coloring of G using $\chi(G)$ colors.) To see how Theorem 5.2 is equivalent to the statement that $g(2, m) = m + 1$ for $m \geq 4$, first assume that Theorem 5.2 holds. Let G be a graph with $\chi(G) = m + 1$; we may assume that every subgraph of G with fewer edges than G has chromatic number at most m . Suppose that G has no 3-connected subgraphs; then G would have a cutset $S = \{u, v\}$ with u, v nonadjacent in G and two nonempty vertex sets A, B with $A \cup B = V(G) \setminus S$ such that there are no edges between A and B . By [3, Theorem 14.9], there would be $H \in \{G[S \cup A], G[S \cup B]\}$ such that $\chi(H) = m$ and every optimal coloring of H assigns to u, v different colors; so H would be a counterexample to Theorem 5.2, a contradiction. Hence G is 3-connected and so $g(2, m) = m + 1$ for all $m \geq 4$.

Now, assume that there exists a counterexample G to Theorem 5.2 with two special vertices u, v . Let G_1, G_2, G_3 be disjoint copies of G with u_i, v_i being the respective copies of u, v for $i = 1, 2, 3$, and let $H_{2,m}$ be the graph constructed in the discussion on Lemma 4.1 with stable set $S = \{x_1, x_2, x_3\}$; then for $i = 1, 2, 3$ identify u_i with x_i and v_i with x_{i+1} (here $x_1 = x_4$), and let H be the resulting graph. If $\chi(H) = m$, then $\chi(G) = m$, and any optimal coloring of H would color u_i, v_i differently for each $i \in \{1, 2, 3\}$ which implies that x_1, x_2, x_3 would receive different colors; but this is a contradiction since each of them is adjacent to the vertices in $H_{2,m} \setminus S$ which is a complete graph on $m - 2$ vertices. This shows that $\chi(H) = m + 1$. Moreover, it is not hard to see that H has no 3-connected subgraphs with chromatic number at least m , which yields $g(2, m) = m + 2$.

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