

# MAT 540 : Problem Set 1

Due Thursday, September 19

1.

- (a). (2 points) In the category **Set**, show that a morphism is a monomorphism (resp. an epimorphism) if and only if it is injective (resp. surjective).
- (b). (2 points) Let  $\mathcal{C}$  be a category and  $F : \mathcal{C} \rightarrow \mathbf{Set}$  be a *faithful* functor, show that any morphism  $f$  of  $\mathcal{C}$  whose such that  $F(f)$  is injective (resp. surjective) is a monomorphism (resp. an epimorphism).
- (c). (2 points) What are the monomorphisms and epimorphisms in  ${}_R\mathbf{Mod}$  ?
- (d). (2 points) What are the monomorphisms in **Top** ? Give an example of a continuous morphism with dense image that is not an epimorphism in **Top**.<sup>1</sup>
- (e). (2 points) Find a category  $\mathcal{C}$ , a faithful  $F : \mathcal{C} \rightarrow \mathbf{Set}$  and a monomorphism  $f$  in  $\mathcal{C}$  such that  $F(f)$  is not injective.
- (f). (1 point) Find an epimorphism in **Ring** that is not surjective.
- (g). The goal of this question is to show that any epimorphism in **Grp** is a surjective map. Let  $\phi : G \rightarrow H$  be a morphism of groups, and suppose that it is an epimorphism in **Grp**. Let  $A = \text{Im}(\phi)$ . Let  $S = \{*\} \sqcup (H/A)$ , where  $\{*\}$  is a singleton, and let  $\mathfrak{S}$  be the group of permutations of  $S$ . We denote by  $\sigma$  the element of  $\mathfrak{S}$  that switches  $*$  and  $A$  and leaves the other elements of  $H/A$  fixed. For every  $h \in H$ , we denote by  $\psi_1(h)$  the element of  $\mathfrak{S}$  that leaves  $*$  fixed and acts on  $H/A$  by left translation by  $H$ ; this defines a morphism of groups  $\psi_1 : H \rightarrow \mathfrak{S}$ . We denote by  $\psi_2 : H \rightarrow \mathfrak{S}$  the morphism  $\sigma\psi_1\sigma^{-1}$ .
  - (i) (2 points) Show that  $\psi_1 = \psi_2$ .
  - (ii) (1 point) Show that  $A = H$ .

*Solution.*

- (a). Let  $X, Y$  be sets and  $f : X \rightarrow Y$  be a map.

Suppose that  $f$  is injective. If  $g_1, g_2 : Z \rightarrow X$  are maps such that  $f \circ g_1 = f \circ g_2$ , then, for every  $z \in Z$ , we have  $f(g_1(z)) = f(g_2(z))$ , hence  $g_1(z) = g_2(z)$ ; so  $g_1 = g_2$ . This shows that  $f$  is a monomorphism.

Conversely, suppose that  $f$  is a monomorphism. Let  $x, x' \in X$  such that  $x \neq x'$ . Let  $\{*\}$  be a singleton, and consider the maps  $g_1, g_2 : \{*\} \rightarrow X$  defined by  $g_1(*) = x$  and  $g_2(*) = x'$ . As  $g_1 \neq g_2$ , we have  $f \circ g_1 \neq f \circ g_2$ , so  $f(x) \neq f(x')$ . This shows that  $f$  is injective.

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<sup>1</sup>In fact, the epimorphisms in **Top** are the surjective continuous maps.

Suppose that  $f$  is surjective. If  $h_1, h_2 : Y \rightarrow Z$  are maps such that  $h_1 \circ f = h_2 \circ f$ , then, for every  $y \in Y$ , there exists  $x \in X$  such that  $f(x) = y$ , and then  $h_1(y) = h_1(f(x)) = h_2(f(x)) = h_2(y)$ ; so  $h_1 = h_2$ . This shows that  $f$  is a monomorphism.

Conversely, suppose that  $f$  is an epimorphism. Let  $y_0 \in Y$ , let  $Z = \{a, b\}$  be a set with two distinct elements, and define  $h_1, h_2 : Y \rightarrow Z$  by  $h_1(y) = a$  for every  $y \in Y$ ,  $h_2(y) = a$  for every  $y \in Y - \{y_0\}$  and  $h_2(y_0) = b$ . We have  $h_1 \neq h_2$ , so  $h_1 \circ f \neq h_2 \circ f$ . As  $h_1$  and  $h_2$  coincide on  $Y - \{y_0\}$ , this implies that  $y_0 \in \text{Im}(f)$ . So  $f$  is surjective.

- (b). Let  $f : X \rightarrow Y$  be a morphism of  $\mathcal{C}$ . Suppose that  $F(f)$  is injective. Let  $g_1, g_2 : Z \rightarrow X$  be morphisms of  $\mathcal{C}$  such that  $f \circ g_1 = f \circ g_2$ . Then  $F(f) \circ F(g_1) = F(f) \circ F(g_2)$ , so  $F(g_1) = F(g_2)$  by a). As  $F$  is faithful, this implies that  $g_1 = g_2$ . So  $f$  is a monomorphism.

Suppose that  $F(f)$  is surjective. Let  $h_1, h_2 : Y \rightarrow Z$  be morphisms of  $\mathcal{C}$  such that  $h_1 \circ f = h_2 \circ f$ . Then  $F(h_1) \circ F(f) = F(h_2) \circ F(f)$ , so  $F(h_1) = F(h_2)$  by a). As  $F$  is faithful, this implies that  $h_1 = h_2$ . So  $f$  is an epimorphism.

- (c). By b), any  $R$ -linear that is injective (resp. surjective) is a monomorphism (resp. epimorphism) in  $R\mathbf{Mod}$ .

Conversely, let  $f : M \rightarrow N$  be a monomorphism in  $R\mathbf{Mod}$ . Consider the inclusion map  $g_1 : \text{Ker}(f) \rightarrow M$  and the map  $g_2 = 0 : \text{Ker}(f) \rightarrow M$ . By definition of the kernel, we have  $f \circ g_1 = f \circ g_2 = 0$ , so  $g_1 = g_2$ , so  $\text{Ker}(f) = 0$ , so  $f$  is injective.

Now let  $f : M \rightarrow N$  be an epimorphism in  $R\mathbf{Mod}$ . Consider the obvious surjection  $h_1 : N \rightarrow \text{Coker}(f)$  and the zero map  $h_2 : N \rightarrow \text{Coker}(f)$ . By definition of the cokernel, we have  $h_1 \circ f = h_2 \circ f = 0$ , so  $h_1 = h_2$ , so  $\text{Coker}(f) = 0$ , so  $f$  is surjective.

- (d). By b), we know that any (continuous) injection is a monomorphism in  $\mathbf{Top}$ . Conversely, let  $f : X \rightarrow Y$  be a monomorphism in  $\mathbf{Top}$ . Let  $x, x' \in X$  such that  $x \neq x'$ . Let  $\{*\}$  be a singleton with the discrete topology, and consider the maps  $g_1, g_2 : \{*\} \rightarrow X$  defined by  $g_1(*) = x$  and  $g_2(*) = x'$ ; these maps are continuous, hence morphisms in  $\mathbf{Top}$ . As  $g_1 \neq g_2$ , we have  $f \circ g_1 \neq f \circ g_2$ , so  $f(x) \neq f(x')$ . This shows that  $f$  is injective.

Let  $X = \{s, \eta\}$  be a set with two distinct points. We put the topology on  $X$  for which the open sets are  $\emptyset$ ,  $X$  and  $\{\eta\}$ . Note that  $\{\eta\}$  is dense in  $X$ . Let  $f : X \rightarrow X$  be the map sending every point of  $X$  to  $\eta$ . Then  $f$  has dense image, but  $f$  is not an epimorphism, because  $\text{id}_X \circ f = f \circ f$ , while  $f \neq \text{id}_X$ .

- (e). Let  $\mathcal{C}$  be the subcategory of  $\mathbf{Set}$  whose objects are  $\{0\}$  and  $\{0, 1\}$ , and whose morphisms are the identities and the unique map  $f$  from  $\{0, 1\}$  to  $\{0\}$ . Then  $f$  is a monomorphism in  $\mathcal{C}$ , but it is not injective. (And the inclusion is a faithful functor from  $\mathcal{C}$  to  $\mathbf{Set}$ .)
- (f). Consider the inclusion  $f : \mathbb{Z} \rightarrow \mathbb{Q}$ . It is an epimorphism in  $\mathbf{Ring}$ . Indeed, let  $R$  be a ring and let  $h_1, h_2 : \mathbb{Q} \rightarrow R$  are morphisms of rings such that  $h_1 \circ f = h_2 \circ f$ . For every  $m \in \mathbb{Z} - \{0\}$ , the image of  $m$  in  $\mathbb{Q}$  is invertible, so  $h_1(m), h_2(m) \in R^\times$ . For every  $x \in \mathbb{Q}$ , we can write  $x = nm^{-1}$  with  $n \in \mathbb{N}$  and  $m \in \mathbb{Z} - \{0\}$ , and then  $h_1(x) = h_1(n)h_1(m)^{-1} = h_2(n)h_2(m)^{-1} = h_2(x)$ .

More generally, if  $A$  is a commutative ring and  $S$  is a multiplicative subset of  $A$ , then the canonical map  $A \rightarrow S^{-1}A$  is an epimorphism in  $\mathbf{Ring}$ .

- (g). <sup>2</sup>

(i) Note that  $\psi_1(h)|_{S-\{*,A\}} = \psi_2(h)_{S-\{*,A\}}$  for every  $h \in H$ .

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<sup>2</sup>This proof comes from [?].

Let  $h \in A = \text{Im}(\phi)$ . We have  $\psi_1(h)(*) = *$ . On the other hand, the action of  $h$  on  $H/A$  by left translation fixes  $A$ , so  $\psi_1(h)(A) = A$ . So  $\psi_1(h)_{\{*,A\}}$  is the identity morphism of  $\{*, A\}$ . This implies that  $\psi_2(h)_{\{*,A\}}$  is also the identity morphism of  $\{*, A\}$ , hence that  $\psi_1(h) = \psi_2(h)$ . So  $\psi_1$  and  $\psi_2$  are equal on the image of  $\phi$ , which implies that  $\psi_1 \circ \phi = \psi_2 \circ \phi$ . As  $\phi$  is an epimorphism, we deduce that  $\psi_1 = \psi_2$ .

- (ii) Let  $h \in A$ . Then  $\psi_1(h)(*) = *$ , and  $\psi_2(h)(*) = \sigma \circ \psi_1(h)(A) = \sigma(hA)$ . By (i), we know that  $\psi_1(h) = \psi_2(h)$ , so  $* = \sigma(hA)$ . This is only possible if  $hA = A$ , i.e. if  $h \in A$ . So  $H = A = \text{Im}(\phi)$ , and  $\phi$  is surjective.

□

2. Let  $F : \mathcal{C} \rightarrow \mathcal{C}'$  be a functor.

- (a). (3 points) If  $F$  has a quasi-inverse, show that it is fully faithful and essentially surjective.  
(b). (4 points) If  $F$  is fully faithful and essentially surjective, construct a functor  $G : \mathcal{C}' \rightarrow \mathcal{C}$  and isomorphisms of functors  $F \circ G \simeq \text{id}_{\mathcal{C}}$  and  $G \circ F \simeq \text{id}_{\mathcal{C}'}$ .

*Solution.*

- (a). Let  $G : \mathcal{C}' \rightarrow \mathcal{C}$  be a quasi-inverse of  $F$ , and let  $u : G \circ F \xrightarrow{\sim} \text{id}_{\mathcal{C}}$  and  $v : F \circ G \xrightarrow{\sim} \text{id}_{\mathcal{C}'}$  be isomorphisms of functors.

Let  $X, Y \in \text{Ob}(\mathcal{C})$ . We denote by  $\beta$  the map  $\text{Hom}_{\mathcal{C}}(X, Y) \rightarrow \text{Hom}_{\mathcal{C}'}(F(X), F(Y))$  given by  $F$ . Consider the map  $\alpha : \text{Hom}_{\mathcal{C}'}(F(X), F(Y)) \rightarrow \text{Hom}_{\mathcal{C}}(X, Y)$  that we get by composing  $G : \text{Hom}_{\mathcal{C}'}(F(X), F(Y)) \rightarrow \text{Hom}_{\mathcal{C}}(G \circ F(X), G \circ F(Y))$  and the map  $\text{Hom}_{\mathcal{C}}(G \circ F(X), G \circ F(Y)) \rightarrow \text{Hom}_{\mathcal{C}}(X, Y)$ ,  $g \mapsto u(Y) \circ g \circ u(X)^{-1}$ . We claim that  $\alpha \circ \beta$  is the identity on  $\text{Hom}_{\mathcal{C}}(X, Y)$ . Indeed, let  $f \in \text{Hom}_{\mathcal{C}}(X, Y)$ . As  $u$  is a morphism of functors, the following diagram is commutative :

$$\begin{array}{ccc} G \circ F(X) & \xrightarrow{G \circ F(f)} & G \circ F(Y) \\ u(X) \downarrow & & \downarrow u(Y) \\ X & \xrightarrow{f} & Y \end{array}$$

This shows that  $u(Y) \circ G \circ F(f) \circ u(X)^{-1} = f$ , i.e. that  $\alpha \circ \beta(f) = f$ . In particular, the map  $\beta$  is injective and the map  $\alpha$  is surjective. This shows that  $F$  is faithful. Applying this result to  $G$  (which is also an equivalence of categories, with quasi-inverse  $F$ ), we see that the map  $\alpha$  is also injective, hence it is bijective, hence  $\beta$  is also bijective. This shows that  $F$  is fully faithful.

Let  $X' \in \text{Ob}(\mathcal{C}')$ . Then  $v : F(G(X')) \xrightarrow{\sim} X'$  is an isomorphism, and  $G(X') \in \text{Ob}(\mathcal{C})$ . This shows that  $F$  is essentially surjective.

- (b). We construct the functor  $G$ . Let  $X' \in \text{Ob}(\mathcal{C}')$ ; we choose an object  $X$  of  $\mathcal{C}$  and an isomorphism  $u(X') : F(X) \xrightarrow{\sim} X'$ , and we set  $G(X') = X$ . Let  $X', Y' \in \text{Ob}(\mathcal{C}')$ , and let  $X = G(X')$  and  $Y = G(Y')$ . We define a map  $\text{Hom}_{\mathcal{C}'}(X', Y') \rightarrow \text{Hom}_{\mathcal{C}'}(F(X), F(Y))$  by  $f' \mapsto u(Y')^{-1} \circ f' \circ u(X')$ . Composing this with the inverse of the bijection  $F : \text{Hom}_{\mathcal{C}}(X, Y) \xrightarrow{\sim} \text{Hom}_{\mathcal{C}'}(F(X), F(Y))$ , we get a map  $\text{Hom}_{\mathcal{C}'}(X', Y') \rightarrow \text{Hom}_{\mathcal{C}}(X, Y)$ , which we denote by  $G$ .

Next we show that  $G$  is a functor. If  $X' \in \text{Ob}(\mathcal{C}')$ , then  $u(X')^{-1} \circ \text{id}_{X'} \circ u(X') = \text{id}_{F(G(X'))}$ , so  $G(\text{id}_{X'}) = \text{id}_{G(X')}$ . Let  $f' : X' \rightarrow Y'$  and  $g' : Y' \rightarrow Z'$  be two morphisms of  $\mathcal{C}'$ ,

and let  $f : X \rightarrow Y$  and  $g : Y \rightarrow Z$  be their images by  $G$ . By definition of  $G$  on morphisms, we have  $F(f) = u(Y')^{-1} \circ f' \circ u(X')$  and  $F(g) = u(Z')^{-1} \circ g' \circ u(Y')$ , so  $F(g \circ f) = F(g) \circ F(f) = u(Z')^{-1} \circ (g' \circ f') \circ u(X') = F(G(g' \circ f'))$ . As  $F$  is faithful, this implies that  $g \circ f = G(g' \circ f')$ , i.e. that  $G(g') \circ G(f') = G(g' \circ f')$ . So  $G$  is a functor.

Finally, we show that  $G$  is a quasi-inverse of  $F$ . For every  $X' \in \text{Ob}(\mathcal{C}')$ , we have by definition of  $G(X')$  an isomorphism  $u(X') : F(G(X')) \xrightarrow{\sim} X'$ . We need to show that this defines an isomorphism of functors  $F \circ G \xrightarrow{\sim} \text{id}_{\mathcal{C}'}$ . So let  $f' : X' \rightarrow Y'$  be a morphism of  $\mathcal{C}'$ . By definition of  $G(f')$ , we have  $u(Y') \circ F(G(f')) = f' \circ u(X')$ , which is what we wanted. We still need to define an isomorphism of functors  $v : G \circ F \xrightarrow{\sim} \text{id}_{\mathcal{C}}$ . Let  $X \in \text{Ob}(\mathcal{C})$ . By definition of  $G$ , we have an isomorphism  $u(F(X)) : F(G(F(X))) \xrightarrow{\sim} F(X)$ . As  $F$  is fully faithful, there is a unique  $v(X) \in \text{Hom}_{\mathcal{C}}(G(F(X)), X)$  such that  $F(v(X)) = u(F(X))$ , and  $v(X)$  is an isomorphism because a fully faithful functor is conservative. It remains to show that this defines a morphism of functors. So let  $f : X \rightarrow Y$  be a morphism of  $\mathcal{C}$ . Then  $F(G(F(f))) = u(F(Y))^{-1} \circ F(f) \circ u(F(X))$ , so

$$F(f) \circ F(v(X)) = F(f) \circ u(F(X)) = u(F(Y)) \circ F(G(F(f))) = F(v(Y)) \circ F(G(F(f))).$$

Using the fact that  $F$  is faithful (and is a functor), we get  $f \circ v(X) = v(Y) \circ G(F(f))$ , which is what we wanted. □

**3.** Let  $\mathcal{C}$  be the full subcategory of  $\mathbf{Ab}$  whose objects are finitely generated abelian groups.

- (a). (2 points) Show that every natural endomorphism of  $\text{id}_{\mathcal{C}}$  is multiplication by some  $n \in \mathbb{Z}$ .
- (b). (3 points) Consider the functor  $F : \mathcal{C} \rightarrow \mathcal{C}$  that sends an abelian group  $A$  to  $A_{\text{tor}} \oplus (A/A_{\text{tor}})$  (and acts in the obvious way on morphisms), where  $A_{\text{tor}}$  is the torsion subgroup of  $A$ . Show that there is no natural isomorphism  $F \xrightarrow{\sim} \text{id}_{\mathcal{C}}$ .

*Solution.*

- (a). Let  $u : \text{id}_{\mathcal{C}} \rightarrow \text{id}_{\mathcal{C}}$  be a morphism of functors. Then  $u(\mathbb{Z}) \in \text{End}_{\mathbf{Ab}}(\mathbb{Z})$ , so  $u(\mathbb{Z})$  is of the form  $n\text{id}_{\mathbb{Z}}$  for some  $n \in \mathbb{Z}$ . Let  $A$  be an arbitrary abelian group. We want to show that  $u(A) = n\text{id}_A$ . Let  $a \in A$ . We consider the morphism of groups  $f : \mathbb{Z} \rightarrow A$  sending 1 to  $a$ . As  $u$  is a morphism of functors, we have a commutative diagram :

$$\begin{array}{ccc} \mathbb{Z} & \xrightarrow{u(\mathbb{Z})} & \mathbb{Z} \\ f \downarrow & & \downarrow f \\ A & \xrightarrow{u(A)} & A \end{array}$$

In particular,  $u(A)(a) = u(A)(f(1)) = f(u(\mathbb{Z})(1)) = f(n) = na$ . So  $u(A) = n\text{id}_A$ .

- (b). Suppose that  $u : F \xrightarrow{\sim} \text{id}_{\mathcal{C}}$  is a natural isomorphism. For every abelian groups  $A$ , consider the morphism  $v(A) : A \rightarrow A/A_{\text{tor}} \oplus A_{\text{tor}}$  that is the composition of the canonical surjection  $A \rightarrow A/A_{\text{tor}}$  and of the injection  $A/A_{\text{tor}} \rightarrow A/A_{\text{tor}} \oplus A_{\text{tor}}$ . It is easy to see that this defines a morphism of functors  $v : \text{id}_{\mathcal{C}} \rightarrow F$ . So  $u \circ v$  is an endomorphism of  $\text{id}_{\mathcal{C}}$ , and, by a), there exists  $n \in \mathbb{Z}$  such that  $u \circ v$  is the multiplication by  $n$ . As  $v(\mathbb{Z}) = \text{id}_{\mathbb{Z}}$  by definition of  $v$  and  $u(\mathbb{Z})$  is an isomorphism, we must have  $n = \pm 1$ . Now take  $A = \mathbb{Z}/2\mathbb{Z}$ . Then  $v(A) = 0$ , so  $u \circ v(A) = 0$ , so  $n$  is divisible by 2. This is a contradiction. □

4. (2 points, extra credit) Let  $k$  be a field, and let  $F : \mathbf{Mod}_k \rightarrow \mathbf{Mod}_k$  be the functor sending a  $k$ -vector space  $V$  to  $V \otimes_k V$  and a  $k$ -linear transformation  $f$  to  $f \otimes f$ . Show that the only morphism of functors from  $\text{id}_{\mathbf{Mod}_k}$  to  $F$  is the zero one, i.e. the morphism  $u : \text{id}_{\mathbf{Mod}_k} \rightarrow F$  such that  $u(V) = 0$  for every  $k$ -vector space  $V$ .

*Solution.* Let  $u : \text{id}_{\mathbf{Mod}_k} \rightarrow F$  be a morphism of functors. Then  $u(k)$  is a  $k$ -linear map from  $k$  to  $k \otimes_k k$ , so there exists a unique  $\lambda \in k$  such that  $u(k)(1) = \lambda(1 \otimes 1)$ .

Let  $V$  be a  $k$ -vector space, and let  $v \in V$ . We denote by  $f; k \rightarrow V$  the unique  $k$ -linear map such that  $f(1) = v$ . As  $u$  is a morphism of functors, we have  $u(V) \circ f = (f \otimes f) \circ u(k)$ , and in particular  $u(V)(v) = u(V)(f(1)) = (f \otimes f)(\lambda(1 \otimes 1)) = \lambda(v \otimes v)$ .

Take  $V = k^2$ , and let  $(e_1, e_2)$  be the canonical basis of  $V$ . We know that  $(e_1 \otimes e_1, e_1 \otimes e_2, e_2 \otimes e_1, e_2 \otimes e_2)$  is a basis of  $V \otimes_k V$ . Using the previous paragraph, we see that

$$u(V)(e_1 + e_2) = \lambda(e_1 + e_2) \otimes (e_1 + e_2) = \lambda(e_1 \otimes e_1) + \lambda(e_1 \otimes e_2) + \lambda(e_2 \otimes e_1) + \lambda(e_2 \otimes e_2).$$

On the other hand, as  $u(V)$  is  $k$ -linear, we have

$$u(V)(e_1 + e_2) = u(V)(e_1) + u(V)(e_2) = \lambda(e_1 \otimes e_1) + \lambda(e_2 \otimes e_2).$$

This is only possible if  $\lambda = 0$ . But then, by the calculation of the previous paragraph, we have  $u(W) = 0$  for every  $k$ -vector space  $W$ .

Note that we did not use the fact that  $k$  is a field, so the result is also true for the category of modules over a commutative ring.

□

5. (4 points) Let  $\mathcal{C}$  be a category. Remember that the category  $\text{PSh}(\mathcal{C})$  of presheaves on  $\mathcal{C}$  is the category  $\text{Func}(\mathcal{C}^{\text{op}}, \mathbf{Set})$ .

Let  $F$  be a presheaf on  $\mathcal{C}$  and  $X$  be an object of  $\mathcal{C}$ . Let  $\Phi : \text{Hom}_{\text{PSh}(\mathcal{C})}(h_X, F) \rightarrow F(X)$  be the map defined by  $\Phi(u) = u(X)(\text{id}_X)$ . Let  $\Psi : F(X) \rightarrow \text{Hom}_{\text{PSh}(\mathcal{C})}(h_X, F)$  be the map sending  $x \in F(X)$  to the morphism of functors  $\Psi(x) : h_X \rightarrow F$  such that  $\Psi(x)(Y) : h_X(Y) = \text{Hom}_{\mathcal{C}}(Y, X) \rightarrow F(Y)$  sends  $f : Y \rightarrow X$  to  $F(f)(x) \in F(Y)$ . Show that  $\Phi$  and  $\Psi$  are bijections that are inverses of each other.

*Solution.* We show that  $\Psi \circ \Phi$  is the identity of  $\text{Hom}_{\text{PSh}(\mathcal{C})}(h_X, F)$ . Let  $u \in \text{Hom}_{\text{PSh}(\mathcal{C})}(h_X, F)$ . Let  $Y$  be an object of  $\mathcal{C}$ . As  $u$  is a morphism of functors, we have a commutative diagram

$$\begin{array}{ccc} \text{Hom}_{\mathcal{C}}(X, X) & \xrightarrow{u(X)} & F(X) \\ h_X(f) \downarrow & & \downarrow F(f) \\ \text{Hom}_{\mathcal{C}}(Y, X) & \xrightarrow{u(Y)} & F(Y) \end{array}$$

In particular, we have

$$F(f)(\Phi(u)) = F(f)(u(X)(\text{id}_X)) = u(Y)(h_X(f)(\text{id}_X)) = u(Y)(f).$$

As  $F(f)(\Phi(u)) = \Psi(\Phi(u))(Y)(f)$  by definition of  $\Psi$ , this shows that  $\Psi(\Phi(u))(Y) = u(Y)$ , hence that  $\Psi(\Phi(u)) = u$ .

Now we show that  $\Phi \circ \Psi$  is the identity of  $F(X)$ . Let  $x \in F(X)$ . Then  $\Phi(\Psi(x)) = \Psi(x)(X)(\text{id}_X) = F(\text{id}_X)(x) = \text{id}_{F(X)}(x) = x$ .

□

6.

- (a). (2 points) Show that the categories **Set** and **Set**<sup>op</sup> are not equivalent. (Hint : If  $F : \mathbf{Set}^{\text{op}} \rightarrow \mathbf{Set}$  is an equivalence of categories, show that  $F(\emptyset)$  is a singleton and that  $F(X) = \emptyset$  for  $X$  a singleton.)
- (b). (1 point) Let  $\mathcal{C}$  be the full subcategory of **Set** whose objects are finite sets. Show that  $\mathcal{C}$  and  $\mathcal{C}^{\text{op}}$  are not equivalent.
- (c). (1 point) Show that **Rel** and **Rel**<sup>op</sup> are equivalent.
- (d). (2 points) Let  $\mathcal{D}$  be the full subcategory of **Ab** whose objects are finite abelian groups. Show that  $\mathcal{D}$  and  $\mathcal{D}^{\text{op}}$  are equivalent.

*Solution.*

- (a). Suppose that there exists an equivalence of categories  $F : \mathbf{Set} \rightarrow \mathbf{Set}^{\text{op}}$ . For every set  $X$ , the set

$$\text{Hom}_{\mathbf{Set}^{\text{op}}}(F(X), F(\emptyset)) = \text{Hom}_{\mathbf{Set}}(F(\emptyset), F(X)) \simeq \text{Hom}_{\mathbf{Set}}(\emptyset, X)$$

is a singleton (because there is a unique map from the empty set into  $X$ ). So  $F(\emptyset)$  is a singleton.

Similarly, if  $X$  is a singleton, then, for every set  $Y$ , the set

$$\text{Hom}_{\mathbf{Set}^{\text{op}}}(F(X), F(Y)) = \text{Hom}_{\mathbf{Set}}(F(Y), F(X)) \simeq \text{Hom}_{\mathbf{Set}}(Y, X)$$

is a singleton. So  $F(X)$  is the empty set.

Now let  $X$  be a singleton and  $Y$  be a set with two elements. Then  $\text{Hom}_{\mathbf{Set}}(X, Y)$  is a set with two elements. But on the other hand, we have

$$\text{Hom}_{\mathbf{Set}}(X, Y) \simeq \text{Hom}_{\mathbf{Set}^{\text{op}}}(F(X), F(Y)) = \text{Hom}_{\mathbf{Set}}(F(Y), \emptyset),$$

and  $\text{Hom}_{\mathbf{Set}}(F(Y), \emptyset)$  has at most one element (it is empty if  $F(Y) \neq \emptyset$ , and it only contains  $\text{id}_{\emptyset}$  if  $F(Y) = \emptyset$ ). This is a contradiction.

- (b). The proof of a) works just as well.
- (c). Let  $F : \mathbf{Rel} \rightarrow \mathbf{Rel}^{\text{op}}$  be defined by  $F(X) = X$  for every set  $X$  and, for all sets  $X, Y$  and every subset  $f$  of  $X \times Y$ ,  $F(f) = \{(y, x) \mid (x, y) \in f\}$ . We want to show that  $F$  is a functor. (Then it will clearly be an equivalence, and even an isomorphism of categories.) Let  $X, Y, Z$  be sets and  $f : X \rightarrow Y$ ,  $g : Y \rightarrow Z$  be morphisms in **Rel**; that is,  $f$  is a subset of  $X \times Y$  and  $g$  is a subset of  $Y \times Z$ . Then, in **Rel**, we have  $g \circ f = \{(x, z) \mid \exists y \in Y, (x, y) \in f \text{ and } (y, z) \in g\}$ . On the other hand, in **Rel**<sup>op</sup>, we have  $F(f) \circ F(g) = \{(z, x) \in Z \times X \mid \exists y \in Y, (y, x) \in F(f) \text{ and } (z, y) \in F(g)\}$ . This is clearly equal to  $F(g \circ f)$ .
- (d). Consider the functor  $F = \text{Hom}_{\mathbf{Ab}}(\cdot, \mathbb{Q}/\mathbb{Z}) : \mathbf{Ab}^{\text{op}} \rightarrow \mathbf{Ab}$ . If  $A$  is a finite abelian group, then so is  $F(A)$ . So  $F$  induces a functor  $\mathcal{D}^{\text{op}} \rightarrow \mathcal{D}$ , which we still denote by  $F$ . We can also see  $F$  as a functor from  $\mathcal{D}$  to  $\mathcal{D}^{\text{op}}$ . We claim that  $F$  is an equivalence of categories, and in fact that it is its own quasi-inverse. To show this, it suffices to construct an

functorial isomorphism  $\text{id}_{\mathcal{C}} \xrightarrow{\sim} F \circ F$ . For every finite abelian group  $A$ , we consider the map  $u : A \rightarrow F(F(A)) = \text{Hom}_{\mathbf{Ab}}(\text{Hom}_{\mathbf{Ab}}(A, \mathbb{Q}/\mathbb{Z}), \mathbb{Q}/\mathbb{Z})$ ,  $a \mapsto (f \mapsto f(a))$ . The fact that this defines a morphism of functors is a straightforward verification. The fact that is an isomorphism if Pontrjagin duality for finite abelian groups. (By the structure theorem for finite abelian groups, it suffices to check that  $u(A)$  is an isomorphism for  $A$  of the form  $\mathbb{Z}/n\mathbb{Z}$ , which is easy.)

□

**7.** (4 points) Let  $\mathcal{C}$  and  $\mathcal{C}'$  and  $F : \mathcal{C} \rightarrow \mathcal{C}'$ ,  $G : \mathcal{C}' \rightarrow \mathcal{C}$  be two functors. We consider the two bifunctors  $H_1, H_2 : \mathcal{C}^{\text{op}} \times \mathcal{C} \rightarrow \mathbf{Set}$  defined by  $H_1 = \text{Hom}_{\mathcal{C}'}(F(\cdot), \cdot)$  and  $H_2 = \text{Hom}_{\mathcal{C}}(\cdot, G(\cdot))$ . Suppose that we are given, for every  $X \in \text{Ob}(\mathcal{C})$  and every  $Y \in \text{Ob}(\mathcal{C}')$ , a bijection  $\alpha(X, Y) : H_1(X, Y) \xrightarrow{\sim} H_2(X, Y)$ . Show that the two following statements are equivalent :

- (i) The family of bijections  $(\alpha(X, Y))_{X \in \text{Ob}(\mathcal{C}), Y \in \text{Ob}(\mathcal{C}'})$  defines an isomorphism of functors  $H_1 \xrightarrow{\sim} H_2$ .
- (ii) For every morphism  $f : X_1 \rightarrow X_2$  in  $\mathcal{C}$ , every morphism  $g : Y_1 \rightarrow Y_2$  in  $\mathcal{C}'$ , and for all  $u \in \text{Hom}_{\mathcal{C}'}(F(X_1), Y_1)$  and  $v \in \text{Hom}_{\mathcal{C}'}(F(X_2), Y_2)$ , the square

$$\begin{array}{ccc} F(X_1) & \xrightarrow{u} & Y_1 \\ F(f) \downarrow & & \downarrow g \\ F(X_2) & \xrightarrow{v} & Y_2 \end{array}$$

is commutative if and only if the square

$$\begin{array}{ccc} X_1 & \xrightarrow{\alpha(X_1, Y_1)(u)} & G(Y_1) \\ f \downarrow & & \downarrow G(g) \\ X_2 & \xrightarrow{\alpha(X_2, Y_2)(v)} & G(Y_2) \end{array}$$

is commutative.

*Solution.* The key is to write explicitly what it means for the  $(\alpha(X, Y))$  to define a morphism of functors. It means that, for every morphism  $f : X_1 \rightarrow X_2$  in  $\mathcal{C}$  (that is, a morphism  $X_2 \rightarrow X_1$  in  $\mathcal{C}^{\text{op}}$ ) and for every morphism  $g : Y_1 \rightarrow Y_2$  in  $\mathcal{C}'$ , the following square commutes :

$$\begin{array}{ccc} \text{Hom}_{\mathcal{C}'}(F(X_2), Y_1) & \xrightarrow{H_1(f, g)} & \text{Hom}_{\mathcal{C}'}(F(X_1), Y_2) \\ \alpha(X_2, Y_1) \downarrow & & \downarrow \alpha(X_1, Y_2) \\ \text{Hom}_{\mathcal{C}}(X_2, G(Y_1)) & \xrightarrow{H_2(f, g)} & \text{Hom}_{\mathcal{C}}(X_1, G(Y_2)) \end{array}$$

The fact that the square commutes says exactly that, for every morphism  $w : F(X_2) \rightarrow Y_1$  in  $\mathcal{C}'$ , we have

$$\alpha(X_1, Y_2)(g \circ w \circ F(f)) = G(g) \circ \alpha(X_2, Y_1)(w) \circ f.$$

Suppose that (i) holds. Using the calculation of the previous, we get :

- (a) Taking  $X_1 = X_2$ ,  $f = \text{id}_{X_1}$  and  $g : Y_1 \rightarrow Y_2$  arbitrary : for every  $u : F(X_1) \rightarrow Y_1$ , we have

$$\alpha(X_1, Y_2)(g \circ u) = G(g) \circ \alpha(X_1, Y_1)(u).$$

(b) Taking  $f : X_1 \rightarrow X_2$  arbitrary,  $Y_1 = Y_2$  and  $g = \text{id}_{Y_2} : Y_1 \rightarrow Y_2$ , for every  $v : F(X_2) \rightarrow Y_2$ , we have

$$\alpha(X_1, Y_2)(v \circ F(f)) = \alpha(X_2, Y_2)(v) \circ f.$$

Suppose that we are in the situation of (ii), that is, we are given morphisms  $f : X_1 \rightarrow X_2$  in  $\mathcal{C}$ ,  $g : Y_1 \rightarrow Y_2$  in  $\mathcal{C}'$ , and  $u \in \text{Hom}_{\mathcal{C}'}(F(X_1), Y_1)$  and  $v \in \text{Hom}_{\mathcal{C}'}(F(X_2), Y_2)$ . We want to show that the top square of (ii) commutes if and only if the bottom square commutes.

Suppose that the top square commutes, that is, that  $v \circ F(f) = g \circ u$ . Applying (a) and (b), we get

$$G(g) \circ \alpha(X_1, Y_1)(u) = \alpha(X_1, Y_2)(g \circ u) = \alpha(X_1, Y_2)(v \circ F(f)) = \alpha(X_2, Y_2)(v) \circ f.$$

This shows that the bottom square commutes.

Conversely, suppose that the bottom square commutes, that is, that  $G(g) \circ \alpha(X_1, Y_1)(u) = \alpha(X_2, Y_2)(v) \circ f$ . Again, applying (a) and (b), we get

$$\alpha(X_1, Y_2)(g \circ u) = G(g) \circ \alpha(X_1, Y_1)(u) = \alpha(X_2, Y_2)(v) \circ f = \alpha(X_1, Y_2)(v \circ F(f)).$$

As  $\alpha(X_1, Y_2)$  is bijective, this implies that  $g \circ u = v \circ F(f)$ , which means that the top square commutes.

Now we assume that (ii) holds, and we want to show that (i) also holds. Let  $f : X_1 \rightarrow X_2$  be a morphism in  $\mathcal{C}$ ,  $g : Y_1 \rightarrow Y_2$  be a morphism in  $\mathcal{C}'$ , and  $w : F(X_2) \rightarrow Y_1$  be a morphism in  $\mathcal{C}'$ . We want to show that  $\alpha(X_1, Y_2)(g \circ w \circ F(f)) = G(g) \circ \alpha(X_2, Y_1)(w) \circ f$ . We apply (i) to  $u = w \circ F(f) : F(X_1) \rightarrow Y_1$  and  $v = g \circ w : F(X_2) \rightarrow Y_2$ . We obviously have  $g \circ u = v \circ F(f)$ , so, by (i), this implies that

$$(*) \quad \alpha(X_2, Y_2)(g \circ w) \circ f = G(g) \circ \alpha(X_1, Y_1)(w \circ F(f)).$$

Applying (\*) to the particular case where  $Y_1 = Y_2$  and  $g = \text{id}_{Y_1}$ , we get:

$$(**) \quad \alpha(X_2, Y_1)(w) \circ f = \alpha(X_1, Y_1)(w \circ F(f)).$$

Applying (\*\*) with  $w$  replaced by  $g \circ w : F(X_2) \rightarrow Y_2$ , we get

$$(***) \quad \alpha(X_2, Y_2)(g \circ w) \circ f = \alpha(X_1, Y_2)(g \circ w \circ F(f)).$$

Putting (\*), (\*\*) and (\*\*\*) together gives

$$\begin{aligned} \alpha(X_1, Y_2)(g \circ w \circ F(f)) &= \alpha(X_2, Y_2)(g \circ w) \circ f = G(g) \circ \alpha(X_1, Y_1)(w \circ F(f)) \\ &= G(g) \circ \alpha(X_1, Y_1)(w \circ F(f)), \end{aligned}$$

which is what we wanted to prove. □

**8.** Remember that a functor  $F : \mathcal{C} \rightarrow \mathbf{Set}$  is called representable if there exists an object  $X$  of  $\mathcal{C}$  and an element  $x$  of  $F(X)$  such that the morphism of functors  $u : \text{Hom}_{\mathcal{C}}(X, \cdot) \rightarrow F$  defined by  $u(Y) : \text{Hom}_{\mathcal{C}}(X, Y) \rightarrow F(Y)$ ,  $(f : X \rightarrow Y) \mapsto F(f)(x)$  is an isomorphism. The couple  $(X, x)$  is then said to represent  $F$ .

The following functors are representable. For each of them, give a couple representing the functor. (If the functor is only defined on objects, it is assumed to act on morphisms in the obvious way.) (1 point per functor)



- (a). The identity endofunctor of **Set**.
- (b). The functor  $F : \mathbf{Grp} \rightarrow \mathbf{Set}$ ,  $G \mapsto G^n$ , where  $n \in \mathbb{N}$ .
- (c). The forgetful functor  $\mathbf{Mod}_R \rightarrow \mathbf{Set}$ , where  $R$  is a ring.
- (d). The forgetful functor  $\mathbf{Ring} \rightarrow \mathbf{Set}$ .
- (e). The functor  $\mathbf{Ring} \rightarrow \mathbf{Set}$ ,  $R \mapsto R^\times$ .
- (f). The functor  $F : \mathbf{Cat} \rightarrow \mathbf{Set}$  that takes a category to its set of objects.
- (g). The functor  $F : \mathbf{Cat} \rightarrow \mathbf{Set}$  that takes a category to its set of morphisms (i.e.  $\bigcup_{X, Y \in \text{Ob}(\mathcal{C})} \text{Hom}_{\mathcal{C}}(X, Y)$ ).
- (h). The functor  $F : \mathbf{Cat} \rightarrow \mathbf{Set}$  that takes a category to its set of isomorphisms.
- (i). The functor  $F : \mathbf{Top}_* \rightarrow \mathbf{Set}$  that takes a pointed topological space  $(X, x)$  to the set of continuous loops on  $X$  with base point  $x$ .
- (j). The functor  $F : \mathbf{Set}^{\text{op}} \rightarrow \mathbf{Set}$  such that  $F(X) = \mathfrak{P}(X)$  and, for every map  $f : X \rightarrow Y$ ,  $F(f) : \mathfrak{P}(Y) \rightarrow \mathfrak{P}(X)$  is the map  $A \mapsto f^{-1}(A)$ .
- (k). The functor  $F : \mathbf{Top}^{\text{op}} \rightarrow \mathbf{Set}$  that sends a topological space to its set of open subsets. (If  $f : X \rightarrow Y$  is a continuous map,  $F(f) : F(Y) \rightarrow F(X)$  is the map  $U \mapsto f^{-1}(U)$ .)
- (l). If  $k$  is a field, the functor  $F : \mathbf{Mod}_k^{\text{op}} \rightarrow \mathbf{Set}$  that sends a  $k$ -vector space to the underlying set of  $V^*$  (so  $F$  is the composition of the duality functor  $\mathbf{Mod}_k^{\text{op}} \rightarrow \mathbf{Mod}_k$  and of the forgetful functor from  $\mathbf{Mod}_k$  to  $\mathbf{Set}$ .)

*Solution.*

- (a). Take  $X = \{x\}$  to be a singleton and  $x$  to be the unique element of  $F(X) = X$ . Then, for every set  $Y$ ,  $u(Y) : \text{Hom}_{\mathbf{Set}}(X, Y) \rightarrow F(Y) = Y$  sends  $f : X \rightarrow Y$  to  $f(x) \in F(Y) = Y$ ; it is clearly bijective.
- (b). Let  $X = F_n$  be the free group on  $n$  generators  $(x_1, \dots, x_n)$ , and  $x = (x_1, \dots, x_n) \in F(F_n) = (F_n)^n$ . For every group  $G$ , the map  $u(G) : \text{Hom}_{\mathbf{Grp}}(F_n, G) \rightarrow G^n$  sends  $f : F_n \rightarrow G$  to  $(f(x_1), \dots, f(x_n)) \in G^n$ . The fact that this is bijective is the universal property of the free group  $F_n$ .
- (c). Take  $X = R$  with the obvious right  $R$ -action, and  $x = 1 \in F(R) = R$ . Then, for every right  $R$ -module  $M$ , the map  $u(M) : \text{Hom}_R(R, M) \rightarrow F(M) = M$  sends  $f : R \rightarrow M$  to  $f(1)$ . This is bijective because  $R$  is a free  $R$ -module with base  $\{1\}$ .
- (d). Take  $X$  equal to the polynomial ring  $\mathbb{Z}[x]$  and  $x \in F(X) = X$  to be the indeterminate. For every ring  $R$ , the map  $u(R) : \text{Hom}_{\mathbf{Ring}}(\mathbb{Z}[x], R) \rightarrow F(R) = R$  sends  $f : \mathbb{Z}[x] \rightarrow R$  to  $f(x) \in R$ . The fact that this is bijective is the universal property of the polynomial ring.
- (e). Take  $X = \mathbb{Z}[x, x^{-1}]$  (the polynomial ring  $\mathbb{Z}[x]$  localized at the indeterminate  $x$ ) and  $x$  to be the indeterminate in  $F(X) = \mathbb{Z}[x, x^{-1}]^\times$ . For every ring  $R$ , the map  $u(R) : \text{Hom}_{\mathbf{Ring}}(\mathbb{Z}[x], R) \rightarrow F(R) = R^\times$  sends  $f : \mathbb{Z}[x] \rightarrow R$  to  $f(x) \in R^\times$ . The fact that this is bijective follows from the universal properties of the polynomial ring of the localization.
- (f). Let  $X$  be the category with only one object  $*$  and such that  $\text{End}_X(*) = \{\text{id}_*\}$ , and let  $x \in F(X) = \{*\}$  be the unique object. (Note that  $X$  is the category corresponding to the poset  $[0]$ .) If  $\mathcal{C}$  is a category, the map  $u(\mathcal{C}) : \text{Func}(X, \mathcal{C}) \rightarrow F(\mathcal{C}) = \text{Ob}(\mathcal{C})$  takes a functor  $G : X \rightarrow \mathcal{C}$  to  $G(*) \in \text{Ob}(\mathcal{C})$ . This map is bijective, with inverse the map

$v(\mathcal{C}) : \text{Ob}(\mathcal{C}) \rightarrow \text{Func}(X, \mathcal{C})$  sending  $c \in \text{Ob}(\mathcal{C})$  to the functor  $G : X \rightarrow \mathcal{C}$  defined by  $G(*) = c$  and  $G(\text{id}_*) = \text{id}_c$ .

- (g). Let  $X$  be the category corresponding to the poset  $[1]$ , that is,  $X$  has two objects  $0$  and  $1$ , and a unique non-identity morphism  $\alpha : 0 \rightarrow 1$ . Let  $x \in F(X)$  be the morphism  $\alpha$ . If  $\mathcal{C}$  is a category, the map  $u(\mathcal{C}) : \text{Func}(X, \mathcal{C}) \rightarrow F(X)$  sends a functor  $G : X \rightarrow \mathcal{C}$  to  $G(\alpha) \in \text{Hom}_{\mathcal{C}}(F(0), F(1))$ . Let  $v(\mathcal{C}) : F(X) \rightarrow \text{Func}(X, \mathcal{C})$  be defined as follows : if  $f : c_0 \rightarrow c_1$  is a morphism of  $\mathcal{C}$ , that is, an element of  $F(\mathcal{C})$ , we defined a functor  $G : X \rightarrow \mathcal{C}$  by  $G(0) = c_0$ ,  $G(1) = c_1$  and  $G(\alpha) = f$ . Then  $v(\mathcal{C})$  is an inverse of  $u(\mathcal{C})$ , so  $u(\mathcal{C})$  is bijective.

Let  $X$  be the category such that  $\text{Ob}(X) = \{0, 1\}$ , and such that the only two non-identity morphisms of  $X$  are morphisms  $\alpha : 0 \rightarrow 1$  and  $\beta : 1 \rightarrow 0$  such that  $\alpha \circ \beta = \text{id}_1$  and  $\beta \circ \alpha = \text{id}_0$ . If  $\mathcal{C}$  is a category, the map  $u(\mathcal{C}) : \text{Func}(X, \mathcal{C}) \rightarrow F(X)$  sends a functor  $G : X \rightarrow \mathcal{C}$  to  $G(\alpha) \in \text{Hom}_{\mathcal{C}}(F(0), F(1))$ , which is an isomorphism with inverse  $G(\beta)$ . Let  $v(\mathcal{C}) : F(X) \rightarrow \text{Func}(X, \mathcal{C})$  be defined as follows : if  $f : c_0 \rightarrow c_1$  is an isomorphism of  $\mathcal{C}$ , that is, an element of  $F(\mathcal{C})$ , we defined a functor  $G : X \rightarrow \mathcal{C}$  by  $G(0) = c_0$ ,  $G(1) = c_1$ ,  $G(\alpha) = f$  and  $G(\beta) = f^{-1}$ . Then  $v(\mathcal{C})$  is an inverse of  $u(\mathcal{C})$ , so  $u(\mathcal{C})$  is bijective.

- (h). Remember that a loop on a topological space  $Y$  with base point  $y$  is just a continuous map  $\gamma$  from  $S^1$  (the unit circle in  $\mathbb{C}$ ) to  $Y$  such that  $\gamma(1) = y$ . In other words, it is a morphism from  $(S^1, 1)$  to  $(Y, y)$  in the category  $\mathbf{Top}_*$ . So we can take  $X = (S^1, 1)$  and  $x = \text{id}_{S^1} \in F(X)$ .
- (i). For every set  $Y$ , we have a bijection  $v(Y) : \mathfrak{P}(Y) \xrightarrow{\sim} \text{Hom}_{\mathbf{Set}}(Y, \{0, 1\})$  sending a subset  $A$  of  $Y$  to its characteristic function. So we can take  $X = \{0, 1\}$  and  $x = \{1\} \in \mathfrak{P}(X)$ . Indeed, if  $Y$  is a set, then the map  $u(Y) : \text{Hom}_{\mathbf{Set}^{\text{op}}}(X, Y) = \text{Hom}_{\mathbf{Set}}(Y, X) \rightarrow \mathfrak{P}(Y)$  sends  $f : Y \rightarrow \{0, 1\}$  to  $f^{-1}(\{1\})$ , which is the inverse of the bijection  $v(Y)$ .
- (j). Let  $X$  be the Sierpinski space, that is, the topological space  $\{s, \eta\}$  where the open subsets are  $\emptyset$ ,  $\{\eta\}$  and  $\{s, \eta\}$ , and let  $x = \{\eta\} \in F(X)$ . Then, if  $Y$  is a topological space, the map  $u(Y) : \text{Hom}_{\mathbf{Set}^{\text{op}}}(X, Y) = \text{Hom}_{\mathbf{Set}}(Y, X) \rightarrow \mathfrak{P}(Y)$  sends  $f : Y \rightarrow \{s, \eta\}$  to the open subset  $f^{-1}(\{\eta\})$  of  $Y$ . Conversely, if  $U$  is an open subset of  $Y$ , then the map  $f : Y \rightarrow \{s, \eta\}$  such that  $f(y) = \eta$  for  $y \in U$  and  $f(y) = s$  for  $y \in Y - U$  is continuous. So  $u(Y)$  is bijective.
- (k). For every  $k$ -vector space  $V$ , we have  $F(V) = \text{Hom}_k(V, k)$ . So we can take  $X = k$  (with the obvious action of  $k$ ) and  $x = \text{id}_k \in \text{Hom}_k(k, k)$ . Indeed, for every  $k$ -vector space  $V$ , the map  $u(V) : \text{Hom}_{\mathbf{Mod}_k^{\text{op}}}(k, V) = \text{Hom}_k(V, k) \rightarrow F(V) = \text{Hom}_k(V, k)$  sends  $f : V \rightarrow k$  to  $\text{id}_k \circ f = f$ . This is the identity of  $F(V)$ , so it is obviously bijective.

□

**9.** (extra credit) The simplicial category  $\Delta$  is defined in Example I.2.1.8(5) of the notes. It is the category whose objects are the finite sets  $[n] = \{0, 1, \dots, n\}$  with their usual order and whose morphisms are the nondecreasing maps between these sets.

The category  $\mathbf{sSet}$  of *simplicial sets* is  $\text{Func}(\Delta^{\text{op}}, \mathbf{Set})$ . So a simplicial set is by definition a functor  $X : \Delta^{\text{op}} \rightarrow \mathbf{Set}$ ; in that case, we write  $X_n$  for  $X([n])$  and, if  $f : [n] \rightarrow [m]$ , we often write  $f^* : X_m \rightarrow X_n$  for  $X(f)$ . For example, for each  $n \in \mathbb{N}$ , the *standard simplex* of dimension  $n$  is the simplicial set  $\text{Hom}_{\Delta}(\cdot, [n])$ .

If  $X$  is a simplicial set, a simplicial subset  $Y$  of  $X$  is the data of a subset  $Y_n$  of  $X_n$ , for every  $n \in \mathbb{N}$ , such that  $\alpha^*(Y_m) \subset Y_n$  for every morphism  $\alpha : [n] \rightarrow [m]$  in  $\Delta$ . We can form images of morphisms of simplicial sets, and unions and intersections of simplicial subsets, in the obvious

way.

If we see each poset  $[n]$  as a category in the usual way, then the morphisms of  $\Delta$  become functors, so this allows us to see  $\Delta$  as a subcategory of  $\mathbf{Cat}$ .

Let  $\mathcal{C}$  be a category. Its *nerve*  $N(\mathcal{C})$  is the restriction to  $\Delta^{\text{op}}$  of the functor  $\text{Hom}_{\mathbf{Cat}}(\cdot, \mathcal{C})$  on  $\mathbf{Cat}^{\text{op}}$ ; it is a functor from  $\Delta^{\text{op}}$  to  $\mathbf{Set}$ , i.e. a simplicial set. As  $\text{Hom}_{\mathbf{Cat}}$  is a bifunctor, this construction is functorial in  $\mathcal{C}$ , and we get a nerve functor  $N : \mathbf{Cat} \rightarrow \mathbf{sSet}$ .

- (a). (3 points) If  $\mathcal{C}$  is a category, show that  $N(\mathcal{C})_0 \simeq \text{Ob}(\mathcal{C})$  and  $N(\mathcal{C})_1 \simeq \coprod_{X, Y \in \text{Ob}(\mathcal{C})} \text{Hom}_{\mathcal{C}}(X, Y)$ . Can you give a similar description of  $N(\mathcal{C})_n$  for  $n \geq 2$ ?
- (b). (1 point) Let  $n \in \mathbb{N}$ . Show that the nerve of  $[n]$  is isomorphic to  $\Delta_n$ .
- (c). (1 point) Let  $n \in \mathbb{N}$ . Show that there exists  $e_n \in \Delta_n([n])$  such that, for every simplicial set  $X$ , the map  $\text{Hom}_{\mathbf{sSet}}(\Delta_n, X) \xrightarrow{\sim} X_n$  sending  $u$  to  $u_n(e_n)$  is bijective.
- (d). (1 point) For every category  $\mathcal{C}$  and every simplicial set  $X$ , if  $u, v : X \rightarrow N(\mathcal{C})$  are two morphisms of simplicial sets such that  $u_i, v_i : X_i \rightarrow N(\mathcal{C})_i$  are equal for  $i \in \{0, 1\}$ , show that  $u = v$ .
- (e). (1 point) We denote by  $\Delta_{\leq 2}$  the full subcategory of  $\Delta$  whose objects are  $[0]$ ,  $[1]$  and  $[2]$ ; if  $X$  is a simplicial set, we denote by  $X_{\leq 2}$  its restriction to  $\Delta_{\leq 2}$  (which is a functor  $\Delta_{\leq 2}^{\text{op}} \rightarrow \mathbf{Set}$ ).

Let  $X$  be a simplicial set and  $\mathcal{C}$  be a category. Show that every morphism  $X_{\leq 2} \rightarrow N(\mathcal{C})_{\leq 2}$  extends to a morphism  $X \rightarrow N(\mathcal{C})$ .

- (f). (2 points) Show that the functor  $N : \mathbf{Cat} \rightarrow \mathbf{sSet}$  is fully faithful.

Let  $n \in \mathbb{N}$ . For every  $k \in [n]$ , we denote by  $\delta_k$  the unique injective increasing map  $[n-1] \rightarrow [n]$  such that  $k \notin \text{Im}(\delta_k)$ . This induces a map  $\Delta_{n-1} \rightarrow \Delta_n$ , that we also denote by  $\delta_k$ ; the image of this map is called the  $k$ th facet of  $\Delta_n$ .

If  $k \in [n]$ , the *horn*  $\Lambda_k^n$  is the union of all the facets of  $\Delta_n$  except for the  $k$ th one; in other words, it is the simplicial subset of  $\Delta_n$  defined by

$$\Lambda_k^n([m]) = \{f \in \text{Hom}_{\Delta}([m], [n]) \mid \exists l \in [n] - \{k\} \text{ and } g \in \text{Hom}_{\Delta}([m], [n-1]) \text{ with } f = \delta_l \circ g\}.$$

- (g). (1 point) Let  $\mathcal{C}$  be a category. If  $n \geq 3$  and  $k \in [n] - \{0, n\}$ , show that every morphism of simplicial sets  $\Lambda_k^n \rightarrow X$  extends uniquely to a morphism  $\Delta_n \rightarrow X$ .
- (h). (1 point) Let  $\mathcal{C}$  be a category. Show that every morphism of simplicial sets  $\Lambda_1^2 \rightarrow X$  extends uniquely to a morphism  $\Delta_2 \rightarrow X$ .
- (i). (2 points) Show that a simplicial set  $X$  is the nerve of a category if and only if, for every  $n \in \mathbb{N}$ , every  $0 < k < n$  and every morphism of simplicial sets  $u : \Lambda_k^n \rightarrow X$ , the morphism  $u$  extends uniquely to a morphism  $\Delta_n \rightarrow X$ .

*Solution.*

- (a). By problem 8(f), the functor  $\mathbf{Cat} \rightarrow \mathbf{Set}$ ,  $\mathcal{C} \mapsto \text{Ob}(\mathcal{C})$  is represented by  $[0]$ . As  $N(\mathcal{C})_0 = \text{Hom}_{\mathbf{Cat}}([0], \mathcal{C})$ , this gives an isomorphism  $N(\mathcal{C})_0 \simeq \text{Ob}(\mathcal{C})$ , natural in  $\mathcal{C}$ . Similarly, by 8(g), the functor  $\mathbf{Cat} \rightarrow \mathbf{Set}$ ,  $\mathcal{C} \mapsto \coprod_{X, Y \in \text{Ob}(\mathcal{C})} \text{Hom}_{\mathcal{C}}(X, Y)$  is represented by  $[1]$ . As  $N(\mathcal{C})_1 = \text{Hom}_{\mathbf{Cat}}([1], \mathcal{C})$ , this gives an isomorphism  $N(\mathcal{C})_1 \simeq \coprod_{X, Y \in \text{Ob}(\mathcal{C})} \text{Hom}_{\mathcal{C}}(X, Y)$ , also natural in  $\mathcal{C}$ . Note that, if  $\delta_0, \delta_1 : [0] \rightarrow [1]$  are the two maps defined by  $\delta_0(0) = 1$  and  $\delta_1(0) = 0$ , then  $\delta_1^* : N(\mathcal{C})_1 \rightarrow N(\mathcal{C})_0$  sends a morphism to its source and  $\delta_0^* : N(\mathcal{C})_1 \rightarrow N(\mathcal{C})_0$  sends a morphism to its target.

Let  $\mathcal{C}$  be category. For  $n \geq 1$ , consider the set  $M_n$  of sequences of  $n$  composable morphisms  $c_0 \xrightarrow{f_1} c_1 \xrightarrow{f_2} \dots \xrightarrow{f_n} c_n$  of  $\mathcal{C}$ , which we will also write as  $(f_1, \dots, f_n)$ . We have a map  $\alpha : N(\mathcal{C})_n \rightarrow M_n$  sending a functor  $F : [n] \rightarrow \mathcal{C}$  to the sequence  $F(0) \rightarrow F(1) \rightarrow \dots \rightarrow F(n)$ , where the morphism  $F(i) \rightarrow F(i+1)$  is the image by  $F$  of the unique morphism  $i \rightarrow i+1$  in  $[n]$ . This uniquely determines the functor  $F$ , because, for  $i \leq j$  in  $[n]$ , the unique morphism  $i \rightarrow j$  is the composition of  $i \rightarrow i+1 \rightarrow i+2 \rightarrow \dots \rightarrow j$ . For the same reason, every element of  $M_n$  comes from a functor  $F : [n] \rightarrow \mathcal{C}$ . So we get a bijection  $N(\mathcal{C})_n \xrightarrow{\sim} M_n$ . (We can easily make  $M_n$  into a functor  $\mathbf{Cat} \rightarrow \mathbf{Set}$ , and then this bijection is an isomorphism of functors.)

We will identify  $N(\mathcal{C})_n$  with  $M_n$  in the rest of this solution. We also write  $M_0 = \text{Ob}(\mathcal{C})$ . (we can think of  $c \in \text{Ob}(\mathcal{C})$  as a length 0 sequence of composable morphisms  $(c)$ .)

Let  $\alpha : [m] \rightarrow [n]$  be a nondecreasing map. We can give an explicit description of the map  $\alpha^* : N(\mathcal{C})_n \rightarrow N(\mathcal{C})_m$  by chasing through the identifications. If  $n = 0$  and  $m \geq 1$ , then  $\alpha^*$  sends  $c \in \text{Ob}(\mathcal{C})$  to the sequence  $(\text{id}_c, \dots, \text{id}_c) \in M_m$ . If  $n \geq 1$  and  $m = 0$ , let  $i = \alpha(0)$ ; then  $\alpha^*$  sends the sequence  $c_0 \xrightarrow{f_1} c_1 \xrightarrow{f_2} \dots \xrightarrow{f_n} c_n$  to  $c_i$ . Suppose that  $n, m \geq 1$ , let  $c_0 \xrightarrow{f_1} c_1 \xrightarrow{f_2} \dots \xrightarrow{f_n} c_n$  be an element of  $M_n$ , and let  $d_0 \xrightarrow{g_1} d_1 \xrightarrow{g_2} \dots \xrightarrow{g_m} d_m$  be its image by  $\alpha^*$ . For  $i \in \{1, \dots, m\}$ , we have :

- if  $\alpha(i-1) = \alpha(i)$ , then  $d_{i-1} = d_i = c_{\alpha(i)}$  and  $g_i = \text{id}_{c_{\alpha(i)}}$ ;
- if  $\alpha(i-1) < \alpha(i)$ , then  $g_i = f_{\alpha(i)} \circ \dots \circ f_{\alpha(i-1)+1}$ .

- (b). As we have identified  $\Delta$  to a subcategory of  $\mathbf{Cat}$ , this is just the definition of  $\Delta_n$ .
- (c). Let  $e_n = \text{id}_{[n]} \in \Delta_n([n]) = \text{Hom}_\Delta([n], [n])$ . The fact that the map of the statement is bijective is exactly the Yoneda lemma (Theorem I.3.2.2 of the notes).
- (d). Suppose that  $u, v : X \rightarrow N(\mathcal{C})$  satisfy the condition of the question. Let  $n \geq 2$ , and let  $x \in X_n$ . Write  $u(x) = (c_0 \xrightarrow{f_1} c_1 \xrightarrow{f_2} \dots \xrightarrow{f_n} c_n)$  and  $v(x) = (d_0 \xrightarrow{g_1} d_1 \xrightarrow{g_2} \dots \xrightarrow{g_n} d_n)$ . We want to show that  $u(x) = v(x)$ , that is, that  $f_i = g_i$  for every  $i \in \{1, \dots, n\}$ . Fix  $i \in \{1, \dots, n\}$ , and consider the map  $\alpha : [1] \rightarrow [n]$  sending 0 to  $i-1$  and 1 to  $i$ . Then  $\alpha$  is a morphism in  $\Delta$ , so we have a commutative diagram

$$\begin{array}{ccc} X_n & \xrightarrow{u_n} & N(\mathcal{C})_n \\ \alpha^* \downarrow & & \downarrow \alpha^* \\ X_1 & \xrightarrow{u_1} & N(\mathcal{C})_1 \end{array}$$

and a similar commutative diagram for  $v$ . By definition of the bijection  $N(\mathcal{C})_n \xrightarrow{\sim} M_n$ , the map  $\alpha^*$  sends a sequence  $e_0 \xrightarrow{h_1} e_1 \xrightarrow{h_2} \dots \xrightarrow{h_n} e_n$  to  $h_i : e_{i-1} \rightarrow e_i$ . So we get  $f_i = \alpha^*(u_n(x)) = u_1(\alpha^*(x)) = v_1(\alpha^*(x)) = \alpha^*(v_n(x)) = g_i$ .

- (e). Let  $u : X_{\leq 2} \rightarrow N(\mathcal{C})_{\leq 2}$ . We want to show that  $u$  extends to a morphism of simplicial sets  $v : X \rightarrow N(\mathcal{C})$ . The solution of question (d) tells us how we must extend  $u$ : Let  $n \geq 2$ , and, for  $i \in \{1, \dots, n\}$ , let  $\alpha_i^n = \alpha_i : [1] \rightarrow [n]$  be the map  $m \mapsto m + i - 1$ . Then, for every  $x \in X_n$ ,  $v_n(x)$  must be the sequence  $(u_1(\alpha_1^n(x)), \dots, u_1(\alpha_n^n(x)))$  of morphisms of  $\mathcal{C}$ . These morphisms are composable : indeed, if we denote by  $\delta_0, \delta_1 : [0] \rightarrow [1]$  the two maps defined by  $\delta_0(0) = 1$  and  $\delta_1(0) = 0$ , then  $\alpha_i \circ \delta_0 = \alpha_{i+1} \circ \delta_1$  for  $1 \leq i \leq n-1$ , so the target  $u_0(\delta_0^* \alpha_i^n(x))$  of  $u_1(\alpha_i^n(x))$  is equal to the source  $u_0(\delta_1^* \alpha_{i+1}^n(x))$  of  $u_1(\alpha_{i+1}^n(x))$ .

We have to check that  $v_2 = u_2$  and that  $v$  is a morphism of simplicial sets. The proof that  $v_2 = u_2$  is exactly as in the solution of (d). To show that  $v$  is a morphism of simplicial sets, we take a nondecreasing map  $\alpha : [m] \rightarrow [n]$  and we show that  $v_m \circ \alpha^* = \alpha^* \circ v_n$ .

We can write  $\alpha = \alpha' \circ \alpha''$  with  $\alpha', \alpha''$  both nondecreasing,  $\alpha'$  injective and  $\alpha''$  surjective, and it suffices to show the statement for  $\alpha'$  and  $\alpha''$ . Moreover, we can write  $\alpha'$  (resp.  $\alpha''$ ) as a composition of injective (resp. surjective) nondecreasing maps  $[p] \rightarrow [p+1]$  (resp.  $[p+1] \rightarrow [p]$ ). So we may assume that  $\alpha$  is injective or surjective and that  $n = m \pm 1$ .

Suppose first that  $\alpha : [n+1] \rightarrow [n]$  is a surjective nondecreasing map. Then there is a unique  $i \in [n]$  such that  $\alpha(i) = \alpha(i+1) = i$ ,  $\alpha(j) = j$  for  $0 \leq j < i$  and  $\alpha(j) = j-1$  for  $i+1 < j \leq n+1$ . Let  $x \in X_n$ , and let  $(f_1, \dots, f_n) = v_n(x)$ . The map  $\alpha^* : N(\mathcal{C})_n \rightarrow N(\mathcal{C})_{n+1}$  sends the sequence of composable morphisms  $(f_1, \dots, f_n)$  to  $(f_1, \dots, f_i, \text{id}_c, f_{i+1}, \dots, f_n)$ , where  $c$  is the target of  $f_i$ . By definition,  $v_{n+1}(\alpha^*(x)) = (g_1, \dots, g_{n+1})$ , with  $g_j = u_1(\alpha_j^{n+1} \alpha^*(x))$ . If  $1 \leq j \leq i$ , then  $\alpha \circ \alpha_j^{n+1} = \alpha_j^n$ , so  $g_j = f_j$ . If  $i+2 \leq j \leq n+1$ , then  $\alpha \circ \alpha_j^{n+1} = \alpha_{j-1}^n$ , so  $g_j = f_{j-1}$ . Finally,  $\alpha_{i+1}^n \circ \alpha : [1] \rightarrow [n]$  is the map sending every element of  $[1]$  to  $i$ , so it is equal to  $\alpha' \circ \alpha''$ , where  $\alpha' : [0] \rightarrow [n]$  sends 0 to  $i$  and  $\alpha'' : [1] \rightarrow [0]$  is the unique map; so  $g_{i+1} = u_1(\alpha''^* \circ \alpha'^*(x)) = \alpha''^* u_0(\alpha'^*(x))$  is  $\text{id}_c$ , where  $c' = u_0(\alpha'^*(x))$ ; as  $\alpha' = \alpha_i^n \circ \delta_0$ , we have  $c' = \delta_0^*(f_i)$ , that is,  $c'$  is the target  $c$  of  $f_i$ , as we wanted.

Now we take  $\alpha : [n-1] \rightarrow [n]$  injective and increasing; we may also assume  $n \geq 3$ , as we already have the result for  $n \leq 2$ . There exists  $i \in [n]$  such that  $\text{Im}(\alpha) = [n] - \{i\}$ , that is, such that  $\alpha$  is the map  $\delta_i$  defined before (g). Let  $x \in X_n$ , and let  $c_0 \xrightarrow{f_1} \dots \xrightarrow{f_n} c_n$  be  $v_n(x)$  and  $(g_1, \dots, g_{n-1})$  be  $v_{n-1}(\alpha^*(x))$ . As we saw in the solution of (a), the map  $\alpha^* : M_n \rightarrow M_{n-1}$  sends the sequence of composable morphisms  $c_0 \xrightarrow{f_1} \dots \xrightarrow{f_n} c_n$  to the sequence :

- $c_0 \xrightarrow{f_1} \dots \xrightarrow{f_{n-1}} c_{n-1}$  if  $i = n$ ;
- $c_1 \xrightarrow{f_2} \dots \xrightarrow{f_n} c_n$  if  $i = 0$ ;
- $c_0 \xrightarrow{f_1} \dots \xrightarrow{f_{i-1}} c_{i-1} \xrightarrow{f_{i+1} \circ f_i} c_{i+1} \dots \xrightarrow{f_n} c_n$  if  $1 \leq i \leq n-1$ .

If  $1 \leq j \leq i-1$ , we have  $\alpha \circ \alpha_j^{n-1} = \alpha_j^n$ , which implies that  $g_j = f_j$ . If  $i+1 \leq j \leq n-1$ , we have  $\alpha \circ \alpha_j^{n-1} = \alpha_{j+1}^n$ , which implies that  $g_j = f_{j+1}$ . To finish the proof that  $\alpha^*(v_n(x)) = v_{n-1}(\alpha^*(x))$ , it remains to consider the case  $j = i \in \{1, \dots, n-1\}$ . Then  $\alpha \circ \alpha_j^{n-1} = \alpha' \circ \delta_1$ , where  $\alpha' : [2] \rightarrow [n]$  is the map  $x \mapsto x + i - 1$  and  $\delta_1 : [1] \rightarrow [2]$  is the map sending 0 to 0 and 1 to 2. Hence  $g_j = u_1(\delta_1^* \alpha'^*(x)) = \delta_1^* u_2(\alpha'^*(x))$ , so if  $u_2(\alpha'^*(x)) = (h_1, h_2)$  then  $g_j = h_2 \circ h_1$ ; but it is easy to see that  $u_2(\alpha'^*(x)) = (f_i, f_{i+1})$  (by looking at the composition of  $\alpha'$  with  $\alpha_1^2, \alpha_2^2 : [1] \rightarrow [2]$ ), so we are done.

- (f). Let  $\mathcal{C}$  and  $\mathcal{C}'$  be categories. We want to show that the map  $N : \text{Func}(\mathcal{C}, \mathcal{C}') \rightarrow \text{Hom}_{\mathbf{sSet}}(N(\mathcal{C}), N(\mathcal{C}'))$  is bijective, so we try to construct an inverse of this map.

Let  $u : N(\mathcal{C}) \rightarrow N(\mathcal{C}')$  be a morphism of simplicial sets. We denote by  $F$  the map  $\text{Ob}(\mathcal{C}) \simeq N(\mathcal{C})_0 \xrightarrow{u_0} N(\mathcal{C}')_0 \simeq \text{Ob}(\mathcal{C}')$ . Let  $f : c_0 \rightarrow c_1$  be a morphism of  $\mathcal{C}$ . We saw in (a) that this morphism corresponds to a functor  $T : [1] \rightarrow \mathcal{C}$ , that is, an element of  $N(\mathcal{C})_1$ . We denote by  $F(f) : d_0 \rightarrow d_1$  the morphism of  $\mathcal{C}'$  corresponding to  $u_1(T) \in N(\mathcal{C}')_1$ . We want to show that  $d_0 = F(c_0)$  and  $d_1 = F(c_1)$ . Let  $i \in \{0, 1\}$ , and consider the map  $\alpha : [0] \rightarrow [1]$  sending 0 to  $i$ . This is a morphism of  $\Delta$ , and  $\alpha^* : N(\mathcal{C})_1 \rightarrow N(\mathcal{C})_0$  sends a morphism of  $\mathcal{C}$  to its source if  $i = 0$  and its target if  $i = 1$ . Using the commutativity of the diagram

$$\begin{array}{ccc} N(\mathcal{C})_1 & \xrightarrow{u_1} & N(\mathcal{C}')_1 \\ \alpha^* \downarrow & & \downarrow \alpha^* \\ N(\mathcal{C})_0 & \xrightarrow{u_0} & N(\mathcal{C}')_0 \end{array}$$

we see that  $d_0 = F(c_0)$  and  $d_1 = F(c_1)$ . Now we show that  $F$  is a functor. There are two conditions to check :

- (1) Consider the unique map  $\alpha : [1] \rightarrow [0]$ . This is a morphism of  $\Delta$ , and  $\alpha^* : N(\mathcal{C})_0 \rightarrow N(\mathcal{C})_1$  sends the element of  $N(\mathcal{C})_0$  corresponding to an object  $c$  of  $\mathcal{C}$  to the element of  $N(\mathcal{C})_1$  corresponding to  $\text{id}_c$ . As  $u_1 \circ \alpha^* = \alpha^* \circ u_0$ , we get that, for every  $c \in \text{Ob}(\mathcal{C})$ ,  $F(\text{id}_c) = \text{id}_{F(c)}$ .
- (2) Consider the map  $\alpha : [1] \rightarrow [2]$  sending 0 to 0 and 1 to 2, and the map  $\sigma_i : [1] \rightarrow [2]$ ,  $m \mapsto m + i$ , for  $i \in \{0, 1\}$ . Then  $\alpha^*$  (resp.  $\sigma_0^*$ , resp.  $\sigma_1^*$ ) sends the element of  $N(\mathcal{C})_2$  corresponding to a sequence  $c_0 \xrightarrow{f_1} c_1 \xrightarrow{f_2} c_2$  to the element of  $N(\mathcal{C})_1$  corresponding to  $f_2 \circ f_1$  (resp.  $f_1$ , resp.  $f_2$ ). (This is clear on the identifications of (a).)

Let  $c_0 \xrightarrow{f_1} c_1 \xrightarrow{f_2} c_2$  be a sequence of composable morphisms of  $\mathcal{C}$ . Using the previous paragraph and the fact that  $u$  is a morphism of functors, we see that the image by  $u$  of the element of  $N(\mathcal{C})_2$  corresponding to this sequence is the sequence  $F(c_0) \xrightarrow{F(f_1)} F(c_1) \xrightarrow{F(f_2)} F(c_2)$ , and using this and the fact that  $\alpha^* \circ u_2 = u_1 \circ \alpha^*$ , we finally get  $F(f_2) \circ F(f_1) = F(f_2 \circ f_1)$ .

So we have constructed a map  $\Phi : \text{Hom}_{\mathbf{sSet}}(N(\mathcal{C}), N(\mathcal{C}')) \rightarrow \text{Func}(\mathcal{C}, \mathcal{C}')$ , and it is clear on the construction that, for every functor  $F : \mathcal{C} \rightarrow \mathcal{C}'$ , we have  $\Phi(N(F)) = F$ . Now let  $u : N(\mathcal{C}) \rightarrow N(\mathcal{C}')$  be a morphism of simplicial sets, and let  $F = \Phi(u)$ . We want to show that  $N(F) = u$ . Again, it is clear from the construction of  $\Phi$  that  $N(F)_0 = u_0$  and  $N(F)_1 = u_1$ . But then the fact that  $N(F) = u$  follows from (c).

- (g). If  $n \geq 4$ , then, for every  $k \in [n]$ , the morphism  $\Lambda_{k, \leq 2}^n \rightarrow \Delta_{n, \leq 2}$  induced by the inclusion  $\Lambda_k^n \subset \Delta_n$  is the identity morphism. So, by (d) and (e), every morphism  $\Lambda_k^n \rightarrow N(\mathcal{C})$  extends uniquely to a morphism  $\Delta_n \rightarrow N(\mathcal{C})$ .

We still need to treat the case  $n = 3$ . Note that the uniqueness of the extension will follow from the fact that  $\Lambda_{k, \leq 1}^3 = \Delta_{3, \leq 1}$ .

Let  $\partial\Delta_3$  be the union of all the faces of  $\Delta_3$ . Then the inclusion  $\partial\Delta_3 \subset \Delta_3$  induces an equality  $\partial\Delta_{3, \leq 2} = \Delta_{3, \leq 2}$ , so it suffices to show that the morphism  $u : \Lambda_k^3 \rightarrow N(\mathcal{C})$  extends to  $\partial\Delta_3$ . As  $\partial\Delta_3 = \Lambda_k^3 \cup \delta_{k*}(\Delta_2)$  and  $\Lambda_k^3 \cap \delta_{k*}(\Delta_2) = \delta_{k*}(\partial\Delta_2)$ , it suffices to extend  $u$  from  $\delta_{k*}(\partial\Delta_2)$  to  $\delta_{k*}(\Delta_2)$ . For  $0 \leq i < j \leq 3$ , let  $\alpha_{i,j} : [1] \rightarrow [3]$  be the map sending 0 to  $i$  and 1 to  $j$ ; note that  $\alpha_{i,j} \in \Lambda_k^3([3])$ . Let  $f_i = u_3(\alpha_{i-1,i})$ , for  $1 \leq i \leq 3$ . We treat the case  $k = 1$ , the case  $k = 2$  is similar. Factoring both  $\alpha_{1,2}$  and  $\alpha_{2,3}$  through the morphism  $\delta_0 : [2] \rightarrow [3]$ , we see that  $f_3$  and  $f_2$  are composable, and that  $f_3 \circ f_2 = u_3(\alpha_{1,3})$ . Factoring both  $\alpha_{0,1} = f_1$  and  $\alpha_{1,3}$  through the morphism  $\delta_2 : [2] \rightarrow [3]$ , we see that  $f_3 \circ f_2 = u_3(\alpha_{1,3})$  and  $f_1$  are composable, and that  $f_3 \circ f_2 \circ f_1 = u_3(\alpha_{0,3})$ . Similarly, using  $\delta_3 : [2] \rightarrow [3]$ , we show that  $f_2 \circ f_1 = u_3(\alpha_{0,2})$ .

In particular, we see that  $u_3(\alpha_{0,3}) = f_3 \circ f_2 \circ f_1 = u_3(\alpha_{3,2}) \circ u_3(\alpha_{0,2})$ . This is exactly the condition that we need to extend  $u$  from  $\delta_{1*}(\partial\Delta_2)$  to  $\delta_{1*}(\Delta_2)$ . (See the solution of the next question.)

- (h). Let  $u : \Lambda_1^2 \rightarrow N(\mathcal{C})$ . We want to extend  $u$  to a morphism  $v : \Delta_2 \rightarrow N(\mathcal{C})$ . Remember that, by the Yoneda lemma, giving  $v$  is the same as giving an element  $e$  of  $N(\mathcal{C})_2$ ; the fact that  $v$  extends  $u$  then says that, for every  $\alpha : [m] \rightarrow [2]$  such that  $\alpha \in \Lambda_1^2([m])$ , we have  $\alpha^*(e) = u_m(\alpha)$ .

Note that the maps  $\delta_2$  and  $\delta_0$  from  $[1]$  to  $[2]$  are in  $\Lambda_1^2([1])$  by definition of the horn  $\Lambda_1^2$ . We set  $f_1 = u_1(\delta_2)$  and  $f_2 = u_1(\delta_0)$ . Comparing the compositions of  $\delta_0$  and  $\delta_2$  with the two maps  $[0] \rightarrow [1]$ , we see that  $(f_1, f_2)$  is a sequence of composable morphisms of  $\mathcal{C}$ , hence an

element of  $e$  of  $N(\mathcal{C})_2$ ; we denote by  $v : \Delta_2 \rightarrow N(\mathcal{C})$  the corresponding morphism, that is, the unique morphism such that  $v_2(e_2) = (f_1, f_2)$ . Using the method of the solution of (d), we see that this is the only possibility for a morphism extending  $u$  (such a morphism must send  $e_2 \in \Delta_2([2])$  to  $(f_1, f_2)$ ).

It remains to show that  $v$  does extend  $u$ . Let  $\alpha : [m] \rightarrow [2]$  be an element of  $\Lambda_1^2([2])$ ; by definition of the horn, this means that we can write  $\alpha = \delta_i \circ \beta$ , with  $\beta : [m] \rightarrow [1]$  nondecreasing and  $i \in \{0, 2\}$ . Then  $v_2(\alpha) = \alpha^*(e) = \beta^*(\delta_i^*(e))$  and  $u_2(\alpha) = \beta^*u_1(\delta_i)$ , so it suffices to show that  $\delta_i^*(e) = u_1(\delta_i)$ ; but this follows from the definition of  $f_1$  and  $f_2$  and the description of  $\delta_i^* : N(\mathcal{C})_2 \rightarrow N(\mathcal{C})_1$  in (a).

(i). Let  $X$  be a simplicial set, and suppose that every morphism  $u : \Lambda_k^n \rightarrow X$  with  $0 < k < n$  extends uniquely to  $\Delta_n$ . We denote by  $d_0, d_1 : [0] \rightarrow [1]$  the two maps sending 0 to 0 and 1 respectively, and by  $s$  the unique map from  $[1]$  to  $[0]$ . We construct a category  $\mathcal{C}$  in the following way :

- (1) We take  $\text{Ob}(\mathcal{C}) = X_0$ .
- (2) If  $c, d \in X_0$ , we have  $\text{Hom}_{\mathcal{C}}(c, d) = \{f \in X_1 \mid d_0^*(f) = c \text{ and } d_1^*(f) = d\}$ .
- (3) For every  $c \in X_0$ , we denote by  $\text{id}_c$  the element  $s^*(c)$  of  $X_1$ .
- (4) Let  $c, d, e \in X_0$  and  $f \in \text{Hom}_{\mathcal{C}}(c, d)$ ,  $g \in \text{Hom}_{\mathcal{C}}(d, e)$ . We want to construct a morphism  $u : \Lambda_1^2 \rightarrow X$ . Let  $\alpha : [m] \rightarrow [2]$  be an element of  $\Lambda_1^2([m])$ . By definition of  $\Lambda_1^2$ , there exists  $\beta : [m] \rightarrow [1]$  and  $j \in \{0, 2\}$  such that  $\alpha = \delta_j \circ \beta$ . We set  $u_m(\alpha) = \beta^*(f_j)$ , with  $f_j = f$  if  $j = 2$  and  $f_j = g$  if  $j = 0$ . We must check that this is well-defined; if  $\alpha$  can be written as  $\beta \circ \delta_0$  and  $\beta' \circ \delta_2$ , with  $\beta : [m] \rightarrow [1]$ , this means that  $\text{Im}(\alpha) = \{1\}$ , so  $\text{Im}(\beta) = \{0\}$  and  $\text{Im}(\beta') = \{1\}$ , so there exists  $\gamma : [m] \rightarrow [0]$  such that  $\beta = d_0 \circ \gamma$  and  $\beta' = d_1 \circ \gamma$ , hence  $\beta^*(g) = \gamma^*(d) = \beta'^*(f)$ . We now check that  $u$  is a morphism of simplicial sets. If  $\alpha : [m] \rightarrow [2]$  is an element of  $\Lambda_1^2([m])$ , write  $\alpha = \delta_j \circ \beta$ , with  $\beta : [m] \rightarrow [1]$  and  $j \in \{0, 2\}$ ; then, for every  $\gamma : [m'] \rightarrow [m]$ , we have  $\alpha \circ \gamma = \delta_j \circ (\beta \circ \gamma)$ , so  $u_{m'}(\gamma) = (\beta \circ \gamma)^*(f_j) = \gamma^*(\beta^*(f_j)) = \gamma^*(u_m(\beta))$ . So  $u$  is a morphism of simplicial sets, and, by assumption, it extends uniquely to a morphism  $v : \Delta_2 \rightarrow X$ . We take  $g \circ f = v_1(\delta_1)$ . It is easy to check that  $g \circ f \in \text{Hom}_{\mathcal{C}}(c, e)$ .

It is easy to check that the identity morphisms are unit elements for the composition.

We check that the composition law of  $\mathcal{C}$  is associative. Let  $(f_1, f_2, f_3)$  be a sequence of composable morphisms in  $\mathcal{C}$ . Remember that we have maps  $\delta_i : [2] \rightarrow [3]$ , inducing morphisms of simplicial sets  $\delta_{i*} : \Delta_2 \rightarrow \Delta_3$ . As in the construction of the composition in (4), we use the pair  $(f_1, f_2)$  to construct a morphism  $\delta_{3*}(\Delta_2) \rightarrow X$ , the pair  $(f_2, f_3)$  to construct a morphism  $\delta_{1*}(\Delta_2) \rightarrow X$ , and the pair  $(f_1, f_3 \circ f_2)$  to construct a morphism  $\delta_{2*}(\Delta_2) \rightarrow X$ . These three morphisms glue to a morphism  $\Lambda_1^3 \rightarrow X$ , which extends uniquely to  $v : \Delta_3 \rightarrow X$ . In particular, if we define maps  $\alpha_{i,j} : [1] \rightarrow [3]$  as in (g), we see as in that question that

$$v_1(\alpha_{0,3}) = v_1(\alpha_{3,2}) \circ v_1(\alpha_{2,0}) = f_3 \circ (f_2 \circ f_1) = v_1(\alpha_{3,1} \circ \alpha_{1,0}) = (f_3 \circ f_2) \circ f_1.$$

For  $n \geq 1$  and  $1 \leq i \leq n$ , let the  $\alpha_i^n : [1] \rightarrow [n]$  be as in the solution of (e).

Let  $n \geq 2$ . If  $1 \leq m \leq n$   $0 \leq i_0 < i_1 < \dots < i_m \leq n$ , we denote by  $\alpha_{i_0, \dots, i_m} : [1] \rightarrow [n]$  the map sending  $r \in [m]$  to  $i_r \in [n]$ . If  $x \in X_n$ , we define morphisms  $f_{1,x}, \dots, f_{n,x}$  in  $\mathcal{C}$  by  $f_i = \alpha_i^{n*}(x)$ . As  $\alpha_i^n \circ d_0 = \alpha_{i+1}^n \circ d_1$  for  $1 \leq i \leq n-1$ , the  $f_i$  form a sequence of composable morphisms, so  $(f_1, \dots, f_n) \in N(\mathcal{C})_n$ . We claim that :

- (A) For every  $(f_1, \dots, f_n) \in N(\mathcal{C})_n$ , there exists  $x \in X$  such that  $(f_{1,x}, \dots, f_{n,x}) = (f_1, \dots, f_n)$ .

(B) If  $x, y \in X_n$  are such that  $f_{i,x} = f_{i,y}$  for  $1 \leq i \leq n$ , then  $x = y$ .

We prove (A). Let  $(f_1, \dots, f_n)$  be a sequence of composable morphisms in  $\mathcal{C}$ . For  $0 \leq i_0 < i_1 \leq n$ , we define a morphism  $u_{i_0, i_1} : \alpha_{i_0, i_1, *}( \Delta_1 ) \rightarrow X$  by sending  $\alpha_{i_0, i_1, *}(e_1)$  to  $f_{i_1} \circ \dots \circ f_{i_0+1} \in X_1$ , where  $e_1 \in \Delta_1([1])$  is the element defined in (c). Suppose that  $0 \leq i_0 < i_1 < i_2 \leq n$ . Then the morphisms  $u_{i_0, i_1}$ ,  $u_{i_1, i_2}$  and  $u_{i_0, i_2}$  agree on the intersections of their domains (because the  $f_i$  are composable), so they glue to a morphism  $u'_{i_0, i_1, i_2} : \alpha_{i_0, i_1, i_2, *}( \partial \Delta_2 ) \rightarrow X$ ; by the property of  $X$ , this morphism extends uniquely to a morphism  $u_{i_0, i_1, i_2} : \alpha_{i_0, i_1, i_2, *}( \Delta_2 ) \rightarrow X$ . Now take  $0 \leq i_0 < i_1 < i_2 < i_3 \leq n$ . Then the morphisms  $u_{i_0, i_1, i_2}$ ,  $u_{i_0, i_1, i_3}$ ,  $u_{i_0, i_2, i_3}$  and  $u_{i_1, i_2, i_3}$  agree on the intersections of their domains (we just recover one of the  $u_{i_r, i_s}$  on such an intersection), so they glue to a morphism  $u'_{i_0, i_1, i_2, i_3} : \alpha_{i_0, i_1, i_2, i_3, *}( \partial \Delta_3 ) \rightarrow X$ ; by the property of  $X$ , this morphism extends uniquely to a morphism  $u_{i_0, i_1, i_2, i_3} : \alpha_{i_0, i_1, i_2, i_3, *}( \Delta_3 ) \rightarrow X$ . We continue in this way until we get a morphism  $u_{0, 1, \dots, n} : \Delta_n \rightarrow X$  extending the original  $u_{i_0, i_1}$ ; the corresponding element of  $X_n$  has the required property.

Now we prove (B). By the Yoneda lemma, the elements  $u, v \in X_n$  correspond to two morphisms  $u_x, u_y : \Delta_n \rightarrow X$ , and the condition of (B) says that  $u_x$  and  $u_y$  agree on  $\alpha_{i_0, i_1, *}( \Delta_1 )$  for all  $i_0, i_1 \in [n]$  such that  $i_0 < i_1$ . But we saw in the proof of (A) that there is a unique way to extend a family of morphisms  $\alpha_{i_0, i_1, *}( \Delta_1 ) \rightarrow X$  (agreeing on the intersections of the  $\alpha_{i_0, i_1, *}( \Delta_1 )$ ) to a morphism  $\Delta_n \rightarrow X$ . So  $u_x = u_y$ , that is,  $x = y$ .

Finally, we define  $u : X \rightarrow N(\mathcal{C})$  by taking  $u_1$  and  $u_0$  to be the obvious bijections and by sending  $x \in X_n$  to the sequence of maps  $(f_{1,x}, \dots, f_{n,x})$ , for every  $n \geq 2$ . This induces a morphism of functors from  $X_{\leq 2}$  to  $N(\mathcal{C})_{\leq 2}$  by the definition of the composition and the description of the maps between the  $N(\mathcal{C})_n$  in (a). Then the solution of (e) shows that  $u$  is a morphism of simplicial sets. Points (A) and (B) imply that  $u$  is an isomorphism. □