

Minors of plane digraphs

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Abstract

A digraph H is a “semi-strong minor” of another, G , if a subdivision of H can be obtained from a subdigraph of G by contracting strongly-connected subdigraphs to single vertices. We will define a width measure of “plane” digraphs (that is, drawn in the plane) based on a kind of branch-composition, and show that for every plane digraph H , all plane digraphs not containing H as a semi-strong minor have bounded width, while plane digraphs in general have unbounded width.

1 Introduction

A theorem of Robertson and Seymour [7] states that:

1.1 *For every $k \geq 0$ there exists $\ell \geq 0$ such that every graph with tree-width at least ℓ contains a $k \times k$ grid as a minor. Conversely, for every $\ell \geq 0$ there exists $k \geq 0$ such that every graph that contains a $k \times k$ grid as a minor has tree-width at least ℓ .*

(We omit the definitions of “grid”, “minor” and “tree-width” since they are well known. See Figure 1.) See [3] for a simpler proof, and [2] for the best dependence between grid size and tree-width known.

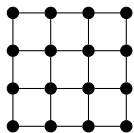


Figure 1: A 4×4 grid.

The result 1.1 is equivalent to saying that for every planar graph H , all graphs not containing H as a minor have bounded tree-width, because every planar graph is a minor of a grid. We are interested in extending this to digraphs.

That raises several issues: what is “minor containment” for digraphs? What is “tree-width” for digraphs? What is a “grid” for digraphs? But at least we can avoid the first of these issues for the moment. One can recast the theorem of [7] as a statement about the undirected graphs that do not contain a subdivision of a large “wall” as a subgraph, and the same idea will work for us in digraphs.

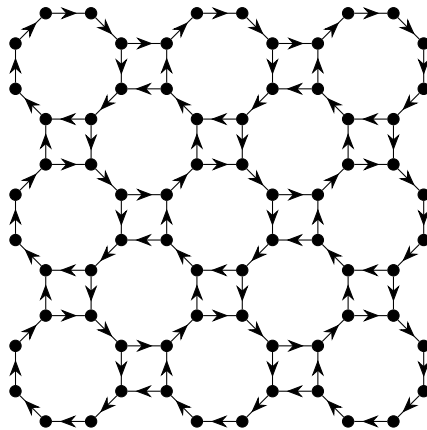


Figure 2: The 6×6 diwall.

Let us say a digraph G *embeds* a digraph H if there is a subdigraph of G that is isomorphic to a subdivision of H . For $k \geq 2$, even, let us say a $k \times k$ *diwall* is a digraph as illustrated in Figure 2. The subdigraph formed by the horizontal and diagonal edges is the disjoint union of k directed paths, alternating in direction, that we call the “horizontal” paths; and similarly there are k “vertical” paths. Each horizontal path intersects each vertical path in exactly one edge (the diagonal edges in the figure). If we contract the diagonal edges we obtain a $k \times k$ grid called an *alternating grid* that we will discuss later.

Our theorem concerns the structure of digraphs drawn in the 2-sphere that do not embed a large diwall.

Before we begin with more precise statements and definitions, let us give an idea of what we mean by “structure” here. For undirected graphs, Menger’s theorem involves a vertex-cutset of bounded size, but for digraphs, it involves a cutset containing a bounded number of vertices and an unbounded number of edges, all crossing the cutset in the wrong direction. We have something similar. For undirected graphs, bounded tree-width means that the graph can be cut up by “non-crossing” vertex-cutsets into trivial pieces, where the cutsets have bounded size. Our theorem is the same, except that the cutsets now contain both a bounded number of vertices and an unbounded number of edges crossing the cutset in prescribed directions. The edges in the cutset fall into a bounded number of “intervals”, and all the edges in one interval cross the cutset in the same direction. For “intervals” to make sense, we assume that we are working with cutsets in a digraph drawn in the plane, so there is an associated circular order on the elements of each cutset, and the intervals are intervals in this order. For this reason, our result only gives information about planar digraphs. We will prove that for every k , every digraph drawn in the plane that does not embed the $k \times k$ diwall can be cut up into trivial pieces by non-crossing cutsets, each consisting of a bounded number of vertices and a bounded number of intervals of edges as discussed.

The proof is in two stages. First we assume that G is drawn in the plane, and does not embed the $k \times k$ diwall, and in addition we have a bound on the “interleaving” of G , that is, for each vertex v , the edges incident with v fall into a bounded number of intervals in their cyclic order around v , and in each interval all edges enter v or they all leave v . In this case, the theorem is simpler, and we do not need to allow the cutsets to contain vertices; we can cut the graph into trivial (that is, one-vertex) pieces by a non-crossing collection of edge-cutsets, where the edges in each cutset fall into a bounded number of intervals as before. Then we deduce the stronger theorem, when there is no bound on interleaving.

Later in the paper, we discuss the connection with the well-known digraph grid theorem of Kawarabayashi and Kreutzer [5, 6]; and also deduce some consequences for excluding general planar digraphs as minors of large planar digraphs.

2 Drawings

We have omitted several definitions, and been sloppy about some others, in order to state the main result quickly; but now we need to go back and define things more carefully. So we will basically start again. Graphs and digraphs in this paper are all finite, and may have parallel edges or loops. Throughout the paper, Σ denotes the 2-sphere, and we assign an orientation “clockwise” to Σ . An *O-arc* in Σ is a subset homeomorphic to a circle, and a *line* in Σ is a subset homeomorphic to $[0, 1]$. We will not consider drawings in general surfaces; in this paper a *drawing* G is a graph, such that:

- $V(G) \subseteq \Sigma$;
- each $e = uv \in E(G)$ is either a line with ends $u \neq v$ in $V(G)$, or an O-arc containing $u = v$;
- if $v \in V(G)$ and $e \in E(G)$ are not incident in G then $v \notin e$; and
- if $e, f \in E(G)$ are distinct, every point in $e \cap f$ equals v for some $v \in V(G)$ incident with both e, f .

We will only consider drawings with no topological pathologies. A *region* is a maximal arc-wise connected subset of Σ disjoint from the drawing, and hence is an open set. If a drawing G is a digraph, we say G is a *didrawing*. Two drawings or didrawings are *homeomorphic* if there is a homeomorphism of Σ to itself that takes one to the other. If G is a digraph, we say a digraph H is a *subdivision* of G if H can be obtained from G by repeatedly replacing an edge uv with a new vertex w and two edges uw, wv (u, v might be equal). Let us say a digraph G *embeds* a digraph H if there is a subdigraph of G that is isomorphic to a subdivision of H . Similarly, a didrawing H is a Σ -*subdivision* of a didrawing G if H can be obtained from G by repeatedly replacing an edge uv with a new vertex w and two edges uw, wv , such that $uw \cup wv = uv$. A didrawing G Σ -*embeds* a didrawing H if there is a subdidrawing of G that is homeomorphic to a Σ -subdivision of H .

If G is a digraph, G^\natural denotes the undirected graph underlying G . If G is a drawing, G^* denotes a dual drawing. If G is a digraph drawn in the plane, let G^* be the digraph obtained from the dual graph $G^{\natural*}$ of G^\natural by directing its edges such that each edge of G^* crosses the corresponding edge of G from left to right (we recall that we fixed some orientation of Σ , so “left to right” makes sense). We call G^* a *right-dual didrawing*, and a *left-dual didrawing* is defined similarly. (Taking right-dual and then left-dual returns us to the original didrawing.) A *dual* drawing is either a right- or left-dual drawing. We say a didrawing G has *interleaving* at most k (where k is an even integer) if for each vertex v , the edges incident with v can be partitioned into at most k intervals in their cyclic order around v , and the edges in each interval either all enter v or all leave v .

3 The decomposition

Next, we want to explain the decomposition more precisely. The *leaves* of a tree are its vertices of degree one, and a tree T is *ternary* if each of its vertices has degree one or three.

Let S be a finite set with $|S| \geq 2$, let T be a ternary tree, and let τ be a bijection from S onto the set of leaves of T . We call (T, τ) a *carving* of S . For each edge e of T let $T_1(e), T_2(e)$ be the two components of $T \setminus e$. We call

$$(\{v \in S : \tau(v) \in V(T_1(e))\}, \{v \in S : \tau(v) \in V(T_2(e))\})$$

and its reverse the *partitions of S corresponding to e under (T, τ)* .

For instance, if G is a graph, we can ask for a carving (T, τ) of $E(G)$ such that for each edge $e \in E(T)$, there are at most k vertices incident with edges in both parts of the partition corresponding to e ; the smallest k for which such a carving exists is called the *branch-width* of G , and was introduced in [8]. It is within a constant factor of tree-width. In [9], Seymour and Thomas showed that one can compute the branch-width of a planar graph in polynomial time, but it is NP-hard for general graphs.

Similarly, if G is a graph, we can ask for a carving (T, τ) of $V(G)$ such that for each edge $e \in E(T)$, there are at most k edges between the two parts of the partition corresponding to e ; the smallest k for which such a pair (T, τ) exists is the *carving-width*, and again it was shown in [9] that one can compute the carving-width of a planar graph in polynomial time, and it is NP-hard for general graphs. The same paper also showed that if G is 2-connected, there is an “optimal” carving (T, τ) such that for each edge $e \in E(T)$, the set of edges joining the two parts of the corresponding partition is a *bond* of G , that is, a minimal edge-cutset of G , and so corresponds to a cycle of the dual graph.

We want something very close to this for didrawings G . We will look for carvings (T, τ) of $V(G)$ where the corresponding edge-cutsets are bonds of G (that is, of G^\natural), but we will measure the “size” of edge-cutsets in a different way.

If D is a cycle of a digraph G , some of its vertices might be the head of both incident edges, or the tail of both these edges. We call such a vertex a *change-vertex* of D , and the number of change-vertices is the *change number* of D . Thus, a cycle D has change number zero if and only if it is a directed cycle; and the change number of every cycle is even. Similarly, if P is a path of a digraph G , its *change number* is the number of vertices v of P such that v is the head of two edges of P or the tail of two edges of P . If G is a didrawing, the change number of a bond of G is the change number of the corresponding cycle in a dual drawing.

For $k \geq 1$, we say a digraph G is *weakly k -connected* or *k -weak* if G^\natural is k -connected. Let us say a didrawing has *diwidth* at most k if there is a carving (T, τ) of $V(G)$, such that for each edge e of T , the set of edges joining the two parts of the corresponding partition is a bond of G with change number at most k . (Didrawings that are not 2-weak do not admit any such “bond” carvings, so diwidth is only defined for 2-weak didrawings.) Then our first theorem says:

3.1 *For all even $k \geq 2$, and every odd integer $\lambda \geq 1$, every 2-weak loopless didrawing with interleaving at most 2λ and with diwidth more than $40\lambda^2 k^2$ embeds the $k \times k$ diwall.*

We remark that our interest in diwidth grew from a paper by Berger, Carter and Seymour [1], where it is shown that a strongly-connected, 2-weak loopless didrawing is “weightable” if and only if it has diwidth at most two. “Weightable” means that one can assign weights to its edges totalling to 1 in every directed cycle. (There are several equivalent conditions, for instance that a 0/1-weighting exists, or that for every triple of vertices, every directed cycle containing them uses them in the same cyclic order.)

Now let us turn to the decomposition used for our main theorem, when there is no bound on the interleaving. Let G be a loopless didrawing. A *dart* of G is a pair (v, e) where $v \in V(G)$, $e \in E(G)$ and e is incident with v . Let F be a simple closed curve in Σ , such that F has at most one point in common with each edge, and for each region r , either $F \cap r = \emptyset$ or $F \cap r$ is the interior of a line. We call such a curve F *good*. Let F be a good curve, and let Δ_1, Δ_2 be the two closed discs in Σ bounded by F . We say a dart (v, e) *belongs to* Δ_i if either v belongs to the interior of Δ_i , or $v \in F$ and $e \setminus \{v\}$ is a subset of the interior of Δ_i . For $i = 1, 2$, let D_i be the set of darts that belong to Δ_i ; we call (D_1, D_2) and (D_2, D_1) the *dart-partitions from* F . A partition (A, B) of the set of all darts is *sensible* if (A, B) is a dart-partition from F for some good curve F .

Let F be a good curve, and let $\Delta_1, \Delta_2, D_1, D_2$ be as before. A *change-region* for F is a region r incident with two edges e, f that both cross F (since F is good, there cannot be three edges crossing F and incident with r), such that Δ_1 contains the head of one of e, f and Δ_2 contains the head of the other. (If r, e, f are as above, it follows that $F \cap (r \cup e \cup f)$ is a line with ends in the interiors of e and f .) The number of change-regions for F is its *change number*. If $F \cap V(G) = \emptyset$, this coincides with our previous definition of change number for bonds. Different good curves F may give rise to the same dart-partition (D_1, D_2) , but it follows that all such curves have the same change number; the number of change-regions for F is determined by the partition (D_1, D_2) . Let us say the *cost* of (D_1, D_2) is the sum of the change number of F and $|F \cap V(G)|$ (again, it is easy to check that the cost of F is also determined by the dart-partition).

We say a loopless didrawing has *dart-width* at most k if there is a carving (T, τ) of the set of all darts, such that for each edge e of T , the corresponding dart-partition is sensible and has cost at most k . A digraph is *weakly 2-edge-connected* if G^\dagger is 2-edge-connected (no didrawing with a cut-edge admits a sensible dart-partition). We will prove:

3.2 *For all even $k \geq 2$, if G is a loopless, weakly 2-edge-connected didrawing that does not embed the $k \times k$ diwall, then G has dart-width at most $3240k^2$.*

If we apply 3.2 to the digraph obtained from an undirected planar graph by replacing each edge by two opposite directed edges, we obtain 1.1 restricted to planar graphs and re-expressed in terms of branch-width.

4 Patterns of directed cuts

Much of the paper is concerned with proving 3.1, and we will begin on that now. A *pattern* is a sequence (π_1, \dots, π_k) where $k \geq 1$ and $\pi_i \in \{+1, -1\}$ for $1 \leq i \leq k$. Let u, v be distinct vertices of a 1-weak digraph G . Let us say a (u, v) -*multicut* is a sequence (A_0, \dots, A_k) of subsets of $V(G)$ with $k \geq 1$, such that

- A_0, \dots, A_k are pairwise disjoint and have union $V(G)$;
- $u \in A_0$ and $v \in A_k$;
- for $0 \leq i, j \leq k$, if $j \geq i + 2$ there are no edges between A_i and A_j ;
- for $1 \leq i \leq k$, either there are no edges from A_{i-1} to A_i or there are no edges from A_i to A_{i-1} .

(See Figure 3.)

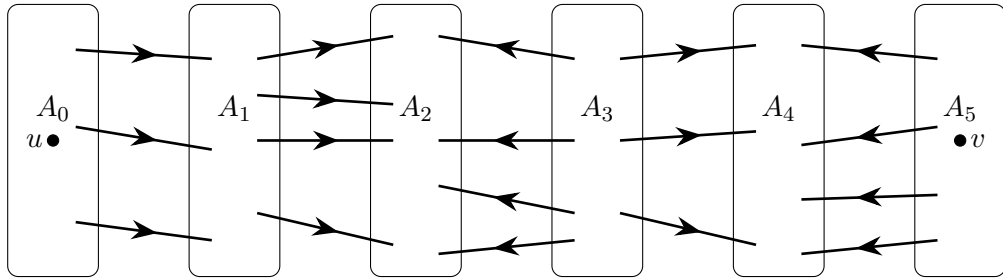


Figure 3: A (u, v) -multicut with pattern $(1, 1, -1, 1, -1)$.

If (A_0, \dots, A_k) is a (u, v) -multicut, let π_1, \dots, π_k be defined by $\pi_i = 1$ if there are no edges from A_i to A_{i-1} , and $\pi_i = -1$ if there are no edges from A_{i-1} to A_i for $1 \leq i \leq k$. (This is well-defined, since there is some edge between A_{i-1}, A_i because G is 1-weak.) We call the sequence $\pi = (\pi_1, \dots, \pi_k)$ the *pattern* of the multicut (A_0, \dots, A_k) .

We are interested in when there is a (u, v) -multicut with a given pattern. The main result of this section is:

4.1 Let u, v be distinct vertices of a 1-weak digraph G , and let $\pi = (\pi_1, \dots, \pi_k)$ be a pattern. Then there is a (u, v) -multicut in G with pattern π if and only if for every path of G between u, v , there is a (u, v) -multicut in P with pattern π .

Proof. The “only if” part is clear. For the “if” part, we proceed by induction on k , and for fixed k , by induction on $|E(G)|$. We assume that there is no (u, v) -multicut in G with pattern π , and we must show that there is a path P between u, v that admits no (u, v) -multicut in P with pattern π . If $k = 1$ the claim is clear, so we assume that $k \geq 2$. We may assume that there are no loops (by deleting them, since loops make no difference). By reversing all edges of G and negating π if necessary, we may assume that $\pi_k = +1$.

(1) If there is an edge of G with tail v , then there is a path P between u, v that admits no (u, v) -multicut in P with pattern π .

Suppose that $e = vw \in E(G)$. If $w = u$, then e makes a path P between u, v that admits no (u, v) -multicut in P with pattern π as required. So we assume that $w \neq u, v$. Let G' be obtained from G by contracting e , and let v' be the vertex of G' made by identifying v, w under contraction. Since there is no (u, v) -multicut in G with pattern π , it follows that there is no (u, v') -multicut in G' with pattern π . From the inductive hypothesis, there is a path P' of G' between u, v' that admits no (u, v') -multicut with pattern π . Hence there is a path P of G between u, v , such that either $E(P) = E(P')$ or $E(P) = E(P') \cup \{e\}$. We claim that in both cases, P satisfies (1). This is clear if $E(P) = E(P')$, so we assume that $E(P) = E(P') \cup \{e\}$, and hence e is the edge of P incident with v , and $P \setminus v$ is a path of G between u, w . Suppose that P admits a (u, v) -multicut (A_0, \dots, A_k) with pattern π . Since $\pi_k = +1$ and e has tail v , it follows that $w \notin A_{k-1}$; and $w \notin A_i$ for $i \leq k-2$ from the definition of a (u, v) -multicut. Thus $w \in A_k$, and so contracting e yields a (u, v') -multicut in P' with pattern π , a contradiction. This proves that P satisfies (1), and so proves (1).

From (1) we may assume that every edge of G incident with v has head v . Let N be the set of all vertices adjacent to v , and let G' be obtained from G by contracting all edges between N, v . Let v' be the vertex of G' made by identifying $N \cup \{v\}$ under contraction. Let π' be the pattern $(\pi_1, \dots, \pi_{k-1})$.

(2) There is no (u, v') -multicut in G' with pattern π' .

Suppose that (A'_0, \dots, A'_{k-1}) is a (u, v') -multicut in G' with pattern π' . Thus $v' \in A'_{k-1}$. Let $A_i = A'_i$ for $0 \leq i \leq k-2$, let $A_{k-1} = (A'_{k-1} \setminus \{v'\}) \cup N$, and let $A_k = \{v\}$. We claim that (A_0, \dots, A_k) is a (u, v) -multicut in G with pattern π . To show this, we must check that

- A_0, \dots, A_k are pairwise disjoint and have union $V(G)$, and $u \in A_0$ and $v \in A_k$ (and this is all clear);
- there are no edges between A_i, A_{k-1} with $i \leq k-3$ (this is true since there are no edges of G' between A'_i, A'_{k-1});
- there are no edges between A_i, A_k with $i \leq k-2$ (this is true since all edges of G incident with v are between v, N);
- every edge between A_{k-1}, A_k is from A_{k-1} to A_k (this is true since no edge has tail v).

This shows that (A_0, \dots, A_k) is a (u, v) -multicut in G with pattern π , a contradiction, and so proves (2).

From the inductive hypothesis on k , there is a path P' of G' between u, v' that admits no (u, v') -multicut with pattern π' . Hence there is a path Q of G between $u, N \cup \{v\}$ with $E(Q) = E(P')$; and since every vertex in N is an in-neighbour of v , there is a path P of G between u, v such that either $E(P) = E(P')$, or $E(P) = E(P') \cup \{e\}$ for some edge $e = qv$, where $q \in N$ is the end of Q in N . If $E(P) = E(P')$ then P admits no (u, v) -multicut with pattern π' , and hence none with pattern π , as required. Thus we assume that $E(P) = E(P') \cup \{e\}$ from some edge $e = qv$. Suppose that (A_0, \dots, A_k) is a (u, v) -multicut in P with pattern π . Hence $(A_0, \dots, A_{k-2}, A_{k-1} \cup A_k)$ is a (u, v) -multicut in P with pattern π' . The edge e has an end in A_k , and so $q \notin A_0, \dots, A_{k-2}$. Consequently, both its ends are in $A_{k-1} \cup A_k$; and therefore contracting e converts $(A_0, \dots, A_{k-2}, A_{k-1} \cup A_k)$ to a (u, v') -multicut in P' with pattern π' , a contradiction. This proves that there is no (u, v) -multicut in P with pattern π , and therefore completes the proof of 4.1. \blacksquare

Here is a companion lemma:

4.2 *Let u, v be distinct vertices of a 1-weak digraph G , and suppose there is a (u, v) -multicut in G with some pattern $\pi = (\pi_1, \dots, \pi_k)$. Then there is a (u, v) -multicut (A_0, \dots, A_k) in G with pattern π , such that for $1 \leq i \leq k$, $G[A_0 \cup A_1 \cup \dots \cup A_{i-1}]$ and $G[A_i \cup \dots \cup A_k]$ are both connected.*

Proof. Choose a (u, v) -multicut (A_0, \dots, A_k) in G with pattern π , such that the sum (over $1 \leq i \leq k$) of the number of edges between A_{i-1}, A_i is minimum. Suppose that for some $i \in \{1, \dots, k\}$, $G[A_0 \cup A_1 \cup \dots \cup A_{i-1}]$ is not connected, and let X, Y be a partition of $A_0 \cup A_1 \cup \dots \cup A_{i-1}$ into two nonempty subsets such that there are no edges between X, Y , and $u \in X$. Then

$$(A_0 \cap X, A_1 \cap X, \dots, A_{i-1} \cap X, A_i \cup Y, A_{i+1}, \dots, A_k)$$

is also a (u, v) -multicut in G with pattern π ; so, from the minimality in the choice of (A_0, \dots, A_k) , there is no edge between Y, A_i , contradicting that G is 1-weak. Thus $G[A_0 \cup A_1 \cup \dots \cup A_{i-1}]$ is connected, and similarly so is $G[A_i \cup \dots \cup A_k]$, for $1 \leq i \leq k$. This proves 4.2. \blacksquare

Thus, 4.1 reduces the general question to the same question when G is a path. Let us look more closely at the path case. We are mostly interested in when the desired sequence (π_1, \dots, π_k) is *alternating*, that is, π_i, π_{i+1} have opposite sign for $1 \leq i \leq k-1$.

Let P be a path between u, v in a digraph G , and let $u = q_0, q_1, \dots, q_t = v$ be vertices of P in order (not necessarily all distinct). For $1 \leq i \leq t$ let P_i be the subpath between q_{i-1}, q_i , and suppose that each P_i is a directed path. We say that (P_1, \dots, P_t) is a *concatenation* for (P, u, v) . Now let $\pi = (\pi_1, \dots, \pi_t)$ be an alternating pattern; we say that the concatenation (P_1, \dots, P_t) is a π -*concatenation* for (P, u, v) if for $1 \leq i \leq t$, P_i is a directed path from q_{i-1} to q_i if and only if $\pi_i = 1$. (A concatenation may be a π -concatenation for more than one pattern π since we permit some of the paths P_i to have length zero.) If $\pi = (\pi_1, \dots, \pi_t)$ is a pattern, $-\pi$ denotes the pattern $(-\pi_1, \dots, -\pi_t)$.

We claim that:

4.3 *Let P be a path between u, v and let $\pi = (\pi_1, \dots, \pi_k)$ be an alternating pattern. Then exactly one of the following holds:*

- P admits a (u, v) -multicut with pattern π ;
- there is a $(-\pi)$ -concatenation for (P, u, v) .

We omit the proof, which is clear.

If $\pi = (\pi_1, \dots, \pi_k)$ is an alternating pattern of length k , let π^{++} denote the alternating pattern obtained by adding a new first and last term, that is,

$$(-\pi_1, \pi_1, \dots, \pi_k, -\pi_k).$$

Note that $-\pi$ is an initial subsequence of π^{++} .

We also need the following:

4.4 *Let G be a 1-weak digraph, let $u, v \in V(G)$ be distinct, let $\pi = (\pi_1, \dots, \pi_k)$ be an alternating pattern, and suppose that (A_0, \dots, A_k) is a (u, v) -multicut with pattern π in G . Let P be a path between u, v such that there is a π^{++} -concatenation for (P, u, v) . Then P contains exactly one edge between A_{i-1}, A_i for $1 \leq i \leq k$.*

Proof. For $1 \leq i \leq k$, there is an edge of P between A_{i-1}, A_i ; let e_i, f_i be the first and last such edges in P . An edge $e = ab$ of P is *forward* in (P, u, v) if u, a belong to the same component of $P \setminus e$, and *backward* otherwise.

(1) for $1 \leq i \leq k$, e_i, f_i are forward edges of P if and only if $\pi_i = 1$.

The component of $P \setminus e_i$ containing u is a path with no edges between A_{i-1}, A_i , and since it has one end u , all its vertices belong to $A_0 \cup \dots \cup A_{i-1}$. Hence e_i is a forward edge of P if and only if all edges of G between A_{i-1}, A_i are from A_{i-1} to A_i , that is, if and only if $\pi_i = 1$. This proves the claim for e_i , and the claim for f_i is proved similarly, using the component of $P \setminus f_i$ containing v . This proves (1),

(2) $e_i = f_i$ for $1 \leq i \leq k$.

Define $\pi_0 = -\pi_1$, and $\pi_{k+1} = -\pi_k$. Thus $\pi^{++} = (\pi_0, \dots, \pi_{k+1})$. Now e_1, \dots, e_k are all distinct and in order in P . Let (P_0, \dots, P_{k+1}) be a π^{++} -concatenation for (P, u, v) . It follows from (1) that $e_1 \notin E(P_0)$, and therefore $e_1, \dots, e_k \in E(P_1), \dots, E(P_{k+1})$. Similarly $e_2 \notin E(P_1)$, so $e_2, \dots, e_k \in E(P_2), \dots, E(P_{k+1})$; and in general, $e_i \in E(P_i), \dots, E(P_{k+1})$. Similarly, $f_i \in E(P_0), \dots, E(P_i)$; and since e_i equals or is earlier than f_i in P , it follows that $e_i, f_i \in E(P_i)$. Suppose that $e_i \neq f_i$ for some i . If $\pi_i = 1$, then there is a directed path of P_i from the head of e_i to the tail of f_i , which is impossible since the head of e_i is in A_i and the tail of f_i is in A_{i-1} , and all edges of G between A_{i-1}, A_i are from A_{i-1} to A_i ; and the other case (when $\pi_i = -1$) is similar and we omit it. This proves (2).

From (2), only one edge of P is between A_{i-1}, A_i for $1 \leq i \leq k$. This proves 4.4. ▀

We deduce:

4.5 *Let G be a 1-weak digraph, let $u, v \in V(G)$ be distinct, let $\pi = (\pi_1, \dots, \pi_k)$ be an alternating pattern, and suppose that (A_0, \dots, A_k) is a (u, v) -multicut with pattern π in G . Let $f \in E(G)$, and*

suppose that there is no (u, v) -multicut with pattern π in G/f (G/f denotes the digraph obtained by contracting f). Then for some i , f is between A_{i-1} and A_i . Moreover, there is a path P between u, v and a π^{++} -concatenation for (P, u, v) , such that f is the only edge of P between A_{i-1}, A_i .

Proof. By 4.1, there is a path P' in G/f between u, v such that (P', u, v) admits a $(-\pi)$ -concatenation (P_1, \dots, P_k) . This no longer exists in G , so there is a path P of G between u, v with $f \in E(P)$, such that P' is obtained from P by contracting f . Hence (P, u, v) admits a π^{++} -concatenation; and hence the claim follows from 4.4. This proves 4.5. ▀

5 Box systems

For $k_1, k_2 \geq 2$, even, let us say a $k_1 \times k_2$ *diwall layout* in a digraph G consists of two sequences P_1, \dots, P_{k_1} and Q_1, \dots, Q_{k_2} of paths of G , with the following properties:

- P_1, \dots, P_{k_1} are vertex-disjoint directed paths of G , and so are Q_1, \dots, Q_{k_2} ;
- for all i, j with $1 \leq i \leq k_1$ and $1 \leq j \leq k_2$, $P_i \cap Q_j$ is a path $R_{i,j}$ say;
- if i is odd, the paths $R_{i,1}, R_{i,2}, \dots, R_{i,k_2}$ are in order in P_i , and if i is even they are in reverse order;
- if j is even, the paths $R_{1,j}, R_{2,j}, \dots, R_{k_1,j}$ are in order in Q_j , and if j is odd they are in reverse order.

Thus, if P_1, \dots, P_k and Q_1, \dots, Q_k form a $k \times k$ diwall layout then their union embeds a $k \times k$ diwall; and conversely, if G embeds a $k \times k$ diwall then it contains a $k \times k$ diwall layout.

We need the $k_1 \times k_2$ *alternating grid*, defined for $k_1, k_2 \geq 2$, even, and illustrated in Figure 4 (which also shows the “semi-grid”, that we do not need yet). It has k_1 horizontal and k_2 vertical paths; they alternate in direction, and the first vertex of the top horizontal path is the last vertex of a vertical path (this last condition is just to make it unique). We denote the $k_1 \times k_2$ alternating grid by Γ_{k_1, k_2} .

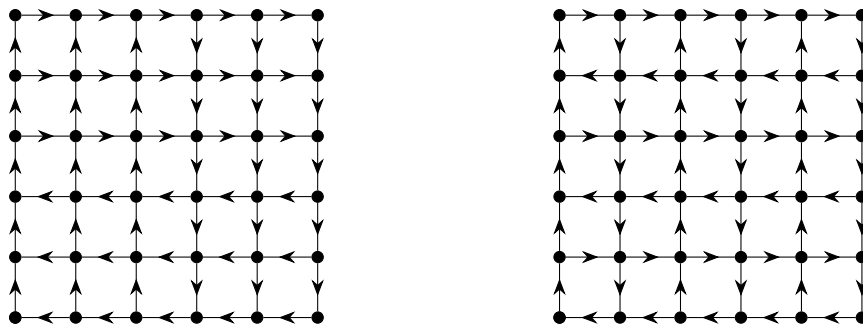


Figure 4: The 6×6 semi-grid and alternating grid.

Third, for $k_1, k_2 \geq 2$, even, let us say a $k_1 \times k_2$ *box system* in a digraph G is a map η from the set of vertices and edges of $H = \Gamma_{k_1, k_2}$ with the following properties:

- for each $v \in V(H)$, $\eta(v)$ is a subset of $V(G)$, and if $u, v \in V(H)$ are distinct then $\eta(u) \cap \eta(v) = \emptyset$;
- for each $e = uv \in E(H)$, $\eta(e) \in E(G)$, and the tail and head of $\eta(e)$ are in $\eta(u), \eta(v)$ respectively.
- for each $v \in V(H)$, if $e, f \in E(H)$ have head v and tail v respectively, then there is a directed path in $G[\eta(v)]$ from the head of $\eta(e)$ to the tail of $\eta(f)$.

Evidently a $k \times k$ box system is some kind of grid-like object in G . One might hope to extract a $k \times k$ diwall layout from it, by choosing for each horizontal path P_i of $\Gamma_{k,k}$ a directed path P'_i in G containing all the edges $\eta(e)$ ($e \in E(P_i)$), joining these edges via appropriate paths in $G[\eta(v)]$ for $v \in V(P_i)$, and defining Q'_j for each vertical path Q_j of $\Gamma_{k,k}$ similarly. But this does not work, because we cannot arrange that $P'_i \cap Q'_j$ is a path for each i, j . Nevertheless, we can extract a $k \times k$ diwall layout from a $k \times 3k$ box system, as we show next.

Figure 5 gives another picture of the $k \times k$ diwall, redrawn without diagonal edges. Each of the “vertical” paths now consists of the edges in two consecutive columns of the picture, together with edges of the horizontal paths joining them.

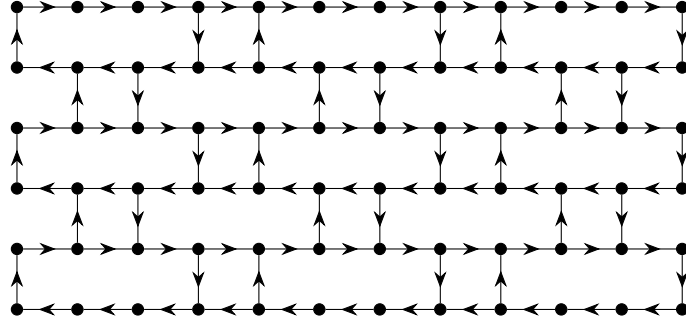


Figure 5: The 6×6 diwall, redrawn. If we subdivide the odd edges of the horizontal paths once, we obtain a subgraph of the $k \times 3k$ alternating grid.

5.1 *Let $k \geq 2$, even, and suppose that there is a $k \times 3k$ box system η in G ; then there is a $k \times k$ diwall layout in G , and hence G embeds a $k \times k$ diwall.*

Proof. Let P_1, \dots, P_k and Q_1, \dots, Q_{3k} be the horizontal and vertical paths of $\Gamma_{k,3k}$, numbered in order, and for $1 \leq i \leq k$ and $1 \leq j \leq 3k$ let v_{ij} be the vertex in both P_i, Q_j . For every pair (e, f) where $e, f \in E(\Gamma_{k,3k})$ and the head of e equals the tail of f , say v , let $L(e, f)$ be a directed path of G with first edge $\eta(e)$, last edge $\eta(f)$, and all internal vertices in $\eta(v)$, chosen subject to the following rule:

- if v is incident with two edges in a horizontal path of $\Gamma_{k,3k}$, say x, y , and $|\{e, f\} \cap \{x, y\}| = 1$, then $L(e, f)$ is chosen such that $L(e, f) \cap L(x, y)$ is a path.

(In other words, $L(e, f)$ consists of a subpath of $L(x, y)$ together with a minimal path between $L(x, y)$ and whichever of e, f is not in $\{x, y\}$.)

For $1 \leq i \leq k$, let the edges of P_i be e_1, \dots, e_{3k-1} in order; then

$$L(e_1, e_2) \cup L(e_2, e_3) \cup \dots \cup L(e_{3k-2}, e_{3k-1})$$

is a directed path of G that we call P'_i , for $1 \leq i \leq k$.

For $i \cong 4 \pmod{6}$, there is a directed path (R_i , say) of $\Gamma_{k,3k}$ with vertices

$$v_{1,i} - v_{2,i-1} - v_{2,i+1} - v_{2,i+2} - v_{3,i+2} - v_{3,i+1} - v_{3,i} - v_{4,i} - v_{4,i+1} - v_{4,i+2} - v_{5,i+2} - \cdots - v_{k,i}$$

in order, and for $i = 1 \pmod{6}$, there is a directed path R_i with vertices

$$v_{k,i} - v_{k-1,i} - v_{k-1,i+1} - v_{k-1,i+2} - v_{k-2,i+2} - v_{k-2,i+1} - v_{k-2,i} - v_{k-3,i} - \cdots - v_{1,i}.$$

For $1 \leq j \leq 3k$ with $j \cong 1, 4 \pmod{6}$, let the edges of R_j be e_1, \dots, e_{3k-5} in order, and let R'_j be the path

$$L(e_1, e_2) \cup L(e_2, e_3) \cup \cdots \cup L(e_{3k-6}, e_{3k-5}).$$

Then P'_1, \dots, P'_k and $R'_1, R'_4, R'_7, \dots, R'_{3k-2}$ form a $k \times k$ diwall layout. (The drawing in Figure 5 might make this easier to see.) This proves 5.1. ■

6 Paths across a disc

We need several results about paths across a didrawing in a disc, so let us first set up some notation. If G is a didrawing, two cycles C_1, C_2 of G are *non-crossing* if C_i bounds two closed discs Δ_i, Δ'_i in Σ for $i = 1, 2$, and one of Δ_1, Δ'_1 includes one of Δ_2, Δ'_2 . Let G be a didrawing, and let $v \in V(G)$, such that no loop is incident with v . By an *alternating ring* through v , we mean a sequence (C_1, \dots, C_k) of directed cycles of G , each containing v , pairwise edge-disjoint and pairwise non-crossing, where $k \geq 2$ is even, such that, if e_i, f_i are the edges of C_i with tail v and head v respectively, then

$$f_1, e_2, f_3, e_4, \dots, e_k, f_k, e_{k-1}, \dots, f_2, e_1$$

are in anticlockwise cyclic order around v . We call k its *size*.

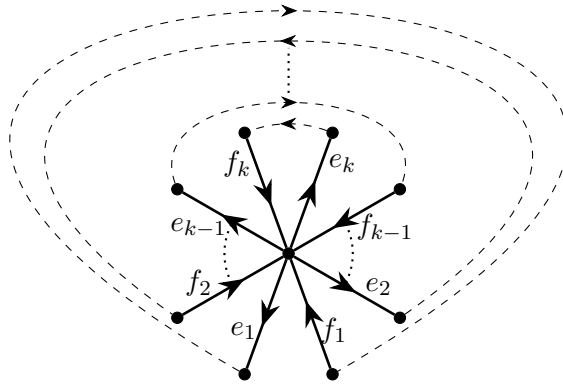


Figure 6: An alternating ring. There may be more edges incident with v that are not part of the ring.

Let G be a 1-weak didrawing, and let $v \in V(G)$, such that no loop is incident with v , and such that $G \setminus v$ is 1-weak. Let G^* be a dual didrawing, and for each vertex, edge or region x of G , let $\phi(x)$

be the corresponding region, edge or vertex of G^* . Consequently, the edges of G incident with v are mapped by ϕ to the edges of a cycle C of G^* that bounds the region $\phi(v)$ of G^* . Let R_1, R_2, R_3, R_4 be four subpaths of C , pairwise edge-disjoint and with $E(R_1 \cup R_2 \cup R_3 \cup R_4) = E(C)$, numbered in anticlockwise circular order in C around $\phi(v)$. An alternating ring (C_1, \dots, C_k) in G through v is from R_2 to R_4 if, with e_i, f_i ($1 \leq i \leq k$) defined as before,

- if i is odd then $\phi(e_i) \in E(R_4)$ and $\phi(f_i) \in E(R_2)$;
- if i is even then $\phi(e_i) \in E(R_2)$ and $\phi(f_i) \in E(R_4)$.

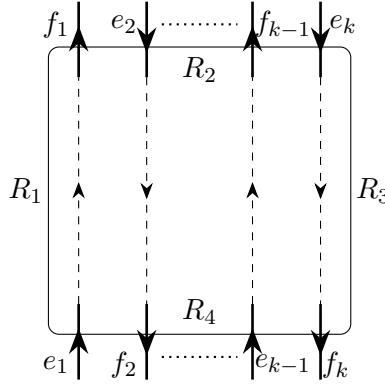


Figure 7: The end-edges of the members of an alternating ring from R_2 to R_4 . The labelled edges are now in clockwise order, since they are all incident with v , which is outside the cycle C and not drawn. The dashed paths are edge-disjoint and non-crossing.

6.1 Let G be a 1-weak didrawing, and let $v \in V(G)$, such that no loop is incident with v , and such that $G \setminus v$ is 1-weak. With notation as above, let $k \geq 2$ be an even integer. Then either:

- there is a path P of G^* between $V(R_1)$ and $V(R_3)$ with change number less than k ; or
- there is an alternating ring (C_1, \dots, C_k) in G through v from R_2 to R_4 .

Proof. Let $\pi = (\pi_1, \dots, \pi_k)$ be the alternating pattern with k terms, with $\pi_1 = -1$ if G^* is a right-dual didrawing of G , and $\pi_1 = 1$ if G^* is a left-dual didrawing of G . We assume the statement of the first bullet is false. Consequently, in the digraph H obtained from G^* by contracting the edges of R_1 (making a vertex r_1) and contracting the edges of R_3 (making a vertex r_3), there is no path P between r_1, r_3 such that (P, r_1, r_3) admits a $(-\pi)$ -concatenation. By 4.1 and 4.2, there is an (r_1, r_3) -multicut (A_0, \dots, A_k) in H with pattern π , such that for $1 \leq i \leq k$, $H[A_0 \cup A_1 \cup \dots \cup A_{i-1}]$ and $H[A_i \cup \dots \cup A_k]$ are both connected. Hence, the set of edges of H between A_{i-1} and A_i is a bond of H , and so its preimage under ϕ is the set of edges of a directed cycle C_i of H^* (and hence of G) containing v . These cycles are pairwise edge-disjoint, and pairwise non-crossing (since each is the boundary of a set of regions of H , and these sets are nested by inclusion) and so (C_1, \dots, C_k) is an alternating ring in G . This proves 6.1. \blacksquare

An alternating ring (C_1, \dots, C_k) through v is *disjointed* if C_1, \dots, C_k are pairwise vertex-disjoint except for v . We deduce:

6.2 *Let G be a 1-weak didrawing, and let $v \in V(G)$, such that no loop is incident with v , and such that $G \setminus v$ is 1-weak. With notation as before, let $k \geq 2$ be an even integer, and suppose that G has interleaving at most 2λ , where λ is an odd integer. Then either:*

- *there is a path P of G^* between $V(R_1)$ and $V(R_3)$ with change number less than λk ; or*
- *there is a disjointed alternating ring (C_1, \dots, C_k) in G through v from R_2 to R_4 .*

Proof. We assume the first bullet is false. From 6.1, there is an alternating ring $(B_1, \dots, B_{k\lambda})$ through v and from R_2 to R_4 . We claim that for $1 \leq i < j \leq \lambda k$, if $j - i \geq \lambda$ then $V(B_i \cap B_j) = \{v\}$; because suppose that some vertex $u \neq v$ belongs to $V(B_i \cap B_j)$. Then u also belongs to $B_{i'}$ for all i' with $i \leq i' \leq j$, and so G has interleaving at least $2(j - i + 1)$, a contradiction. Let $C_i = B_{(i-1)\lambda+1}$ for $1 \leq i \leq k$; then C_1, \dots, C_k are pairwise vertex-disjoint except for v , and since λ is odd, it follows that (C_1, \dots, C_k) is an alternating ring through v from R_2 to R_4 . This proves 6.2. ■

We deduce:

6.3 *Let G be a 1-weak didrawing, and let $v \in V(G)$, such that no loop is incident with v , and such that $G \setminus v$ is 1-weak. With notation as before, suppose also that the didrawing G has interleaving at most 2λ where λ is odd. Let $k_1, k_2 \geq 2$ be even integers, and let $n = 2\lambda k_1 k_2 + 3k_2 - 2\lambda k_1 - 2$. Suppose that every path of G^* between $V(R_1)$ and $V(R_3)$ has change number at least λk_1 , and there is a disjointed alternating ring (B_1, \dots, B_n) through v from R_3 to R_1 . Then there is a disjointed alternating ring (C_1, \dots, C_{k_1}) through v from R_2 to R_4 , and a disjointed alternating ring (D_1, \dots, D_{k_2}) through v from R_3 to R_1 , such that for $1 \leq h \leq k_1$ and $1 \leq i < j \leq k_2$, there is a subpath of C_h with one end v that contains all vertices of $C_h \cap D_i$ and no vertices of $C_h \cap D_j$ except v .*

Proof. We proceed by induction on $|E(G)|$. By 6.1 there is an alternating ring of size λk_1 through v and from R_2 to R_4 , say $(F_1, \dots, F_{\lambda k_1})$. (See Figure 8.) We may assume:

- (1) *Let $1 \leq h \leq \lambda k_1$, and let $1 \leq j, j' \leq n$ with $j' - j > 2(\lambda k_1 + 1)$. Then there is a subpath of F_h with one end v that contains all vertices of $F_h \cap B_j$ and no vertices of $F_h \cap B_{j'}$ except v .*

Since (B_1, \dots, B_n) is a disjointed alternating ring, and the path $F_h \setminus v$ contains a vertex of $B_{j+\lambda k_1+1}$ and a vertex of $B_{j+\lambda k_1+2}$, there is a subpath of $F_h \setminus v$ between $V(B_{j+\lambda k_1+1})$ and $V(B_{j+\lambda k_1+2})$ with no internal vertex in either of these sets. Let e be an edge of this subpath. Thus e belongs to none of B_1, \dots, B_n . Since $k_1, k_2 \geq 2$, it follows that e belongs to a cycle of $G \setminus v$ (in the union of two B_i 's and two F_h 's); and since G and $G \setminus v$ are 1-weak, it follows that $G \setminus e$ and $(G \setminus e) \setminus v$ are 1-weak. From the inductive hypothesis on $|E(G)|$, we may assume that there is no alternating ring in $G \setminus e$ of size λk_1 through v and from R_2 to R_4 . Hence, from 4.5, there is a path P of G^* between $V(R_1)$ and $V(R_3)$ with change number at most λk_1 such that e is the only edge in $E(F_h)$ mapped into $E(P)$ by ϕ . Consequently, there is a partition (A, B) of $V(G) \setminus \{v\}$, such that the edges of $G \setminus \{v\}$ between A, B are precisely the edges of $G \setminus \{v\}$ mapped by ϕ into $E(P)$. In particular, e is the only edge of F_h between A, B . Hence $F_h[A \cup \{v\}]$ is a path Q_1 say, and $F_h[B \cup \{v\}]$ is a path Q_2 . From the

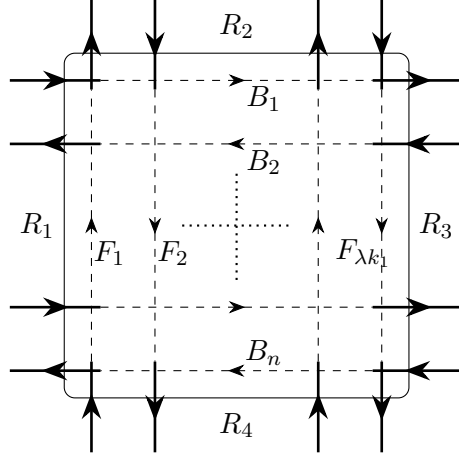


Figure 8: The B_i 's are vertex-disjoint, and the F_j 's are edge-disjoint and non-crossing. The paths are in the orders shown, but the intersections might be much more complicated than indicated.

choice of e , each of $B_1, \dots, B_{j+\lambda k_1+1}$ has a vertex in $V(Q_1) \setminus \{v\}$, and each of $B_{j+\lambda k_1+2}, \dots, B_n$ has a vertex in $V(Q_2) \setminus \{v\}$.

Let I be the set of all $i \in \{1, \dots, n\}$ such that some edge of B_i is between A, B . It follows from planarity that I is an interval (that is, if $i_1 \leq i_2 \leq i_3$ and $i_1, i_3 \in I$ then $i_2 \in I$). Since P has change number at most λk_1 , it follows that $|I| \leq \lambda k_1 + 1$. If $j \in I$ then, since the path P intersects each of $B_j, \dots, B_{j+\lambda k_1+1}$, it follows that $j, j+1, \dots, j+\lambda k_1+1 \in I$, contradicting that $|I| \leq \lambda k_1 + 1$. Thus $j \notin I$. Similarly, if $j' \in I$, then $j', j'-1, \dots, j+\lambda k_1+2 \in I$, contradicting that $|I| \leq \lambda k_1 + 1$ since $j' - (j + \lambda k_1 + 2) + 1 > \lambda k_1 + 1$; and hence $j' \notin I$. Since no edge of B_j is between A, B , and B_j has a vertex in $V(Q_1) \setminus \{v\} \subseteq A$, it follows that $V(B_j) \subseteq A \cup \{v\}$, and in particular all vertices of $F_h \cap B_j$ belong to Q_1 . Similarly, all vertices of $F_h \cap B_{j'}$ belong to Q_2 and therefore do not belong to Q_1 . This proves (1).

Let $\mu = 2(\lambda k_1 + 1) + 1$, and for $1 \leq i \leq k_2$ let $D_i = B_{(i-1)\mu+1}$ (this is well-defined, from the choice of n .) It follows that for $1 \leq h \leq \lambda k_1$ and $1 \leq i < j \leq k_2$, there is a subpath P of F_h with one end v that contains all vertices of $F_h \cap D_i$ and no vertices of $F_h \cap D_j$ except v . For $1 \leq i \leq k_1$ let $C_i = F_{(i-1)\lambda+1}$; then (C_1, \dots, C_{k_1}) and (D_1, \dots, D_{k_2}) are disjoint alternating rings. This proves 6.3. \blacksquare

6.4 Let G be a 1-weak didrawing, and let $v \in V(G)$, such that no loop is incident with v , and such that $G \setminus v$ is 1-weak. With notation as before, let $k \geq 2$ be even, and suppose that there is a disjoint alternating ring (C_1, \dots, C_k) through v from R_2 to R_4 , and a disjoint alternating ring (D_1, \dots, D_{3k}) through v from R_3 to R_1 , such that for $1 \leq h \leq k$ and $1 \leq i < j \leq 3k$, there is a subpath P of C_h with one end v that contains all vertices of $C_h \cap D_i$ and no vertices of $C_h \cap D_j$ except v . Then $G \setminus v$ embeds the $k \times k$ diwall.

Proof. This is immediate from 5.1. Alternatively, for all $i \in \{1, \dots, k-1\}$ and all $j \in \{1, \dots, 3k\}$, let $L_{i,j}$ be a minimal subpath of $D_j \setminus \{v\}$ between $V(C_i), V(C_{i+1})$. Then the paths $C_1 \setminus \{v\}, \dots, C_k \setminus \{v\}$ and the paths $L_{i,j}$ where $1 \leq i \leq k-1$ and $1 \leq j \leq 3k$ and either

- i is odd and j is congruent to 0 or 1 modulo 6, or
- i is even and j is congruent to 3 or 4 modulo 6

form a $k \times k$ diwall layout. This proves 6.4. ■

Putting these pieces together, we deduce:

6.5 *Let $k \geq 2$ be even. Let G be a 1-weak didrawing, and let $v \in V(G)$, such that no loop is incident with v , and such that $G \setminus v$ is 1-weak. Suppose that G has interleaving at most 2λ , where λ is an odd integer. Let G^* be a dual didrawing, and for each vertex, edge or region x of G , let $\phi(x)$ be the corresponding region, edge or vertex of G^* . Let C be the cycle of G^* with edge set all edges $\phi(e)$ where $e \in E(G)$ is incident with v in G . Let $n = 6\lambda k^2 - 2\lambda k + 9k - 2$, and suppose that:*

- R_1, R_2, R_3, R_4 are edge-disjoint subpaths of C , numbered in cyclic order in C ;
- every path of G^* between $V(R_1)$ and $V(R_3)$ has change number at least λk ; and
- every path of G^* between $V(R_2)$ and $V(R_4)$ has change number at least λn .

Then $G \setminus v$ embeds the $k \times k$ diwall.

Proof. By 6.2, there is a disjointed alternating ring (B_1, \dots, B_n) through v from R_3 to R_1 . By 6.3, there is a disjointed alternating ring (C_1, \dots, C_k) through v from R_2 to R_4 , and a disjointed alternating ring (D_1, \dots, D_{3k}) through v from R_3 to R_1 , such that for $1 \leq h \leq k$ and $1 \leq i < j \leq 3k$, there is a subpath of C_h with one end v that contains all vertices of $C_h \cap D_i$ and no vertices of $C_h \cap D_j$ except v . By 6.4, $G \setminus v$ embeds the $k \times k$ diwall. This proves 6.5. ■

7 Shortcuts across a cycle

7.1 *Let $k, \lambda \geq 1$ be integers. Let G be a 2-weak loopless didrawing, and let C be a cycle of G with change number at most n , where n is even and $n \geq 4k + 2\lambda + 10$. Suppose that G^* has interleaving at most 2λ , and let Δ be a disc bounded by C . Suppose that for every path P of $G[\Delta]$ with distinct ends in $V(C)$ and with no internal vertex or edge in C , there is a subpath Q of C between the ends of P such that $P \cup Q$ has change number more than n . Then either:*

- the interior of Δ is a region of G ; or
- there are four subpaths R_1, R_2, R_3, R_4 of G , pairwise edge-disjoint, with union C , and numbered in order in C , such that there is no path in $G[\Delta]$ between $V(R_1), V(R_3)$ with change number at most k , and there is no path in $G[\Delta]$ between $V(R_2), V(R_4)$ with change number at most k .

Proof. We assume the first bullet is false. For a path or cycle P we define $\text{ch}(P)$ to be the change number of P .

(1) $\text{ch}(C) \geq n - 2\lambda$.

Each edge of C is incident with a region of G included in Δ . If all these regions are equal, then the interior of Δ is a region of G (because G is 2-weak and loopless) and the first bullet holds, a contradiction. Thus they are not all equal; and hence there is a subpath P of C and a region $r \subseteq \Delta$ of G , such that some edge of C is incident with r , and every vertex or edge of C incident with r belongs to P . Choose such P, r with P minimal. It follows that every edge of P is incident with r (because if some edge of P is incident with a different region $r' \subseteq \Delta$, then all vertices and edges of C incident with r' belong to P , from planarity, and for the same reason, not both ends of P are incident with both r, r' ; this contradicts the minimality of P). The boundary of r is a cycle, since G is 2-weak and loopless, and hence consists of the union of P and a path Q of $G[\Delta]$ joining the ends of P with no internal vertex or edge in C .

Let $P' \neq P$ be the second path of C between the ends of P . Now $\text{ch}(P \cup Q)$ is at most $2\lambda \leq n$, from the interleaving hypothesis, and so $\text{ch}(P' \cup Q) > n$ (and hence at least $n + 2$). But $\text{ch}(P' \cup Q) \leq \text{ch}(P') + \text{ch}(Q) + 2$, and so $\text{ch}(C) \geq \text{ch}(P') \geq n - 2\lambda$. This proves (1).

From (1), there are four subpaths R_1, \dots, R_4 of C , pairwise edge-disjoint and with union C , numbered in cyclic order in C , and each with change number at least $(n - 2\lambda)/4 - 2 > k$. We may assume (for a contradiction) that R_1, \dots, R_4 do not satisfy the theorem; let P be a path in $G[\Delta]$, either between $V(R_2), V(R_4)$ or between $V(R_1), V(R_3)$, with

- change number as small as possible (and hence at most k);
- subject to that, with $C \cup P$ minimal; and
- subject to that, with length as small as possible.

We may assume that P is between $V(R_1), V(R_3)$ from the symmetry. From the minimality of the length of P , the only vertex of P in $V(R_1)$ is one of its ends, say r_1 , and its only vertex in $V(R_3)$ is its other end r_3 . Moreover, not both $V(R_2), V(R_4)$ contain internal vertices of P , again from the minimality of the length of P , so we may assume that every internal vertex of P in $V(C)$ belongs to R_2 . Let the components of $P \cap C$ be Q_0, \dots, Q_m , numbered in order in P , where $r_1 \in V(Q_0)$ and $r_3 \in V(Q_m)$. Thus, $m \geq 1$, and either Q_0 has only one vertex, or r_1 is the common end of R_1, R_2 , and the same for Q_m . For $1 \leq i \leq m$, let P_i be the minimal subpath of P between $V(Q_{i-1}), V(Q_i)$. Thus, $Q_0, P_1, Q_1, P_2, \dots, P_m, Q_m$ are edge-disjoint subpaths of P with union P , numbered in order in P , and P_1, \dots, P_m each have at least one edge. For $1 \leq i \leq m$, let S_i be the subpath of the path $R_1 \cup R_2 \cup R_3$ between the ends of P_i , and let S'_i be the other subpath of C between these two vertices. Thus, $Q_0, S_1, Q_1, \dots, S_m, Q_m$ are edge-disjoint subpaths of $R_1 \cup R_2 \cup R_3$ with union a path between R_1, R_3 , numbered in order in this path, and S_1, \dots, S_m each have at least one edge.

For $1 \leq i \leq m$, no internal vertex or edge of P_i belongs to C , and so from the hypothesis, one of $\text{ch}(P_i \cup S_i), \text{ch}(P_i \cup S'_i) > n$. Since $S_i \cup S'_i = C$ has change number at most n , and the change numbers of cycles are even, the change number of P_i is at least that of one of S_i, S'_i . But $\text{ch}(S'_i) \geq (n - 2\lambda)/4 - 2$ since S'_i includes R_4 , and since $\text{ch}(P_i) \leq k < (n - 2\lambda)/4 - 2$, it follows that $\text{ch}(P_i) \geq \text{ch}(S_i)$. (Consequently the latter is at most k , and so $R_2 \not\subseteq S_i$, and therefore $m \geq 2$.)

If v is a vertex of a path or cycle T , we say $\delta(v, T) = 1$ if v is a change-vertex of T , and $\delta(v, T) = 0$ otherwise. For $1 \leq i \leq m$, let the ends of P_i be a_i, b_i , where a_i is between r_1, b_i in P . (Thus $a_i \neq b_i$, but possibly $b_i = a_{i+1}$ if $i < m$.)

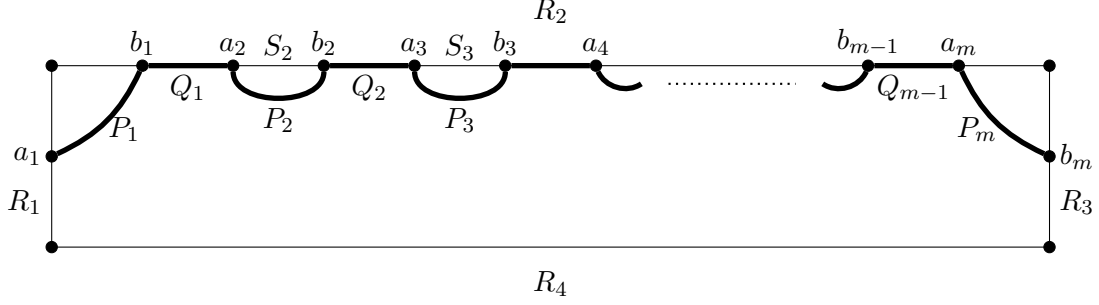


Figure 9: The path P (thick) and various derived paths. The picture might be misleading, because P_1 might have both ends in R_2 , and then Q_0 might have edges; and the same for P_m, Q_m .

(2) For $1 \leq i \leq m$,

$$\text{ch}(P_i) + \delta(a_i, P_i \cup S'_i) + \delta(b_i, P_i \cup S'_i) \geq \text{ch}(S_i) + \delta(a_i, C) + \delta(b_i, C) + 2.$$

As we have seen, $\text{ch}(P_i \cup S'_i) \geq \text{ch}(C) + 2$; but

$$\text{ch}(P_i \cup S'_i) = \text{ch}(P_i) + \delta(a_i, P_i \cup S'_i) + \delta(b_i, P_i \cup S'_i) + \text{ch}(S'_i)$$

and

$$\text{ch}(C) = \text{ch}(S_i) + \delta(a_i, C) + \delta(b_i, C) + \text{ch}(S'_i),$$

and the claim follows. This proves (2).

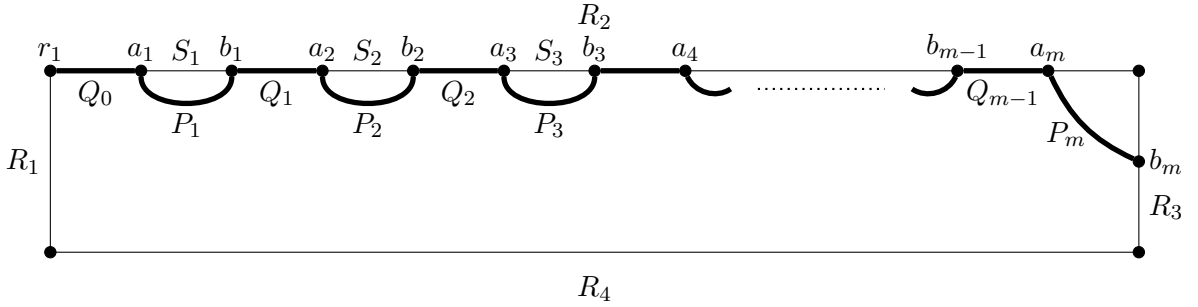


Figure 10: When Q_0 has an edge.

(3) Q_0 has no edges, and Q_m has no edges.

Suppose that Q_0 has an edge. (See Figure 10.) Then r_1 is the common end of R_1, R_2 , and Q_0 is a subpath of R_2 between r_1 and a_1 . Since $m \geq 2$, a_1 is an internal vertex of R_2 . From the minimality of $C \cup P$, the union of $Q_0 \cup S_1$ and the subpath of P between b_1 and r_3 has change number more (and hence at least two more, since $E(Q_0) \neq \emptyset$) than that of P ; and since the union of $Q_0 \cup P_1$ and the subpath of P between b_1 and r_3 equals P , we deduce that

$$\delta(a_1, C) + \text{ch}(S_1) + \delta(b_1, C) \geq \delta(a_1, P) + \text{ch}(P_1) + \delta(b_1, P) + 2.$$

But from (2),

$$\text{ch}(P_1) + \delta(a_1, P_1 \cup S'_1) + \delta(b_1, P_1 \cup S'_1) \geq \text{ch}(S_1) + \delta(a_1, C) + \delta(b_1, C) + 2.$$

Adding, we deduce:

$$\begin{aligned} & \delta(a_1, C) + \text{ch}(S_1) + \delta(b_1, C) + \text{ch}(P_1) + \delta(a_1, P_1 \cup S'_1) + \delta(b_1, P_1 \cup S'_1) \\ & \geq \delta(a_1, P) + \text{ch}(P_1) + \delta(b_1, P) + 2 + \text{ch}(S_1) + \delta(a_1, C) + \delta(b_1, C) + 2, \end{aligned}$$

which simplifies to

$$\delta(a_1, P_1 \cup S'_1) + \delta(b_1, P_1 \cup S'_1) \geq \delta(a_1, P) + \delta(b_1, P) + 4;$$

and this is impossible since $\delta(a_1, P_1 \cup S'_1), \delta(b_1, P_1 \cup S'_1) \leq 1$. Thus Q_0 has no edges and similarly Q_m has no edges. This proves (3).

Let $I = \{i \in \{1, \dots, m-1\} : b_i = a_{i+1}\}$ and $I' = \{1, \dots, m-1\} \setminus I$. It follows that

$$k \geq \text{ch}(P) \geq \sum_{1 \leq i \leq m} \text{ch}(P_i) + \sum_{i \in I'} \text{ch}(Q_i).$$

Moreover, the change number of $S_1 \cup Q_1 \cup S_2 \cup Q_2 \cup \dots \cup Q_{m-1} \cup S_m$ (which includes R_2) equals

$$\sum_{1 \leq i \leq m} \text{ch}(S_i) + \sum_{i \in I} \delta(a_i, C) + \sum_{i \in I'} (\delta(a_i, C) + \text{ch}(Q_i) + \delta(b_i, C)) \geq \text{ch}(R_2) \geq (n - 2\lambda)/4 - 2 > k,$$

and by summing these two inequalities we obtain

$$\sum_{1 \leq i \leq m} \text{ch}(S_i) + \sum_{i \in I} \delta(a_i, C) + \sum_{i \in I'} (\delta(a_i, C) + \delta(b_i, C)) > \sum_{1 \leq i \leq m} \text{ch}(P_i).$$

On the other hand, (2) implies that

$$\text{ch}(P_i) \geq \text{ch}(S_i) + \delta(a_i, C) + \delta(b_i, C)$$

for $1 \leq i \leq m$, and by summing for $1 \leq i \leq m$ we deduce:

$$\sum_{1 \leq i \leq m} \text{ch}(P_i) \geq \sum_{1 \leq i \leq m} \text{ch}(S_i) + \sum_{1 \leq i \leq m} (\delta(a_i, C) + \delta(b_i, C)).$$

Combining with the previous inequality, we deduce:

$$\begin{aligned} & \sum_{1 \leq i \leq m} \text{ch}(S_i) + \sum_{i \in I} \delta(a_i, C) + \sum_{i \in I'} (\delta(a_i, C) + \delta(b_i, C)) + \sum_{1 \leq i \leq m} \text{ch}(P_i) \\ & > \sum_{1 \leq i \leq m} \text{ch}(P_i) + \sum_{1 \leq i \leq m} \text{ch}(S_i) + \sum_{1 \leq i \leq m} (\delta(a_i, C) + \delta(b_i, C)), \end{aligned}$$

which simplifies to a contradiction.

This proves that R_1, \dots, R_4 satisfy the theorem, and so proves 7.1. ■

By combining 7.1 and 6.5, we deduce:

7.2 *Let $k \geq 2$ be even. Let G be a 2-weak loopless didrawing with interleaving at most 2λ , where λ is an odd integer. Let $\ell = 40\lambda^2k^2$. Let G^* be a dual didrawing, let C be a cycle of G^* with change number at most ℓ , and let Δ be a disc bounded by C . Suppose that, for every path P of $G^*[\Delta]$ with distinct ends in $V(C)$ and with no internal vertex or edge in C , there is a subpath Q of C between the ends of P such that $P \cup Q$ has change number more than ℓ . Then either:*

- *the interior of Δ is a region of G^* ; or*
- *G has a subdigraph that is a directed subdivision of the $k \times k$ diwall.*

Proof. Let $n = 6\lambda k^2 - 2\lambda k + 9k - 2$, and $K = \lambda n$. Thus, $\ell \geq 4K + 2\lambda + 10$. We assume the first bullet is false, and so by 7.1 applied to G^* , there are four subpaths R_1, R_2, R_3, R_4 of G^* , pairwise edge-disjoint, with union C , and numbered in order in C , such that there is no path in $G^*[\Delta]$ between $V(R_1), V(R_3)$ with change number at most K , and there is no path in $G^*[\Delta]$ between $V(R_2), V(R_4)$ with change number at most K . Since $K \geq \lambda \max(n, k)$, we deduce from 6.5, applied to the didrawing obtained from G by contracting all the edges of G that are disjoint from Δ , that G has a subdigraph that is a directed subdivision of the $k \times k$ diwall. This proves 7.2. ■

Now we can deduce 3.1, which we restate:

7.3 *For all even $k \geq 2$, and every odd integer $\lambda \geq 1$, every loopless 2-weak didrawing with interleaving at most 2λ and with diwidth more than $40\lambda^2k^2$ embeds the $k \times k$ diwall.*

Proof. Since G is 2-weak, it has a vertex v , and the set of edges incident with v is a bond of G^\natural . Hence there is a ternary tree T and a surjective map τ from $V(G)$ to the set of leaves of T , such that for each edge e of T , if T_1, T_2 are the components of $T \setminus e$, then the set of edges of G between

$$(\{v \in V(G) : \tau(v) \in V(T_1)\}, \{v \in V(G) : \tau(v) \in V(T_2)\})$$

is a bond of G with change number at most $40\lambda^2k^2$. (This is because we can take T to be a two-vertex tree, and map v to one vertex of T and all others to the other vertex of T , and $2\lambda \leq 40\lambda^2k^2$.) Choose T maximal with this property. We claim that τ is a bijection. Suppose not; then for some leaf t of T , at least two vertices of G are mapped to t by τ . Let e be the edge of T incident with t , and let $A = \{v \in V(G) : \tau(v) = t\}$ and $B = V(G) \setminus A$. Thus, $|A| \geq 2$, and the set of edges of G between A, B is a bond of G with change number at most $40\lambda^2k^2$. Let G^* be a dual didrawing, and let C be the cycle of G^* corresponding to this bond. Thus C has change number at most $40\lambda^2k^2$. Let Δ be the closed disc in Σ bounded by C that includes A . Since $|A| \geq 2$, Δ includes more than one region of G^* , and so by 7.2, we may assume that there is a path P of $G^*[\Delta]$ with distinct ends in $V(C)$ and with no internal vertex or edge in C , such that $P \cup Q$ has change number at most $40\lambda^2k^2$ for each of the subpaths Q of C between the ends of P . Let Q_1, Q_2 be these two subpaths, and for $i = 1, 2$, let Δ_i be the closed disc in Δ bounded by $P \cup Q_i$, and let A_i be the set of vertices of G in Δ_i (that is, that correspond under the duality to a region of G^* included in Δ_i). Add two new vertices t_1, t_2 to T , both adjacent to t , forming a ternary tree T' , and define τ' by

$$\begin{aligned} \tau'(v) &= \tau(v) \text{ if } v \in B \\ \tau'(v) &= t_1 \text{ if } v \in A_1 \\ \tau'(v) &= t_2 \text{ if } v \in A_2. \end{aligned}$$

Since $P \cup Q_i$ is a cycle of G^* with change number at most $40\lambda^2 k^2$, it follows that the set of edges between A_i and $V(G) \setminus A_i$ is a bond of G with change number at most $40\lambda^2 k^2$, contrary to the maximality of T . This proves 7.3. \blacksquare

8 The main result

Now we use 7.3 to deduce our main result, 3.2. We need the following construction. Let G be a weakly-2-edge-connected loopless didrawing, and let J be obtained from G as follows. Subdivide each edge uv twice, so it becomes a directed path of length three. This makes a 2-weak didrawing G_1 say. Next, for each vertex v of G , let a_1, \dots, a_n be its neighbours in G_1 in clockwise order (more exactly, since we allow parallel edges and so a_1, \dots, a_n may not all be distinct, in the clockwise order of the corresponding edges around v). Add edges to G_1 from a_i to a_{i+1} for $1 \leq i \leq n$ (where $a_{n+1} = a_1$). This makes G_2 . Now delete all the vertices of G from G_2 , forming J . Then J is a 2-weak didrawing (this is easily checked), and each vertex v of G corresponds to a directed cycle C_v of J that bounds a region of J in a clockwise direction; and each edge of G corresponds to an edge of J between the corresponding two cycles. Each vertex of J has indegree two and outdegree one, or vice versa. We can recover G from J by contracting the edges of the cycles C_v ($v \in V(G)$).

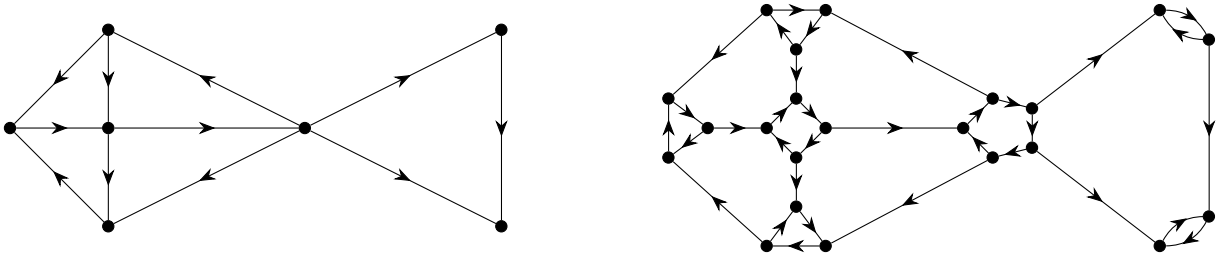


Figure 11: Constructing J from G .

The didrawing J is 2-weak and loopless, and has interleaving at most two; so by 7.3 applied to it, we deduce that for every even $k \geq 2$, either J has diwidth at most $40(9k)^2$ or J embeds the $9k \times 9k$ diwall. Thus, to complete the proof of 3.2, it suffices to prove the following two statements, the first of which is trivial.

8.1 *The dart-width of G is at most the diwidth of J .*

Proof. Let (T, τ) be a carving of $V(J)$ such that for each edge e of T , the set of edges of J joining the two parts of the corresponding partition is a bond of J with change number at most k . There is a natural bijection between the darts of G and the vertices of J , so we can view (T, τ) as a carving of the set of darts of G ; and as such, it shows that the dart-width of G is at most the diwidth of J , as required. \blacksquare

8.2 *For all even $k \geq 2$, if J embeds the $3k \times 9k$ diwall, then G embeds the $k \times k$ diwall.*

Proof. Suppose that J embeds the $3k \times 9k$ diwall, and let H be a subdigraph of J that is a subdivision of this diwall. Let P_1, \dots, P_{3k} and Q_1, \dots, Q_{9k} be its horizontal and vertical paths, numbered in order. Thus they form a $3k \times 9k$ diwall layout. For each $v \in V(G)$, the cycle C_v of J bounds a region of J , which is therefore a subset of one of the regions of the (unique, up to homeomorphism) didrawing of H . One of the regions of H is the “outer” region (in the natural sense), but the boundaries of all its “inner” regions consist of subpaths of two consecutive of P_1, \dots, P_{3k} and two consecutive of Q_1, \dots, Q_{9k} . Thus, for $v \in V(G)$, either C_v intersects only some of P_1, P_{3k}, Q_1, Q_{9k} , or it intersects at most two consecutive members of P_1, \dots, P_{3k} and two consecutive members of Q_1, \dots, Q_{9k} . Hence, the paths $P_1, P_4, P_7, \dots, P_{3k-2}$ and $Q_1, Q_4, Q_7, \dots, Q_{9k-2}$ form a $k \times 3k$ diwall layout, and each C_v intersects at most one of $P_1, P_4, P_7, \dots, P_{3k-2}$ and at most one of $Q_1, Q_4, Q_7, \dots, Q_{9k-2}$. The union of these paths is not a subdivision of a diwall, because (for instance) the last vertex of P_1 is not in any of $Q_1, Q_4, Q_7, \dots, Q_{9k-2}$; but by removing some vertices from the ends of these paths, we obtain a $k \times 3k$ diwall layout R_1, \dots, R_k and S_1, \dots, S_{3k} whose union is a $k \times 3k$ diwall K say.

When we contract the edges of all the cycles C_v , we may no longer contain a $k \times 3k$ diwall layout; but we claim that we still contain a $k \times 3k$ box system and hence a $k \times k$ diwall layout, by 5.1. To see this, let us examine in more detail the effect on R_1, \dots, R_k and S_1, \dots, S_{3k} of contracting the edges of some C_v . In any didrawing, if a region r is incident with an edge e , we can speak of r being to the left or right of e in the natural sense (recall that we have assigned an orientation “clockwise” to Σ). In our case, if C_v meets some R_i or S_j , then it shares at least one edge with that path, since K has maximum total degree three.

Now, C_v bounds a region r_v of J that is a subset of a region r'_v of the didrawing of K . Since C_v bounds r_v in the clockwise direction from the construction, it follows that r_v (and hence r'_v) is on the right of any edge of K that belongs to C_v . Since the outer region of K is on the left of all edges of K incident with it, r'_v is an inner region if C_v shares any vertices at all with R_1, \dots, R_k or S_1, \dots, S_{3k} .

There are four kinds of finite region in the didrawing K (see Figure 2); those bounded by a clockwise cycle (coming from a cycle of length eight in the diwall before subdivision), those bounded by an anticlockwise cycle of “unsubdivided length” eight, and those bounded by cycles of unsubdivided length four (and there are two kinds of the last, depending whether the region is to the right or left of the edges in horizontal paths incident with it). We can assume that C_v shares an edge with at least one of R_1, \dots, R_k or S_1, \dots, S_{3k} , and so r'_v is not bounded by an anticlockwise cycle; and if r'_v is bounded by a cycle with unsubdivided length four, then it only intersects one of R_1, \dots, R_k or S_1, \dots, S_{3k} and will not cause us a problem. The awkward case is when r'_v is bounded by a clockwise cycle of unsubdivided length eight (see Figure 12.)

Let r' be some such region of K . There are (in general) eight vertices of total degree three incident with r' , dividing its boundary into eight paths labelled as in the figure.

Let r' have boundary formed by subpaths of $R_i, R_{i+1}, S_j, S_{j+1}$. Thus i, j are both odd, since r' is bounded by a clockwise cycle of unsubdivided length eight. In the figure,

- R_i contains M_1, L_2, L_3, L_4, M_4 ;
- S_{j+1} contains M_3, L_4, L_5, L_6, M_6 ;
- R_{i+1} contains M_5, L_6, L_7, L_8, M_8 ; and
- S_j contains M_7, L_8, L_1, L_2, M_2 .

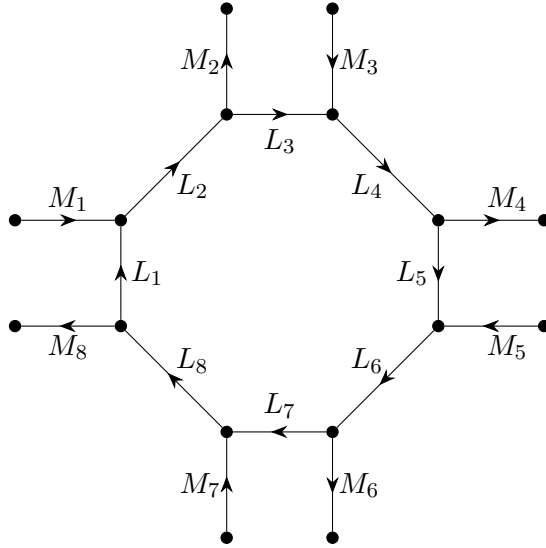


Figure 12: A piece of a subdivided diwall.

Let X be the set of all $v \in V(G)$ with $r'_v = r'$. For each $v \in X$, since C_v does not intersect both R_i, R_{i+1} , and does not intersect both S_j, S_{j+1} , it follows that the set of edges of $L_1 \cup \dots \cup L_8$ in C_v is included in one of $L_1 \cup L_2 \cup L_3$, $L_3 \cup L_4 \cup L_5$, $L_5 \cup L_6 \cup L_7$, $L_7 \cup L_8 \cup L_1$. From planarity, for $h = 1, 3, 5, 7$ we can choose an edge of L_h that divides it into two vertex-disjoint subpaths L'_h, L''_h , such that $L'_1 \cup L_2 \cup L'_3$, $L''_3 \cup L_4 \cup L'_5$, $L''_5 \cup L_6 \cup L'_7$, $L''_7 \cup L_8 \cup L'_1$ are directed paths, and for all $v \in X$, the set of edges of $L_1 \cup \dots \cup L_8$ in C_v is included in one of these four paths.

For $1 \leq i \leq k$ and $1 \leq j \leq 3k$ with i, j both odd,

- let $A_{i,j}$ be the union of $V(L'_1 \cup L_2 \cup L'_3)$ and all internal vertices of M_1 ;
- let $A_{i,j+1}$ be the union of $V(L''_3 \cup L_4 \cup L'_5)$ and all internal vertices of M_3 ;
- let $A_{i+1,j+1}$ be the union of $V(L''_5 \cup L_6 \cup L'_7)$ and all internal vertices of M_5 ; and
- let $A_{i+1,j}$ be the union of $V(L''_7 \cup L_8 \cup L'_1)$ and all internal vertices of M_7 .

This defines $A_{i,j}$ for all i, j with $1 \leq i \leq k$ and $1 \leq j \leq 3k$ (with a few natural adjustments when $i = 1, k$ or $j = 1, 3k$ which we omit). Hence there is a $k \times 3k$ box system in G where the vertex in the i th horizontal path and j th vertical path of $\Gamma_{k,3k}$ is mapped to $A_{i,j}$. This box system has the property that for each $v \in V(G)$, one of the boxes $A_{i,j}$ includes the set of all vertices in $C_v \cap K$, and so this remains a box system after contracting the edges of all the C_v . This proves 8.2, and hence proves 3.2. ■

9 Grids

We all agree what an undirected grid is, but there are several ways to define a directed grid. And while every undirected planar graph is a minor of a big enough undirected grid, this is by no means true for planar digraphs and directed minor containment and directed grids.

A $k \times k$ grid has k “horizontal” paths and k “vertical” paths, defined in the natural way, and let us make them all directed paths, but how should we direct them?

Three types that we care about are:

- All the horizontal paths are directed from left to right; we call this the $k \times k$ *acyclic* grid.
- Assume k is even, and the top $k/2$ horizontal paths are directed from left to right, and the others from right to left, and similarly the $k/2$ leftmost vertical paths are directed from bottom to top, and the others from top to bottom. (See Figure 4.) Let us call this the $k \times k$ *semi-grid*.
- The $k \times k$ *alternating* grid, defined earlier.

For the acyclic grid, in a sense it makes no difference how we direct the vertical paths, because for any acyclic grid H , if k is large enough then every $k \times k$ acyclic grid embeds H . To see this, let A_k be the $k \times k$ acyclic grid when all vertical paths are directed from top to bottom, and let B_k be the one when they alternate in direction. If $k' \geq 2k - 1$ then clearly, every $k' \times k'$ acyclic grid embeds A_k ; and for every $k \times k$ acyclic grid H , if $k' \geq 2k - 1$ then $B_{k'}$ embeds H . But one can show that for all k there exists $\ell = O(k^2)$ such that A_ℓ embeds B_k . Thus, up to embedding, these types of grid are all equivalent. But we do not know any analogue of 1.1 for acyclic grids. (That is an interesting question, perhaps: is there a version of 1.1 for acyclic digraphs?)

The semi-grid is equivalent (up to embedding, as before) to the grid when we direct the horizontal paths as before (that is, the top $k/2$ horizontal paths are directed from left to right, and the others from right to left), but have the vertical paths alternate in direction. It is also equivalent in the same sense to what is called the $k \times 2k$ *cylindrical grid*, defined as follows. Take k disjoint cycles drawn in the plane, each inside the next and each rotating clockwise; and then take $2k$ directed paths in the plane between the outside cycle and the inside one, alternating in direction and each intersecting each of the cycles in just one vertex. Again, a large enough semi-grid embeds a large cylindrical grid, and vice versa. For this one there is an analogue of 1.1, that we discuss next.

10 Minor containment, and the Kawarabayashi-Kreutzer theorem

Minor containment in digraphs is much more complicated than in undirected graphs. We agree what a minor of an undirected graph means, but for digraphs there are at least two definitions of minor containment of interest, “butterfly” minors and “strong” minors, and here we want a third, “semi-strong” minors.

- A *butterfly contraction* is the operation on a digraph of contracting a non-loop edge $e = uv$ with the property that either u has out-degree one or v has in-degree one. A digraph H is a *butterfly minor* of a digraph G if a digraph isomorphic to H can be obtained from a subdigraph of G by a sequence of butterfly contractions.
- A *strong contraction* is the operation on a digraph of contracting a non-null strongly-connected subdigraph to a single vertex. A digraph H is a *strong minor* of a digraph G if a digraph isomorphic to H can be obtained from a subdigraph of G by a sequence of strong contractions.
- H is a *semi-strong minor* of a digraph G if some directed subdivision of H is a strong minor of G .

Not every didrawing is a strong minor of a diwall, or a butterfly minor of a diwall, but each is a semi-strong minor of a large enough diwall (we omit the proof), and that is the reason we want semi-strong minors. So a consequence of our main theorem is that:

10.1 *For every didrawing H there exists k such that every weakly 2-edge-connected loopless didrawing that does not contain H as a semi-strong minor has dart-width at most k .*

It is easy to show the weak converse, that large enough diwalls have large dart-width, so our result 10.1 is thus a sort of analogue of the theorem of Robertson and Seymour [7] for planar digraphs.

There is at least one other way to extend that theorem to digraphs, due to Kawarabayashi and Kreutzer [5, 6], using butterfly minors and the semi-grid of Figure 4. It needs “directed tree-width” which is a complicated parameter introduced in [4]. (We omit the definition.) Kawarabayashi and Kreutzer [5, 6] proved a conjecture of [4], that:

10.2 *For every $k \geq 0$ there exists $\ell \geq 0$ such that every digraph with directed tree-width at least ℓ contains a $k \times k$ semi-grid as a butterfly minor. Conversely, for every $\ell \geq 0$ there exists $k \geq 0$ such that every digraph that contains a $k \times k$ semi-grid as a butterfly minor has directed tree-width at least ℓ .*

Consequently, we have the same conclusion if we exclude a digraph H that is a butterfly minor of a semi-grid. Such digraphs H are the digraphs that can be drawn in the plane without crossings, such that every edge subtends a positive angle at the origin, and for every vertex v , the out-edges at v and in-edges at v fall into two disjoint intervals in the circular order around v (we omit the proof). Let us call such digraphs *circular*. Thus, the theorem of Kawarabayashi and Kreutzer says:

10.3 *For every circular digraph H , there exists k such that every digraph that does not contain H as a butterfly minor has directed tree-width at most k .*

Our result 10.1 says something similar to this, but is different in three ways, one gain, one loss and one tie:

- it involves semi-strong minors instead of butterfly minors (a tie);
- the digraph H we exclude can be any planar digraph H (a gain);
- it only tells us how to decompose *planar* digraphs that do not contain H as a semi-strong minor (a loss).

Finally, let us say a didrawing with interleaving at most two is *normal*. Thus, circular didrawings are normal. The didrawings that are butterfly minors of diwalls are precisely the normal didrawings (again, we omit the proof), so there is another consequence of our theorem:

10.4 *For every normal didrawing H , there exists k such that every loopless, 2-weak didrawing not containing H as a butterfly minor has dart-width at most k .*

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