

Random Cayley graphs and random sumsets

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Abstract

We prove that any finite abelian group G contains a collection \mathcal{F} of not too many subsets with a special structure, so that for every subset A of G with a small doubling, there is a member $F \in \mathcal{F}$ that is fully contained in the sumset $A + A$ and is not much smaller than it. Using this result we obtain improved bounds for the problem of estimating the typical independence number of sparse random Cayley or Cayley-sum graphs, and for the problem of estimating the smallest size of a subset of G which is not a sumset. We also obtain tight bounds for the typical maximum length of an arithmetic progression in the sumset of a sparse random subset of G .

1 Introduction

Given an abelian group G and a finite subset A of G , define the sumset of A

$$A + A = \{a + b : a, b \in A\},$$

and the associated doubling constant $K = \frac{|A+A|}{|A|}$. Sumsets and sets with small doubling K are of fundamental interest in additive combinatorics. Over the years, multiple aspects of the structure of sumsets and sets with small doubling have been studied. A notable result is the influential Freiman-Ruzsa theorem [9, 10, 23], generalized by Ruzsa [22] and Green and Ruzsa [16], which shows that sets A of constant doubling K must be dense subsets of certain structured objects known as coset progressions.

While structural results typically provide information about the set A given its doubling K , they provide relatively weak information about the sumset $A + A$. Motivated by fundamental applications to the investigation of sparse random Cayley graphs [1, 2, 13, 15], we study in this paper new perspectives on the structure of sumsets $A + A$ of sets with small doubling. Our main result provides an answer to the following basic question: Does every sumset $A + A$ of a set A with small doubling K fully contain a *dense* structured

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subset F ? In exploring the inherent structure of the sumset $A + A$, we also study natural questions about sumsets $A + A$ of binomial random sets A . Before describing the precise results, we first describe the motivating problems for our work.

Sparse random Cayley graphs. Given an abelian group G and a symmetric subset of it S , the Cayley graph $\Gamma(G; S)$ of G with generating set S has vertex set G , and two group elements x and y are connected if and only if $y - x \in S$. A random Cayley graph $G(p)$ is obtained by selecting each equivalence class $\{x, -x\}$ to be in the generating set S independently at random with probability p .

We also consider here the Cayley sum graph $\Gamma^+(G; S)$ which, given a generating set S (not necessarily symmetric), has vertex set G and two group elements x and y are connected if and only if $x + y \in S$. A random Cayley sum graph $G^+(p)$ is obtained by selecting each element x to be in S independently at random with probability p .

The independence number $\alpha(G(p))$ of random Cayley graphs has been extensively studied. In the dense case $p = 1/2$ (and more generally $p = \Theta(1)$), the first author and Orlitsky [4] showed that $\alpha(G(1/2)) = O((\log |G|)^2)$ with high probability. Here and in what follows we say that an event holds with high probability (whp, for short) if the probability it holds tends to 1 as the relevant parameter tends to infinity. Green [13] showed that whp $\alpha(G(1/2)) = \Theta(\log |G|)$ for cyclic groups $G = \mathbb{Z}_n$ and that $\alpha(G(1/2)) = \Theta(\log |G| \log \log |G|)$ for finite field vector spaces $G = \mathbb{F}_2^d$ (a similar result holds for any finite field vector space $G = \mathbb{F}_p^d$ with p fixed). Green and Morris [15] later sharpened Green's result to show that whp $\alpha(G(1/2)) = (2 + o(1)) \log_2 |G|$ for $G = \mathbb{Z}_n$. The problem is significantly harder in the sparse case $p = o_{|G|}(1)$. The best general result in this direction is the following theorem of the first author [1, 2].

Theorem 1. *Let G be a group of size n . The independence number of the random Cayley graph $G(p)$ is at most $O(\min(p^{-2}(\log n)^2, \sqrt{n(\log n)/p}))$ whp.*

It is natural to conjecture that, in terms of the independence number, random Cayley graphs behave similarly to random regular graphs of the same degree. Here and throughout the paper, we often hide polylogarithmic factors in $|G|$ using the \tilde{O} and $\tilde{\Omega}$ notation.

Conjecture 2 ([2]). *Let G be a group of size n . The independence number of the random Cayley graph $G(p)$ is at most $\tilde{O}(p^{-1})$ whp.*

Unlike random regular graphs, random Cayley graphs, and especially sparse random Cayley graphs, are significantly harder to analyze due to the limited randomness in their definition which causes significant dependencies.

In a recent work [7], motivated by Ramsey-theoretic applications, Conlon, Fox, the second author and Yepremyan gave an improvement of Theorem 1 for general groups G .

Theorem 3. *Let G be a group of size n . The independence number of the random Cayley graph $G(p)$ is at most $O\left(p^{-2} \log n \max\left(\log p^{-1}, p^{-1} \log\left(\frac{\log n}{\log p^{-1}}\right)\right)\right)$ whp.*

This result is generally tight for the dense case $p = \Theta(1)$. On the other hand, for sparse p , such as $p = n^{-c}$ for some $c > 0$, the result only gives lower order improvements. No improvement over the exponent p^{-2} in Theorem 1 has been obtained so far.

As one application of our key result, we obtain the first improvement in the exponent of p .

Theorem 4. *Let G be an abelian group of size n and let $p \leq 1/2$. Then the independence number of the random Cayley graph $G(p)$ and the random Cayley sum graph $G^+(p)$ is at most $\tilde{O}(p^{-3/2})$ whp.*

Largest non-sumsets in \mathbb{Z}_n . One of the applications of the independence number of polynomially sparse random Cayley sum graphs is toward the following question in additive combinatorics first considered by Green [11]. Let $f(n)$ be the largest integer such that every subset of \mathbb{Z}_n with size larger than $n - f(n)$ can be represented as a sumset $A + A$. Green asked to determine or estimate $f(n)$ and showed that $f(n) \geq \Omega(\log n)$.

The first author proved that $f(n) \geq \tilde{\Omega}(n^{1/2})$ and, via the upper bound $\alpha(G^+(p)) = O(p^{-2}(\log n)^2)$, that $f(n) \leq \tilde{O}(n^{2/3})$. Using Theorem 4, we obtain an improvement in the exponent of the upper bound to $f(n)$.

Theorem 5. *In the notation above, $f(n) \leq \tilde{O}(n^{3/5})$. Specifically, let G be an abelian group of size n . Then there exists a subset of G of size at least $n - \tilde{O}(n^{3/5})$ which cannot be represented as a sumset $A + A$ for $A \subset G$.*

The key result. Our key input to Theorem 4 is a structural result showing that every sumset $A + A$ of a set A with small doubling must fully contain a *dense structured* subset F .

Theorem 6. *There exist $C, c > 0$ such that the following holds. Let G be an abelian group of order n and let $s \leq n$. There exist collections \mathcal{F}_ℓ of subsets of G such that*

$$|\mathcal{F}_\ell| \leq \exp \left(C \left(\min(2^{2\ell}(\log n)^2, \sqrt{2^\ell s(\log n)^{3/2}}) \right) \right),$$

and

$$\min_{F \in \mathcal{F}_\ell} |F| \geq c2^\ell s / \ell^2,$$

so that the following property holds.

Let $A \subseteq G$ be such that $|A + A| \leq K|A|$ and $|A| = s$. Then there exists $\ell \leq \log_2 K$ and $F \in \mathcal{F}_\ell$ such that

$$A + A \supseteq F.$$

Informally speaking, the theorem yields that for every $s \leq n$ and every K , there exists a dyadic scale $h = 2^\ell \leq K$ and a collection of sets \mathcal{F} with complexity $\log |\mathcal{F}| \leq \tilde{O}(\min(h^2, \sqrt{hs}))$, such that for every subset A of size s and doubling at most K , $A + A$

fully contains a set from \mathcal{F} of size at least $\tilde{\Omega}(h|A|)$. A particularly surprising aspect of this result is the case the doubling K is bounded. In this case we obtain that $A + A$ contains a structured subset F which is dense in $A + A$, where F lies in a collection of sets of complexity $\tilde{O}(1)$ (which, crucially, is independent of A).

Indeed, Lovett [19] asked if it is possible to find a small collection \mathcal{F} of dense sets such that, for every dense subset $A \subseteq G$, the sumset $A + A$ contains a member of \mathcal{F} . For every dense A , the doubling of A is clearly bounded. As a special case of Theorem 6, we thus resolve Lovett's question.

Theorem 7. *Let G be an abelian group of order n . For any $\delta > 0$, there are $\epsilon > 0$ and $C > 0$ such that the following holds. There exists a collection $\mathcal{F}_\delta \subseteq 2^G$ consisting of sets of size at least ϵn with $|\mathcal{F}_\delta| \leq \exp(C(\log n)^2)$ so that for every $|A| \geq \delta n$, $A + A$ fully contains a set $F \in \mathcal{F}_\delta$.*

It is easy to see that any collection \mathcal{F}_δ satisfying the property in Theorem 7 must have size at least $\exp(\omega_\delta(1)(\log n))$. For example, if $G = \mathbb{F}_2^d$, consider subsets A given by codimension $\log_2(1/\delta)$ subspaces. Each set F of size $|F| \geq \epsilon n$ can be a subset of at most $2^{\log_2(1/\epsilon) \log_2(1/\delta)}$ subspaces of codimension $\log_2(1/\delta)$. On the other hand, the number of subspaces of codimension $\log_2(1/\delta)$ is at least $2^{d \log_2(1/\delta) - (\log_2(1/\delta))^2}$. Hence we get

$$|\mathcal{F}_\delta| \geq 2^{d \log_2(1/\delta) - \log_2(1/\epsilon) \log_2(1/\delta) - \log_2(1/\delta)^2}.$$

Theorem 6 connects directly with sparse random Cayley graphs via providing an efficient *union bound obstruction* or *small cover* to the existence of large independent sets in random Cayley graphs. In probabilistic combinatorics language, a cover for a collection \mathcal{H} of sets is a collection \mathcal{G} such that every set in \mathcal{H} contains a set in \mathcal{G} as a subset. In our case the collection \mathcal{H} is the collection of sumsets $A + A$ over sets A of suitable size. In this language, the collection $\bigcup_\ell \mathcal{F}_\ell$ describes a cover for the collection \mathcal{H} . By the union bound over $\bigcup_\ell \mathcal{F}_\ell$, the random Cayley sum graph $G^+(p)$ typically does not contain an independent set of size s if this cover is small, that is, if

$$\sum_\ell \sum_{F \in \mathcal{F}_\ell} (1-p)^{|F|} = o(1).$$

Union bound obstructions play a crucial role in the study of thresholds. In particular, the Kahn-Kalai conjecture [17], proved in [21], implies direct connections between thresholds and union bound obstructions. We expect that the independence number of random Cayley graphs can be accurately determined via an optimally efficient cover for the collection \mathcal{H} of sumsets (see Conjecture 15 in the final section).

Arithmetic progressions in random sumsets. Our results here suggest that, even for arbitrary sets A , the sumset $A + A$ contains large nontrivial structures of low complexity. In the investigation of the structure of the collection of sumsets $A + A$, it is natural to

study the behavior of typical sumsets, that is, the sumset $A + A$ of a random set A . In the cyclic group \mathbb{Z}_p , we consider a random set A obtained by including each element of \mathbb{Z}_p independently with probability q . Here we study specifically the length of the longest arithmetic progression contained in $A + A$.

For arbitrary dense A , the length of the longest arithmetic progression in $A + A$ has been extensively studied [12, 8]. In a recent paper of Kohayakawa and Miyazaki [18] the authors consider the problem of estimating the typical maximum length $ap(A + A)$ of an arithmetic progression of the sumset $A + A$ of a random subset A of $[n] = \{1, 2, \dots, n\}$.

The authors of [18] show that if every element is chosen to lie in A , randomly and independently, with probability $q = q(n)$, then $ap(A + A)$ exhibits a sharp change of behavior around $q = \frac{1}{\sqrt{n}}$. In particular, they prove that for

$$q = \frac{1}{\sqrt{n(\log n)^{\Theta(1)}}}$$

the maximum length is $\Theta(\frac{\log n}{\log \log n})$ whp, and that for $q \geq \sqrt{\frac{4 \log n}{n}}$ the maximum is $\Theta(n)$ whp. For $q = \frac{1}{\sqrt{n}}$ their arguments only suffice to establish an $\Omega(\log n / \log \log n)$ lower bound and an $O(n)$ upper bound.

Here we prove that for this probability the correct answer is $\Theta(\log n)$ whp. It is more convenient to study the closely related analogous problem for the finite cyclic group \mathbb{Z}_p for prime p . The same arguments provide a similar result for a random subset of $[n]$.

Let p be a large prime, and let A be a random subset of \mathbb{Z}_p , where each $a \in \mathbb{Z}_p$ lies in A , randomly and independently, with probability $q = 1/\sqrt{p}$. Let $ap(A + A)$ denote, as above, the maximum length of an arithmetic progression in $A + A$.

Theorem 8. *There exist two absolute positive constants c_1, c_2 so that for A as above, the maximum length $ap(A + A)$ of an arithmetic progression in $A + A$ satisfies, whp,*

$$c_1 \log p \leq ap(A + A) \leq c_2 \log p.$$

Paper organization. In Section 2 we prove our key result, Theorem 6 and Theorem 7. In Section 3 we prove Theorems 4 and 5 as applications of the main covering lemma Theorem 6. In Section 4 we study random sumsets and prove Theorem 8. The final Section 5 contains some concluding remarks and open problems.

2 The main covering lemma

2.1 Preliminaries on discrete Fourier analysis

Let G be an abelian group of order n . Denote by \widehat{G} the group of characters $\chi : G \rightarrow \mathbb{C}$. Given $f : G \rightarrow \mathbb{C}$, the Fourier transform of f is defined as

$$\widehat{f}(\chi) = \mathbb{E}[f(x)\chi(x)].$$

For the convolution

$$f * g(x) = \mathbb{E}_y[f(y)g(x-y)],$$

we have that

$$\widehat{f * g}(\chi) = \widehat{f}(\chi)\widehat{g}(\chi).$$

Parseval's identity asserts that

$$\sum_{\chi} |\widehat{f}(\chi)|^2 = \mathbb{E}[|f(x)|^2].$$

Finally, we have the Fourier inversion formula

$$f(x) = \sum_{\chi \in \widehat{G}} \widehat{f}(\chi) \chi(-x).$$

2.2 The covering theorem

In this section we prove the main covering results, Theorem 6 and Theorem 7. Throughout the section, G is an abelian group of size n and A is a subset of G . Identifying A with its characteristic function, observe that $A + A = \text{supp}(A * A)$ and $A - A = \text{supp}(A * (-A))$. Our main tool is the following result giving Fourier sparse pointwise approximation to the convolutions $A * A$ and $A * (-A)$.

Lemma 9. *Let $\alpha \in (0, 1)$, $\eta \in (0, 1)$ and $m = 4\eta^{-2}(\log n)$. Then for all $A \subseteq G$ of size αn , there exist χ_1, \dots, χ_m such that*

$$\max_x \left| A * (-A)(x) - \frac{\alpha}{m} \sum_{i=1}^m \chi_i(-x) \right| \leq \eta\alpha. \quad (1)$$

Similarly, there exist χ_1, \dots, χ_m such that

$$\max_x \left| A * A(x) - \frac{\alpha}{m} \sum_{i=1}^m \frac{\widehat{A}(\chi_i)^2}{|\widehat{A}(\chi_i)|^2} \chi_i(-x) \right| \leq \eta\alpha. \quad (2)$$

Proof. We have

$$A * (-A)(x) = \sum_{\chi} |\widehat{A}(\chi)|^2 \chi(-x).$$

Note that $\sum_{\chi} |\widehat{A}(\chi)|^2 = \alpha$ by Parseval's identity. For each $i \in [m]$, pick χ_i independently such that

$$\mathbb{P}(\chi_i = \chi) = \alpha^{-1} |\widehat{A}(\chi)|^2.$$

Then

$$\mathbb{E} \left[\frac{\alpha}{m} \sum_{i=1}^m \chi_i(-x) \right] = A * (-A)(-x).$$

Furthermore, noting that $|\chi(-x)| \leq 1$, by the Hoeffding bound:

$$\mathbb{P} \left[\left| \frac{\alpha}{m} \sum_{i=1}^m \chi_i(-x) - A * (-A)(x) \right| > t \frac{\alpha}{\sqrt{m}} \right] \leq 2 \exp(-t^2/2).$$

By the union bound, we then have that with positive probability,

$$\max_x \left| A * (-A)(x) - \frac{\alpha}{m} \sum_{i=1}^m \chi_i(-x) \right| \leq 2\sqrt{\log n} \cdot \frac{\alpha}{\sqrt{m}} \leq \eta\alpha.$$

For (2), we follow the same proof, noting that

$$\mathbb{E} \left[\frac{\alpha}{m} \sum_{i=1}^m \frac{\widehat{A}(\chi_i)^2}{|\widehat{A}(\chi_i)|^2} \chi_i(-x) \right] = A * A(x). \quad \square$$

We are now ready to prove the main covering theorem. Informally, the theorem asserts that the collection of difference sets $\{A - A : |A| = \alpha n, |A - A| \leq K|A|\}$ admits a small cover. For $A \subseteq G$ and an integer ℓ , let

$$A_\ell = \{x : A * (-A)(x) \in (2^{-\ell-1}\alpha, 2^{-\ell}\alpha]\}.$$

We start with the following simple lemma.

Lemma 10. *There is an absolute constant $c > 0$ such that the following holds. Let G be an abelian group of order n . Let $\alpha \in (0, 1)$. Let $A \subseteq G$ be such that $|A - A| \leq K|A|$ and $|A| = \alpha n$. Then there exists $\ell \leq \log_2 K$ such that $|A_\ell| \geq c2^\ell \alpha n / \ell^2$.*

Proof. Note that $\mathbb{E}_x A * (-A)(x) = \alpha^2$. Thus, if for all $\ell \leq \log_2 K$,

$$|A_\ell| < c2^\ell \alpha n / \ell^2,$$

then

$$\alpha^2 n < \sum_{\ell \leq \log_2 K} 2^{-\ell} \alpha \cdot c2^\ell \alpha n / \ell^2 + 2^{-(\log_2 K + 1)} \alpha K |A| < \alpha^2 n,$$

which is a contradiction for a suitable constant c . \square

For $A \subseteq G$ such that $|A - A| \leq K|A|$ and $|A| = \alpha n$, we denote by $\ell(A)$ the smallest index $\ell \leq \log_2 K$ such that $|A_\ell| \geq c2^\ell \alpha n / \ell^2$.

Theorem 11. *There are absolute constants $C, c > 0$ such that the following holds. Let G be an abelian group of order n . Let $\alpha \in (0, 1)$. There exist collections \mathcal{F}_ℓ of subsets of G such that*

$$|\mathcal{F}_\ell| \leq \exp(C2^{2\ell}(\log n)^2),$$

and

$$\min_{F \in \mathcal{F}_\ell} |F| \geq c2^\ell \alpha n / \ell^2,$$

and the following property holds.

Let $A \subseteq G$ be such that $|A - A| \leq K|A|$ and $|A| = \alpha n$. Then for $\ell = \ell(A)$ there is $F \in \mathcal{F}_\ell$ such that

$$A - A \supseteq F \supseteq A_\ell.$$

Proof. For each ℓ , we define \mathcal{G}_ℓ as the collection of functions of the form

$$\frac{\alpha}{m} \sum_{\chi \in S} \chi(-x),$$

for $S \subseteq \widehat{G}$ of size $2^{2\ell+6}(\log n)$. Let \mathcal{F}_ℓ denote the collection of sets of the form

$$\{x : \widehat{f}(x) > 2^{-\ell-2}\alpha\},$$

for $\widehat{f} \in \mathcal{G}_\ell$.

For each set A with $|A - A| \leq K|A|$, let $\ell = \ell(A)$ be as in Lemma 10. For $\eta = 2^{-\ell-2}$, we then find $S \subseteq \widehat{G}$ of size $4\eta^{-2}(\log n)$ such that

$$\left| A * (-A)(x) - \frac{\alpha}{m} \sum_{\chi \in S} \chi(-x) \right| \leq \eta\alpha.$$

Let $\widehat{f}_A(x) = \frac{\alpha}{m} \sum_{\chi \in S} \chi(-x)$ and note that $\widehat{f}_A \in \mathcal{G}_\ell$. Let

$$F(A) = \{x : \widehat{f}_A(x) > \eta\alpha\},$$

so $F(A) \in \mathcal{F}_\ell$. We then have

$$F(A) \supseteq \{x : A * (-A)(x) > 2\eta\alpha\} \supseteq A_\ell,$$

as desired. \square

A similar proof yields an analogous result for sumsets $A + A$. For $A \subseteq G$ and an integer ℓ , let

$$A_\ell^+ = \{x : A * A(x) \in (2^{-\ell-1}\alpha, 2^{-\ell}\alpha]\}.$$

Theorem 12. *There are constants $C, c > 0$ such that the following holds. Let G be an abelian group of order n . Let $\alpha \in (0, 1)$. There exists collections \mathcal{F}_ℓ of subsets of G such that*

$$|\mathcal{F}_\ell| \leq \exp(C2^{2\ell}(\log n)^2),$$

and

$$\min_{F \in \mathcal{F}_\ell} |F| \geq c2^\ell \alpha n / \ell^2,$$

and the following property holds.

Let $A \subseteq G$ be such that $|A + A| \leq K|A|$ and $|A| = \alpha n$. Then for $\ell = \ell(A)$ there is $F \in \mathcal{F}_\ell$ such that

$$A + A \supseteq F \supseteq A_\ell.$$

Proof. Let $\eta = c/n$ and let \mathcal{U} be an $\eta/32$ -net for $\{x \in \mathbb{C} : |x| = 1\}$ of size $O(\eta^{-1})$. For each ℓ , we define \mathcal{G}_ℓ as the collection of functions of the form

$$\frac{\alpha}{m} \sum_{\chi \in S} t_\chi \chi(-x),$$

for $S \subseteq \widehat{G}$ of size $2^{2\ell+6}(\log n)$ and $t_\chi \in U$ for each $\chi \in S$. Let \mathcal{F}_ℓ denote the collection of sets of the form

$$\{x : \widehat{f}(x) > 2^{-\ell-2}\alpha\},$$

for $\widehat{f} \in \mathcal{G}_\ell$. The rest of the proof is very similar to the proof of Theorem 11. \square

We also observe that in the range of large doubling, we have the following version of the covering theorem.

Theorem 13. *There are $C, c > 0$ satisfying the following. Let G be an abelian group of order n . Let $\alpha \in (0, 1)$. There exists collections \mathcal{H}_ℓ of subsets of G such that*

$$|\mathcal{H}_\ell| \leq \exp\left(C\sqrt{2^\ell|A|(\log n)^{3/2}}\right),$$

and

$$\min_{H \in \mathcal{H}_\ell} |H| \geq c2^\ell \alpha n / \ell^2,$$

and the following property holds.

Let $A \subseteq G$ be such that $|A - A| \leq K|A|$ and $|A| = \alpha n$. Then for $\ell = \ell(A)$ there is $H \in \mathcal{H}_\ell$ such that

$$A - A \supseteq H \supseteq A_\ell.$$

Proof. We construct

$$\mathcal{H}_\ell = \left\{ A' - A' : |A'| \leq \sqrt{2^{\ell+4}|A|\log n}, |A' - A'| \geq c2^\ell \alpha n / \ell^2 \right\}.$$

Let $|A| = \alpha n$ with $|A - A| \leq K|A|$. As in the proof of Theorem 11, there exists ℓ such that

$$|A_\ell| \geq c2^\ell \alpha n / \ell^2.$$

Consider a random subset A' of A where each element is sampled independently with probability $q = (2^{\ell+3}(\log n)/|A|)^{1/2}$. Then we have, for $x \in A_\ell$,

$$\mathbb{P}(x \notin A' - A') \leq (1 - q^2)^{2^{-\ell-1}|A|} \leq n^{-2}.$$

Thus, combining with the standard Chernoff bound, with high probability, $A' - A' \supseteq A_\ell$, and $|A'| \leq \sqrt{2^{\ell+4}|A|\log n}$.

In particular, there exists $H \in \mathcal{H}_\ell$ such that

$$A - A \supseteq H \supseteq A_\ell. \quad \square$$

By the same proof, we also obtain the sumset analog of Theorem 13.

Theorem 14. *There are $C, c > 0$ satisfying the following. Let G be an abelian group of order n . Let $\alpha \in (0, 1)$. There exists collections \mathcal{H}_ℓ of subsets of G such that*

$$|\mathcal{H}_\ell| \leq \exp \left(C \sqrt{2^\ell |A| (\log n)^{3/2}} \right),$$

and

$$\min_{H \in \mathcal{H}_\ell} |H| \geq c 2^\ell \alpha n / \ell^2,$$

and the following property holds.

Let $A \subseteq G$ be such that $|A + A| \leq K|A|$ and $|A| = \alpha n$. Then for $\ell = \ell(A)$ there is $H \in \mathcal{H}_\ell$ such that

$$A + A \supseteq H \supseteq A_\ell.$$

Theorems 11, 12, 13 and 14 imply Theorem 6. Theorem 7 is a direct corollary of Theorem 12, noting that $K \leq 1/\delta$ and $\ell(A) \leq \log_2 K$.

3 Applications of the main covering lemma

3.1 Independence number of random Cayley graphs

As mentioned earlier, the efficient cover constructed for the collection of difference sets or sumsets in Theorem 6 immediately yields improved bounds for the independence number of random Cayley graphs or random Cayley sum graphs, Theorem 4. Here, we deduce Theorem 4 based on Theorem 6 by an application of the union bound.

Proof of Theorem 4. Consider $A \subseteq G$ with $|A| = s := \xi p^{-3/2} (\log n)^{19/4}$. Let $\ell = \ell(A)$. Let ℓ_0 be so that 2^{ℓ_0} is within a factor 2 of $p^{-1/2} (\log n)^{3/4}$.

For $\ell \geq \ell_0$, let \mathcal{E}_ℓ denote the event that the complement of the random generating set of the Cayley graph S contains H for some $H \in \mathcal{H}_\ell$. For $\ell < \ell_0$, let \mathcal{E}_ℓ denote the event that the complement of the random generating set of the Cayley graph S contains F for some $F \in \mathcal{F}_\ell$. Then the event that the independence number of $\Gamma(G; S)$ is at least $s = \xi p^{-3/2} (\log n)^{19/4}$ is contained in the union of the events \mathcal{E}_ℓ for $\ell \geq \ell_0$ and \mathcal{E}_ℓ for $\ell < \ell_0$.

By the union bound, we have

$$\begin{aligned} & \mathbb{P}[\alpha(\Gamma(G; S)) > s] \\ & \leq \sum_{\ell < \ell_0} |\mathcal{F}_\ell| (1-p)^{\min_{F \in \mathcal{F}_\ell} |F|} + \sum_{\ell \geq \ell_0} |\mathcal{H}_\ell| (1-p)^{\min_{H \in \mathcal{H}_\ell} |H|} \\ & \leq \sum_{\ell < \ell_0} \exp \left(C 2^{2\ell} (\log n)^2 \right) \exp \left(-p c 2^\ell s / \ell^2 \right) \\ & \quad + \sum_{\ell \geq \ell_0} \exp \left(C \sqrt{2^\ell (\log n)^{3/2} s} \right) \exp \left(-p c 2^\ell s / \ell^2 \right). \end{aligned}$$

We next check that for a sufficiently large absolute constant ξ each of the two sums above is small. To bound the first sum it suffices to check that for every $\ell < \ell_0$

$$pc2^\ell s/\ell^2 \geq 2 \cdot C2^{2\ell}(\log n)^2$$

Since $\ell \leq \log n$ this is the case provided

$$pcs/(\log n)^2 > 2C2^\ell(\log n)^2$$

Substituting the values of $s = \xi p^{-3/2}(\log n)^{19/4}$ it follows that this is equivalent to the inequality $c\xi p^{-1/2}(\log n)^{3/4} > 2C2^\ell$ which clearly holds for each $\ell < \ell_0$, since 2^{ℓ_0} is within a factor of 2 of $p^{-1/2}(\log n)^{3/4}$.

Note that the main contribution for this sum is given by $\ell = 0$, implying that for an appropriate ξ this sum is bounded by $\exp(-\Omega(p^{-1/2}(\log n)^{11/4}))$.

In order to bound the second sum it suffices to check that for $\ell \geq \ell_0$,

$$2^\ell pcs/(\log n)^2 > 2C\sqrt{2^\ell(\log n)^{3/2}s}.$$

This is equivalent to the inequality

$$2^\ell \geq \frac{4C^2}{c^2\xi} p^{-1/2}(\log n)^{3/4},$$

which clearly holds for all $\ell \geq \ell_0$ by the choice of ℓ_0 , provided ξ is a sufficiently large constant. Here the main contribution for the sum is obtained for $\ell = \ell_0$, showing that the second sum is bounded by $\exp(-\Omega(p^{-1}(\log n)^{7/2}))$.

Altogether, this computation shows that

$$\mathbb{P}[\alpha(\Gamma(G; S)) > s] \leq \exp(-\Omega(p^{-1/2}(\log n)^{11/4})). \quad \square$$

Via a similar argument, we obtain similar bounds for the upper bound to the independence number of random Cayley sum graphs.

It is worth noting that the polylogarithmic factor above can be improved, we make here no serious attempt to optimize it.

3.2 Large sets which are not sumsets

Given Theorem 4, we give a short proof of Theorem 5, showing that there exists a subset of \mathbb{Z}_n of size $n - \tilde{O}(n^{3/5})$ that cannot be represented as a sumset.

Proof of Theorem 5. For a suitable value of p which we will choose later, we let S_1 be a random subset of G where each element is chosen with probability p . We will then let $A = G \setminus (S_1 \cup S_2)$, and show that there exists a choice of S_2 of size $\Theta(pn)$ for which A

cannot be written as a sumset $B + B$. Note that if $A = B + B$ then $B + B$ is disjoint from S_1 and hence B is an independent set in $\Gamma^+(G; S_1)$.

By Theorem 4 (and its proof above), with high probability, $|S_1| \leq 2pn$ and any independent set B in $\Gamma^+(G; S_1)$ has size at most $\xi p^{-3/2}(\log n)^{19/4}$. Hence, given S_1 , for any choice of S_2 , the number of potential choices for B such that $B + B = G \setminus (S_1 \cup S_2)$ is at most

$$\sum_{i \leq \xi p^{-3/2}(\log n)^{19/4}} \binom{|G|}{i} \leq \exp\left(\xi p^{-3/2}(\log n)^{23/4}\right).$$

On the other hand, the number of choices for S_2 is at least $\exp(\Omega(pn \log n))$. Hence, for $p = \xi' n^{-2/5}(\log n)^{19/10}$ for a sufficiently large constant ξ' , there must exist a choice of S_2 so that no B satisfies $B + B = G \setminus (S_1 \cup S_2)$. \square

4 Arithmetic progressions in sumsets of random sets

In this section, we study $A + A$ for a random subset A of \mathbb{Z}_p where each element is included independently with probability $q = 1/\sqrt{p}$ and prove Theorem 8 regarding the length of the longest arithmetic progression contained in $A + A$. The proof is described in the following two subsections. The upper bound is proved by applying Talagrand's Inequality (with a twist). The lower bound is established by a second moment argument.

4.1 The upper bound

In this subsection we prove the upper bound in Theorem 8. Put $k = c \log p$, where c is a constant, to be chosen later.

Claim 1: With high probability, the set A does not contain 5 elements inside any progression of length k .

We apply a simple union bound. There are less than p^2 arithmetic progressions b_1, b_2, \dots, b_k . In each progression there are less than k^5 possibilities to choose 5 elements, and the probability all of them belong to A is $q^5 = 1/p^{2.5}$. By the union bound the probability that A contains 5 elements inside a progression of length k is thus at most $p^2 k^5 / p^{2.5}$ which is $O(\frac{\log^5 p}{p^{1/2}}) = o(1)$, proving the claim

Note that if the assertion of Claim 1 holds then $A + A$ does not contain 5 elements in a progression as above, with all of them being the sum of a single element a of A with 5 elements of A (possibly including a itself).

Claim 2: For the right choice of the constant c , whp there is no progression of $k = c \log p$ terms, so that each of its terms is a sum of two elements of A , where no element of A is used more than 4 times in these k sums.

Note that if the assertions of both claims hold, then A contains no k -term Arithmetic Progression.

To prove Claim 2 we fix a progression P of k terms and use Talagrand's Inequality to bound the probability of getting all elements of P in $A + A$ by sums of the required form (that is, sums where no element of A is used more than 4 times). We show that this probability is much smaller than $1/p^2$ and then conclude that Claim 2 holds by the union bound.

Let X be the random variable which is the maximum cardinality of a subset T of P that can be expressed by sums of pairs of elements of A , with no element used more than 4 times. Clearly this is a 4-Lipschitz function. It is not difficult to check that the expectation of X is at most $0.4k$ (even without the extra constraint that no $a \in A$ is used more than 4 times.) Indeed, for each fixed element $g \in Z_p$, the probability that g is not in $A + A$ is precisely

$$(1 - q)(1 - q^2)^{(p-1)/2} = (1 + o(1))1/\sqrt{e} > 0.6$$

Therefore the probability that $g \in A + A$ is smaller than 0.4 and the desired estimate follows by linearity of expectation.

Another simple fact is that for any integer t , if $X \geq t$, then this can be certified by at most $2t$ coordinates of the random vector of the p random choices determining which $g \in Z_p$ belong to A .

Therefore, Talagrand's inequality (see, for example, [5], Section 7.7) implies that for every b and t

$$\text{Prob}[X \leq b - 4t\sqrt{2b}] \cdot \text{Prob}[X \geq b] \leq e^{-t^2/4}. \quad (3)$$

Taking in the inequality above b as the median of X , this implies that the probability that X is much smaller than the median is very small. Similarly, taking b and t so that $b - 4t\sqrt{2b}$ is the median shows that the probability that X is much larger than the median is also very small. This implies that the median is close to the expectation, thus it is at most $(0.4 + o(1))k < k/2$. We can now substitute $b = k$ and choose t so that $k - 4t\sqrt{2k}$ is the median. This gives $t = \Theta(\sqrt{k})$. Thus (3) provides an upper bound for the probability that $X = k$. This upper bound is of the form $e^{-\Theta(k)} = e^{-\Theta(c \log p)}$, which is (much) less than $1/p^2$ for an appropriate choice of c . This completes the proof of Claim 2 and hence that of the upper bound in Theorem 8.

4.2 The lower bound

In the proof of the lower bound in Theorem 8 it will be convenient to restrict attention only to arithmetic progressions of length k in which each term is a sum of two distinct elements of A , all the relevant $2k$ elements of A are pairwise distinct, and moreover all the $\binom{2k}{2}$ sums of pairs of these $2k$ elements are distinct.

For every fixed arithmetic progression $P = (b_1, b_2, \dots, b_k)$ of length k in Z_p , where the difference is smaller than $p/2$, and for every ordered sequence $C = (\{c_1, c_2\}, \dots, \{c_{2k-1}, c_{2k}\})$

of k unordered pairs of elements of Z_p satisfying $c_{2i-1} + c_{2i} = b_i$ for all $1 \leq i \leq k$, where all the elements c_i are distinct and all pairs of their sums are distinct, let $X(P, C)$ be the indicator random variable whose value is 1 iff all elements c_j belong to the random set A . Note that if this random variable is 1 then $A + A$ contains the arithmetic progression P .

It is not difficult to check that the number of these indicator random variables is

$$(1 - o(1)) \frac{p^2}{2} \left(\frac{p-1}{2} \right)^k$$

Indeed, there are $p(p-1)/2$ progressions $P = (b_1, b_2, \dots, b_k)$ (obtained by choosing b_1, b_2). For each i there are $(p-1)/2$ unordered pairs of distinct elements c_{2i-1}, c_{2i} whose sum is b_i . The number of choices in which not all the elements c_j are distinct or two sums of pairs of them are identical is negligible with respect to the total number of choices.

Let X be the sum of all these indicator random variables. Note that if $X > 0$ then $A + A$ contains an arithmetic progression of length k .

For a set C of k pairs as above let us denote by \overline{C} the subset of size $2k$ of Z_p consisting of all elements in the union of the k pairs in C . Note that for each fixed indicator $X(P, C)$ as above, the probability it is 1 is exactly q^{2k} , since it is 1 if and only if all the $2k$ elements in \overline{C} lie in A . Therefore, by linearity of expectation and using the fact that $k = \Theta(\log p)$ and $q = 1/\sqrt{p}$, the expected value $E(X) = \mu$ of X satisfies

$$E(X) = \mu = (1 - o(1)) \frac{p^2}{2} \left(\frac{p-1}{2} \right)^k q^{2k} = (1 - o(1)) \frac{p^2}{2^{k+1}}.$$

Our objective is to show that for $k = b \log p$ with the right choice of the constant b , the variance of this random variable is $o(E(X)^2)$. This will show that $X > 0$ whp. Since $E(X) = \mu$ tends to infinity as p tends to infinity, it suffices to prove, using the second moment method as described, for example, in [5], Chapter 4, that $\Delta = o(\mu^2)$, where Δ is defined as follows. It is the sum, over all pairs $X(P, C)$ and $X(P', C')$ that are not independent, of the probability that

$$X(P, C) = X(P', C') = 1.$$

We proceed to upper bound Δ . Fixing an indicator random variable $X(P, C)$ we bound the sum of the conditional probabilities $\sum \text{Prob}[X(P', C') = 1 \mid X(P, C) = 1]$ where $X(P, C)$ is fixed and $X(P', C')$ ranges over all the indicators that are not independent of $X(P, C)$.

Note, first, that if $\overline{C}' \cap \overline{C} = \emptyset$ then the above indicator random variables are independent, hence we may and will assume that $\overline{C} \cap \overline{C}'$ is nonempty. Put

$$C = (\{c_1, c_2\}, \dots, \{c_{2k-1}, c_{2k}\}) \quad \text{and} \quad C' = (\{c'_1, c'_2\}, \dots, \{c'_{2k-1}, c'_{2k}\}).$$

Let ℓ be the number of pairs c'_{2j-1}, c'_{2j} in which both elements belong to \overline{C} . Similarly, let m be the number of additional indices j so that c'_j belongs to \overline{C} . Therefore $|\overline{C} \cap \overline{C}'| = 2\ell + m$. We consider three possible cases, as follows.

Case 1: $\ell \geq 2$.

There are $\binom{k}{\ell}$ ways to choose the indices j for which both elements of the pair c'_{2j-1}, c'_j belong to \overline{C} . For each such choice, there are less than $(2k)^4$ possibilities to choose the specific items of \overline{C} which are equal to the items in the first two pairs above. Given these choices, and given C , we can compute the values of the corresponding two terms of the progression P' , and hence get all other terms of the progression. Now each additional term of the progression P' which is the sum of two elements of A corresponding to another pair of indices of C' that lie in \overline{C} , has to be a sum of two distinct elements c_j . There is at most one way to choose these two c_j (since all sums of pairs c_j are distinct, by assumption). Next, there are at most $(2k)^{2m}$ ways to choose the additional ordered set of m identical elements in \overline{C} and in \overline{C}' and decide about the bijection between them. It remains to choose the additional $2k - 2\ell - m$ elements of \overline{C}' that have not been determined yet. Note that at this point the progression P' is determined, and $2\ell + m$ of the values c'_j are determined as well. This also determines uniquely the values of additional m elements c'_j , since we know the sum of each of the m pairs of elements in which one summand is known already. There are still $2k - 2\ell - 2m$ yet undetermined values c'_j that are partitioned into $k - \ell - m$ pairs, where in each of these pairs the sum of elements is known. This gives $(\frac{p-1}{2})^{k-\ell-m}$ ways to choose the remaining pairs.

Summarizing, we have seen that there are at most

$$\binom{k}{\ell} (2k)^4 (2k)^{2m} \left(\frac{p-1}{2} \right)^{k-\ell-m}$$

indicators $X(P', C')$ corresponding to Case 1 with parameters ℓ and m , where, conditioned on $X(P, C) = 1$ it is still possible that $X(P', C') = 1$. This last event happens if and only if the $2k - 2\ell - m$ required elements c'_j that are not in \overline{C} lie in A . This probability is $\frac{1}{p^{(2k-2\ell-m)/2}}$. It follows that the total contribution of the sum $\sum \text{Prob}[X(P', C') = 1 \mid X(P, C) = 1]$ for a fixed $X(P, C)$, that correspond to Case 1 over all choices of the parameters $\ell \geq 2$ and $m \geq 0$ is bounded by

$$\sum_{\ell \geq 2, m \geq 0} \binom{k}{\ell} (2k)^4 (2k)^{2m} \left(\frac{p-1}{2} \right)^{k-\ell-m} \frac{1}{p^{k-\ell-m/2}} < 2^k (2k)^4 \sum_{m \geq 0} \left[\frac{(2k)^4}{p} \right]^{m/2} \leq 32k^4 2^k.$$

Multiplying by the probability that $X(P, C) = 1$ and summing over all our indicators, we get a bound for the total contribution to Δ arising from pairs that correspond to Case 1. This total contribution is bounded by

$$\Delta_1 \leq \mu \cdot 32k^4 2^k.$$

Case 2: $\ell = 1$.

The discussion here is similar, so we only describe it briefly. There are k possibilities to choose the index j of the pair c'_{2j-1}, c'_{2j} of indices contained in \overline{C} and less than $(2k)^2$ ways to choose the corresponding indices in C . This determines one term of the progression P' , so there are less than p ways to choose the whole progression. Next, there are at most $(2k)^{2m}$ ways to choose the additional elements of \overline{C}' that lie in \overline{C} and the indices of them in C . Then there are at most $(\frac{p-1}{2})^{k-1-m}$ to select the additional elements of C' . This gives a total of at most

$$k(2k)^2 p (2k)^{2m} \left(\frac{p-1}{2} \right)^{k-1-m}$$

relevant indicators $X(P', C')$. The conditional probability of each of them to be 1 assuming that $X(P, C) = 1$ is $\frac{1}{p^{(2k-2-m)/2}}$. The total contribution from these terms to Δ is therefore at most

$$\Delta_2 \leq \mu 4k^3 p \cdot \sum_{m \geq 0} (2k)^{2m} \frac{1}{p^{m/2}} = \mu (4k^3) p \sum_{m \geq 0} \left[\frac{(2k)^4}{p} \right]^{m/2} < \mu \cdot 8k^3 p.$$

Case 3: $\ell = 0$.

Following the same reasoning as before the contribution to Δ from this case is bounded by

$$\Delta_3 \leq \mu \cdot \sum_{m \geq 1} (2k)^{2m} \left(\frac{p-1}{2} \right)^{k-m} \frac{1}{p^{(2k-m)/2}} < \mu \sum_{m \geq 1} \left[\frac{(2k)^4}{p} \right]^{m/2} = \mu \cdot o(1).$$

Summing Δ_1, Δ_2 and Δ_3 we conclude that

$$\Delta \leq \mu [32 \cdot 2^k k^4 + 8k^3 p + o(1)].$$

Since $\mu = (1 + o(1)) \frac{p^2}{2^{k+1}}$ it follows that if, say, $k = 0.999 \log_2 p$, then $\Delta = o(\mu^2)$ and then $X > 0$ whp. This completes the proof of the lower bound, and the assertion of Theorem 8 follows.

5 Concluding remarks

The main covering lemma shows that sumsets of sets with small doubling contain dense subsets which have low complexity. We believe that this result will have further potential applications beyond those discussed in the present paper. In particular:

- In follow-up work, leveraging the main covering lemma and the combinatorial approach to sets with small doubling [7], we derive sharp asymptotics for the independence number of random Cayley graphs in \mathbb{Z}_n with density $(\log n)^{-2+o(1)}$. This improves and surpasses barriers in earlier works [15, 6, 20].

- Combining the probabilistic perspective on sumsets of sets with small doubling with further inputs, we will address a question of the first author, Balogh, Morris, and Samotij [3] about accurate estimates of the number of sets with small doubling and the typical structure of sets with small doubling.

An optimal dependence on the doubling K in the main covering lemma, Theorem 6, would be very interesting. In particular, the following conjecture, if true, would give optimal obstructions that characterize the independence number of sparse random Cayley graphs up to logarithmic factors.

Conjecture 15. *There exist collections of sets \mathcal{F}_ℓ such that*

$$\log |\mathcal{F}_\ell| = \tilde{O}(2^\ell),$$

and

$$\min_{F \in \mathcal{F}_\ell} |F| = \tilde{\Omega}(2^\ell s),$$

such that for every $|A| = s$ with $|A + A| \leq K|A|$, there exists $\ell \leq \log_2 K$ and $F \in \mathcal{F}_\ell$ such that $A + A \supseteq F$.

We believe that it is also interesting to study further properties of the low-complexity subsets of sumsets arising in Theorem 6. In particular, this would address a version of Lovett's question appearing in Green's list of open problems [14], which asks if sumsets $A + A$ of dense sets contain large subsets which are iterated sumsets $B + B + B + B$.

Regarding arithmetic progressions in random sumsets $A + A$:

- For random subsets $A \subset \mathbb{Z}_p$ obtained by picking each element with probability $q = 1/\sqrt{p}$ it is also possible to prove the upper bound established here using the first moment method, where one can show that the main contribution to the expectation is given by k -progressions expressed by sums of $2k$ distinct terms. Indeed, to bound the contribution of the other possibilities one can consider the connected components of the graph whose edges represent the pairs that sum to the required elements. This gives an upper bound of $(2 + o(1)) \log_2 p$. For random sets obtained with larger probability $q = C\sqrt{p}$, where $C > 1$, this first moment argument does not seem to provide any nontrivial upper bound, but the argument described here using Talagrand's Inequality does provide a bound of some $B(C) \log p$ where $B(C)$ is a finite constant for every fixed C .
- Let $A = A(q) \subset \mathbb{Z}_p$ be a random subset of \mathbb{Z}_p obtained by picking each element of \mathbb{Z}_p , randomly and independently, with probability $q = q(p)$. Combining the methods here with the Brun Sieve it may be possible to determine the typical asymptotic behavior of $ap(A + A)$ up to a constant factor, or even up to a $(1 + o(1))$ factor, for every possible $q = q(p)$. For sufficiently small values of q this is essentially done (using a different method) in [18]. We hope to return to this problem in the future.

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