Excluding the fork and antifork

Maria Chudnovsky\textsuperscript{1}, Linda Cook, and Paul Seymour\textsuperscript{2}
Princeton University, Princeton, NJ 08544

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Abstract

The fork is the tree obtained from the claw $K_{1,3}$ by subdividing one of its edges once, and the antifork is its complement graph. We give a complete description of all graphs that do not contain the fork or antifork as induced subgraphs.
1 Introduction

Graphs in this paper are finite, and without loops or parallel edges. The fork is the tree obtained from a four-vertex path by adding a vertex adjacent to the second vertex of the path. The antifork is its complement graph; thus, the antifork is obtained from a four-vertex path by adding one more vertex adjacent to the first three vertices of the path. (See figure 1.)

Figure 1: The fork and the antifork.

Let us say $G$ is uncluttered if no induced subgraph of $G$ is a fork or antifork. Our goal in this paper is to give a complete description of all uncluttered graphs. The line graph of a triangle-free graph is uncluttered, and so is its complement; and our main theorem says that every uncluttered graph can be obtained by piecing together line graphs of triangle-free graphs and their complements. Before we can make a precise statement we need to explain the “piecing together” process, and that is the content of the next section. We state our main result in 2.1.

This question was motivated by discussions with T. Karthick, who told us several results about colouring graphs not containing forks and/or antiforks [1, 2, 3] (note, however, that in [2] “fork” has a different meaning from its meaning here). In particular, he asked for the best possible “$\chi$-bounding function” for uncluttered graphs. The answer follows from our main result 2.1 (easily, by induction, and we omit the details):

1.1 For every uncluttered graph, its chromatic number is at most twice its clique number.

This is asymptotically best possible, since if $H$ is a triangle-free graph with largest stable set of cardinality $k$, then the complement of the line graph of $H$ is uncluttered, with clique number at most $|V(H)|/2$ and with chromatic number $|V(H)| - k$; and we can choose $H$ and $k$ with $(|V(H)| - k)/|V(H)|$ arbitrarily close to 1.

2 Some safe operations

There are several ways to make larger uncluttered graphs from smaller ones. The most obvious is: if $G_1, G_2$ are both uncluttered, then so is their disjoint union. The complete join of $G_1, G_2$ is obtained from their disjoint union by adding edges between every vertex of $G_1$ and every vertex of $G_2$. Since the complement of an uncluttered graph is also uncluttered, and the complete join of $G_1, G_2$ is the complement of the disjoint union of $\overline{G_1}, \overline{G_2}$, it follows that if $G_1, G_2$ are both uncluttered, then so is their complete join. Let us say $G$ is anticonnected if its complement graph $\overline{G}$ is connected.

If $X, Y \subseteq V(G)$ are disjoint, we say $X$ is complete to $Y$ if every vertex in $X$ is adjacent to every vertex in $Y$; and $X$ is anticomplete to $Y$ if there are no edges between $X$ and $Y$. If $v \in V(G)$, we say “$v$ is complete to $Y$” meaning that $\{v\}$ is complete to $Y$, and so on.
Two (distinct) vertices $u, v$ of $G$ are twins if $u, v$ have the same neighbours in $V(G) \setminus \{u, v\}$; $(u, v$ may or may not be adjacent). A homogeneous set in $G$ means a set $X \subseteq V(G)$ such that every vertex of $G$ not in $X$ is either complete or anticomplete to $X$; and the homogeneous set is nontrivial if $|X| \geq 2$ and $X \neq V(G)$. If $X$ is a homogeneous set in $G$, let $Y$ be the set of vertices in $V(G) \setminus X$ that are complete to $X$. Take another vertex $v$, and let $H$ be the graph formed from $G \setminus X$ by adding the vertex $v$ and making $v$ adjacent to the vertices in $Y$. Then we say $G$ is obtained from $H$ by substituting $G[X]$ for $v$.

Say a vertex $v$ is simplicial if the set of its neighbours is a clique. If $v$ is a simplicial vertex in an uncluttered graph $G$, then we may substitute a complete graph for $v$, and the new graph we obtain is also uncluttered. Consequently, if $G$ has adjacent simplicial twins, then $G$ can be obtained from a smaller graph by the operation just described. A vertex $v$ is antisimplicial if it is simplicial in the complement graph, that is, if the set of all vertices nonadjacent to $v$ is a stable set.

Let $G$ be a graph, let $k \geq 1$ be an integer, and let $Y_1, \ldots, Y_k, Z_1, \ldots, Z_k$ be disjoint nonempty sets of $V(G)$ with union $V(G)$, such that

- $Y_1, \ldots, Y_k$ are cliques, and $Z_1, \ldots, Z_k$ are stable sets;
- for $1 \leq i < j \leq k$, $Y_i$ is anticomplete to $Y_j$, and $Z_i$ is complete to $Z_j$;
- for $1 \leq i, j \leq k$, $Y_i$ is complete to $Z_j$ if $i = j$, and otherwise $Y_i$ is anticomplete to $Z_j$.

We call such a graph $G$ a candelabrum, with base $Z_1 \cup \cdots \cup Z_k$. (See figure 2.)

Figure 2: A candelabrum. (Lines indicate complete pairs, and $K, S$ mark cliques and stable sets.)

Candelabra are uncluttered, but they can also be used to make larger uncluttered graphs. Let $H_1$ be an uncluttered graph, and let $H_2$ be a candelabrum with base $Z$. Take the disjoint union of $H_1, H_2$, and add edges to make $V(H_1)$ complete to $Z$. Let $G$ be the graph we produce. Then $G$ is uncluttered (we leave checking this to the reader), and $V(H_1)$ is a homogeneous set of $G$. We say $G$ is candled if it can be constructed by this process; that is, if $G$ has an induced subgraph $H$ that is a candelabrum, with base $Z$ say, and $Z$ is complete to $V(G) \setminus V(H)$, and $V(H) \setminus Z$ is anticomplete to $V(G) \setminus V(H)$. (Thus $V(G) \setminus V(H)$ is a homogeneous set, but it might not be nontrivial; it might even be empty.)

Now we can state our main theorem:

2.1 Let $G$ be an uncluttered graph. Then either
• one of $G, \overline{G}$ is disconnected; or
• one of $G, \overline{G}$ has adjacent simplicial twins; or
• one of $G, \overline{G}$ is candled; or
• one of $G, \overline{G}$ is the line graph of a triangle-free graph.

The proof of 2.1 will occupy the remainder of the paper, and will be completed in 4.7.

3 Homogeneous sets

In this section we will prove:

3.1 Let $G$ be an uncluttered graph. Suppose that

• $G$ is connected and anticonnected;
• $G$ has no adjacent simplicial twins, and no nonadjacent antisimplicial twins; and
• $G, \overline{G}$ are not candled.

Then $G$ has no nontrivial homogeneous set.

The proof requires several steps, that we carry out in this section. If $X \subseteq V(G)$ we denote by $G[X]$ the subgraph induced on $X$. Let us say an anticomponent of a graph $G$ is a induced subgraph whose complement graph is a component of $\overline{G}$. We begin with:

3.2 Let $G$ be uncluttered, and connected and anticonnected, with a nontrivial homogeneous set that is not a clique or stable set. Then one of $G, \overline{G}$ is candled.

Proof. Let $X$ be a nontrivial homogeneous set that is not a clique or stable set. By replacing $X$ by a superset if necessary, we may assume that no proper superset of $X$ is a nontrivial homogeneous set. Let $Z$ be the set of vertices in $V(G) \setminus X$ that are complete to $X$, and let $Y$ be the set of vertices in $V(G) \setminus X$ that are anticomplete to $X$. Let $Y_1, \ldots, Y_k$ be the vertex sets of the components of $G[Y]$, and let $Z_1, \ldots, Z_\ell$ be the vertex sets of the anticomponents of $G[Z]$.

(1) $Y_1, \ldots, Y_k$ and $Z_1, \ldots, Z_\ell$ are homogeneous sets of $G$.

Suppose that $Y_1$ is not a homogeneous set; then there exists $v \in V(G) \setminus Y_1$ that has a neighbour and a non-neighbour in $Y_1$, and since $G[Y_1]$ is connected, there is an edge $yy'$ of $G[Y_1]$ such that $v$ is adjacent to $y$ and not to $y'$. Since $X$ is anticomplete to $Y_1$, it follows that $v \notin X$, and similarly $v \notin Y_2, \ldots, Y_k$; so $v \in Z$. Now $X$ is not a clique, so there exist nonadjacent $x_1, x_2 \in X$. But then $G[\{x_1, x_2, v, y, y'\}]$ is a fork, a contradiction. Thus $Y_1$ is a homogeneous set, and similarly so are $Y_2, \ldots, Y_k$; and so are $Z_1, \ldots, Z_\ell$, by applying the same argument in the complement graph. This proves (1).

(2) $Y, Z$ are both nonempty. Moreover, for $1 \leq i \leq k$, there exist $j, j' \in \{1, \ldots, \ell\}$ such that $Y_i$ is complete to $Z_j$ and anticomplete to $Z_{j'}$, and for $1 \leq j \leq \ell$, there exist $i, i' \in \{1, \ldots, k\}$ such that
Since X is a nontrivial homogeneous set, X \neq V(G) and so Y \cup Z \neq \emptyset. Since G is connected, it follows that Y \neq \emptyset, and similarly Z \neq \emptyset since G is anticonnected. Since G is connected, for 1 \leq i \leq k there exists j \in \{1, \ldots, \ell\} such that Y_i is complete to Z_j; and also, since X \cup \{Y_i\} is not a homogeneous set from the maximality of X, it follows that there exists j' \in \{1, \ldots, \ell\} such that Y_i is anticomplete to Z_{j'}.

The same argument in the complement shows the final statement. This proves (2).

Let i_1, \ldots, i_t \in \{1, \ldots, k\} be distinct, and let j_1, \ldots, j_t \in \{1, \ldots, \ell\} be distinct. We say the pairs (i_1, j_1), \ldots, (i_t, j_t) form a matching of order t if for 1 \leq r, s \leq t, Y_{i_r} is complete to Z_{j_s} if r = s and otherwise Y_{i_r} is anticomplete to Z_{j_s}. Similarly the pairs form an antimatching of order t if for 1 \leq r, s \leq t, Y_{i_r} is anticomplete to Z_{j_s} if r = s and otherwise Y_{i_r} is complete to Z_{j_s}.

(3) There exist pairs (i, j), (i', j') that form a matching of order two.

Choose i \in \{1, \ldots, k\} such that Y_i is complete to Z_j for as many j \in \{1, \ldots, \ell\} as possible; by (2), we can choose j' \in \{1, \ldots, \ell\} such that Y_i is anticomplete to Z_{j'}; by (2) we can choose i' \in \{1, \ldots, k\} such that Y_{i'} is complete to Z_{j'}; and then from the choice of i, there exists j \in \{1, \ldots, \ell\} such that Z_j is complete to Y_i and not to Y_{i'}. Then (i, j), (i', j') forms a matching of order two. This proves (3).

Choose t \geq 2 maximum such that there are t pairs (i, j) (with i \in \{1, \ldots, k\} and j \in \{1, \ldots, \ell\}) that form a matching or antimatching; and by taking complements if necessary, we may assume the pairs form a matching. By renumbering, we may assume that (1,1), \ldots, (t, t) form a matching. For 1 \leq i \leq t, choose y_i \in Y_i and z_i \in Z_i.

(4) Every vertex in Y_{t+1} \cup \cdots \cup Y_k is complete or anticomplete to Z_1 \cup \cdots \cup Z_t, and every vertex in Z_{t+1} \cup \cdots \cup Z_\ell is complete or anticomplete to Y_1 \cup \cdots \cup Y_t.

If v \in Y_{t+1} \cup \cdots \cup Y_k has a neighbour in Z_1 and a nonneighbour in Z_2 say, then G[\{v, z_1, z_2, y_1, y_2\}] is a fork, a contradiction. Similarly, if v \in Z_{t+1} \cup \cdots \cup Z_\ell has a neighbour in Y_1 and a nonneighbour in Y_2, then G[\{v, z_1, z_2, y_1, y_2\}] is an antifork, a contradiction. This proves (4).

Let P be the set of v \in Y_{t+1} \cup \cdots \cup Y_k such that v is complete to Z_1 \cup \cdots \cup Z_t, and let P' be the set of v \in Y_{t+1} \cup \cdots \cup Y_k such that v is anticomplete to Z_1 \cup \cdots \cup Z_t. Let Q be the set of v \in Z_{t+1} \cup \cdots \cup Z_\ell such that v is complete to Y_1 \cup \cdots \cup Y_t, and let Q' be the set of v \in Z_{t+1} \cup \cdots \cup Z_\ell such that v is anticomplete to Y_1 \cup \cdots \cup Y_t.

(5) X \cup Y_1 \cup \cdots \cup Y_t \cup Z_1 \cup \cdots \cup Z_t \cup P \cup Q' is a homogeneous set.

Let us call this set A. We will show that P' is anticomplete to A, and Q is complete to A. Let v \in P'; then v is anticomplete to X since v \in Y; v is anticomplete to Y_1 \cup \cdots \cup Y_t since Y_1, \ldots, Y_t are vertex sets of components of G[Y_i]; and v is anticomplete to Z_1 \cup \cdots \cup Z_t by definition of P'. Since v belongs to some Y_i, and Y_i is a homogeneous set by (1), and so Y_i \subseteq P', it follows that v is
anticomplete to $P$. Also from the maximality of $t$, $v$ is anticomplete to $Q'$. This proves that $P'$ is anticomplete to $A$.

Now let $v \in Q$. We must show that $v$ is complete to $A$. Certainly $v$ is complete to $X$ and to $Z_1 \cup \cdots \cup Z_t$, and to $Y_1 \cup \cdots \cup Y_t$ from the definition of $Q$. Since $v$ belongs to an anticomponent of $G[Z]$ with vertex set in $Q$, it follows that $v$ is complete to $Q'$. It remains to show that $v$ is complete to $P$. Suppose not, and let $u \in P$ be nonadjacent to $v$. For $1 \leq i \leq t$, choose $y_i \in Y_i$ and choose $z_i \in Z_i$. If $t \geq 3$, $G\{v, y_1, y_2, u, z_3\}$ is a fork, a contradiction, and so $t = 2$. Let $u \in Y_2$ and $v \in Z_3$ say; then $Y_3$ is anticomplete to $Z_3$ by (1), and so the three pairs $(1, 2), (2, 1), (3, 3)$ form an antimatching, contrary to the maximality of $t$. This proves (5).

From the maximality of $X$, (5) implies that the set of (5) is not a nontrivial homogeneous set, and so $P', Q = \emptyset$. But then $X \cup P \cup Q'$ is a nontrivial homogeneous set, and the maximality of $X$ implies that $P, Q' = \emptyset$. If there exist nonadjacent $y, y' \in Y_1$, then $G\{y, y', z_1, z_2, y_2\}$ is a fork, a contradiction; so $Y_1, \ldots, Y_t$ are cliques. If there exist adjacent $z, z' \in Z_1$, then $G\{z, z', y_1, z_2, y_2\}$ is an antifork, a contradiction. Thus $Z_1, \ldots, Z_t$ are stable sets. But then $G$ is candled, as required. This proves 3.2.

Let $P$ be a four-vertex induced path in a graph $G$. A centre for $P$ means a vertex of $V(G) \setminus V(P)$ that is complete to $V(P)$, and an anticentre for $P$ is a vertex of $V(G) \setminus V(P)$ that is anticomplete to $V(P)$.

3.3 Let $G$ be uncluttered, and let $P$ be a four-vertex induced path in $G$. If there is a centre and an anticentre for $P$ then there is a nontrivial homogeneous set in $G$ that is not a clique or stable set.

**Proof.** Let $A$ be the set of all centres for $P$, and let $C$ be the set of all centres for $P$. Thus $A, C \neq \emptyset$. Let $B$ be the set of $v \in V(G) \setminus (V(P) \cup A \cup C)$; thus $v$ is the set of all vertices not in $V(P)$ with a neighbour and a nonneighbour in $V(P)$. Now either every vertex in $C$ has a neighbour in $A$, or every vertex in $A$ has a nonneighbour in $C$; and by taking complements if necessary, we assume the first. Let $P$ have vertices $p_1-p_2-p_3-p_4$ in order.

(1) $B$ is complete to $C$.

Let $b \in B$ and $c \in C$, and suppose that $b, c$ are nonadjacent. Choose $a \in A$ adjacent to $c$. Suppose first that $a, b$ are nonadjacent. Since $P$ is anticonnected, and $b \in B$, there exist nonadjacent $p, p' \in V(P)$ such that $b$ is adjacent to $p$ and not to $p'$; but then $G\{p, p', a, b, c\}$ is a fork, a contradiction. So $a, b$ are adjacent. Let $I$ be the set of $i \in \{1, \ldots, 4\}$ such that $b, p_i$ are adjacent. Since $G\{a, b, c, p_1, p_4\}$ is not a fork, one of $1, 4 \in I$, and we may assume $1 \in I$ without loss of generality. If $4 \notin I$, then $2 \in I$ since $G\{a, b, c, p_2, p_4\}$ is not a fork; so $3 \notin I$ since the subgraph induced on $G\{b, p_1, p_2, p_3, p_4\}$ is not an antifork; but then $G\{b, c, p_1, p_2, p_4\}$ is an antifork, a contradiction. Thus $4 \in I$. Since $b$ is not a centre, one of $2, 3 \notin I$, and we assume $3 \notin I$ without loss of generality. But then $G\{a, b, p_1, p_3, p_4\}$ is a fork, a contradiction. This proves (1).

Let $A'$ be the union of the vertex sets of the components of $G[A]$ that are not anticomplete to $B$, and let $A'' = A \setminus A'$.

(2) $A'$ is complete to $C$.
Let \( c \in C \), and suppose that \( c \) is not complete to \( A' \). From the definition of \( A' \), there is an induced path \( Q \) with one end in \( B \), and all other vertices in \( A' \), such that some vertex of \( Q \) is not adjacent to \( c \). Choose \( Q \) minimal, with vertices \( q_1, \ldots, q_k \) in order, where \( q_k \in B \) and \( q_1, \ldots, q_{k-1} \in A \). From (1), \( k \geq 2 \). From the minimality of \( Q \), it follows that \( c \) is nonadjacent to \( q_1 \) and adjacent to all of \( q_2, \ldots, q_k \). Choose adjacent \( q_{k+1}, q_{k+2} \) of \( P \) such that \( q_k \) is adjacent to \( q_{k+1} \) and not to \( q_{k+2} \). (This is possible since \( q_k \in B \).) Then \( G[\{c, q_1, q_2, q_3, q_4\}] \) is an antifork, a contradiction. This proves (2).

Let \( X = V(P) \cup B \cup A' \). From (1) and (2), \( C \) is complete to \( X \); and from the definition of \( A'', A'' \) is anticomplete to \( X \). Thus \( X \) is a nontrivial homogeneous set satisfying the theorem. This proves 3.3.

3.4 Let \( G \) be uncluttered and connected, and let \( A, B, C, D \) be disjoint subsets of \( V(G) \), with union \( V(G) \), and with the following properties:

- \( A \) is a clique and \( A \neq \emptyset \);
- \( B \) is a stable set and \(|B| \geq 2\);
- \( A \) is complete to \( B \) and anticomplete to \( C, D \);
- \( B \) is complete to \( C \) and anticomplete to \( D \).

Then \( G \) is candled.

**Proof.** Let \( C_1, \ldots, C_k \) be the vertex sets of the anticomponents of \( G[C] \), and let \( D_1, \ldots, D_\ell \) be the vertex sets of the components of \( G[D] \).

(1) \( C_1, \ldots, C_k \) and \( D_1, \ldots, D_\ell \) are homogeneous sets.

Suppose \( C_1 \) is not a homogeneous set; then there exist \( d \in V(G) \setminus C_1 \) and nonadjacent \( c, c' \in C_1 \) with \( d \) adjacent to \( c \) and not to \( c' \). Choose \( a \in A \) and \( b \in B \); then \( G[\{a, b, c, c', d\}] \) is a fork, a contradiction. Thus \( C_1, \ldots, C_k \) are all homogeneous sets.

Now suppose \( D_1 \) is not a homogeneous set; then similarly there exists \( c \in C \) and adjacent \( d, d' \in D_1 \) such that \( c \) is adjacent to \( d \) and not to \( d' \). Since \(|B| \geq 2 \) and \( B \) is stable, there exist nonadjacent \( b, b' \in B \); but then \( G[\{b, b', c, d, d'\}] \) is a fork, a contradiction. This proves (1).

For \( 1 \leq i \leq k \) choose \( c_i \in C_i \), and for \( 1 \leq j \leq \ell \) choose \( d_j \in D_j \). Choose \( a \in A \) and \( b \in B \).

(2) For \( 1 \leq j \leq \ell \), there is a unique value of \( i \in \{1, \ldots, k\} \) such that \( D_j \) is complete to \( C_i \).

We assume \( j = 1 \) without loss of generality. Since \( G \) is connected and \( A \cup B \cup D_2 \cup \cdots \cup D_\ell \) is anticomplete to \( D_1 \), it follows from (1) that there exists \( i \in \{1, \ldots, k\} \) such that \( D_1 \) is complete to \( C_i \). Suppose there are two such values of \( i \), say \( i = 1, 2 \); then \( G[\{a, b, c_1, c_2, d_1\}] \) is an antifork, a contradiction. This proves (2).
(3) For $1 \leq i \leq k$ there is at most one value of $j \in \{1, \ldots, \ell\}$ such that $C_i$ is complete to $D_j$.

For let $i = 1$ say, and suppose that $C_1$ is complete to $D_1, D_2$ say. Then $G[\{a, b, c_1, d_1, d_2\}]$ is a fork, a contradiction. This proves (3).

From (2) and (3) it follows that $k \geq \ell$, and we may renumber such that $D_i$ is complete to $C_i$ for $1 \leq i \leq \ell$. For $1 \leq i \leq k$, $D_i$ is a clique, since if $d, d' \in D_i$ are nonadjacent then $G[\{a, b, c_i, d, d'\}]$ is a fork; and also $C_i$ is a stable set, since if $c, c' \in C_i$ are adjacent then $G[\{a, b, c, c', d_i\}]$ is an antifork.

Thus the restriction of $G$ to $A \cup B \cup C_1 \cup \cdots \cup C_\ell \cup D_1 \cup \cdots \cup D_\ell$

is a candelabrum with base $B \cup C_1 \cup \cdots \cup C_\ell$. Since $C_{\ell+1} \cup \cdots \cup C_k$ is complete to $B \cup C_1 \cup \cdots \cup C_\ell$ and anticomplete to $A \cup D_1 \cup \cdots \cup D_\ell$, it follows $G$ is candled. This proves 3.4.

Now we can prove the main result of this section, which we restate:

3.5 Let $G$ be an uncluttered graph. Suppose that

- $G$ is connected and anticonnected;
- $G$ has no adjacent simplicial twins, and no nonadjacent antisimplicial twins;
- $G, \overline{G}$ are not candled.

Then $G$ has no nontrivial homogeneous set.

Proof. For each nontrivial homogeneous set $X$, we define its “score” as follows. By 3.2, $X$ is either a clique or a stable set. If $X$ is stable, its score is the number of components of $G[Y]$, where $Y$ is the set of all vertices in $V(G) \setminus X$ that are anticomplete to $X$. If $X$ is a clique, its score is the number of anticomponents of $G[Z]$, where $Z$ is the set of vertices in $V(G) \setminus X$ that are complete to $X$.

Suppose that there is a nontrivial homogeneous set $X$, and choose $X$ with minimum score. Let $Z$ be the set of vertices in $V(G) \setminus X$ that are complete to $X$, and $Y$ the set that are anticomplete to $X$. Let $Y_1, \ldots, Y_k$ be the vertex sets of the components of $G[Y]$, and let $Z_1, \ldots, Z_\ell$ be the vertex sets of the anticomponents of $G[Z]$. By replacing $G$ by its complement if necessary, we may assume that $X$ is a stable set.

(1) $Y_1, \ldots, Y_k$ are homogeneous sets, and cliques.

If $Y_1$ is not a homogeneous set, then there exists $z \in Z$ and $y, y' \in Y_1$, such that $z-y-y'$ is an induced path (because $G[Y_1]$ is connected). Let $x, x' \in X$ be nonadjacent; then $G[\{x, x', z, y, y'\}]$ is a fork, a contradiction. This proves that $Y_1$ is a homogeneous set. By 3.2, $Y_1$ is either a stable set or a clique; but $G[Y_1]$ is connected, and so $Y_1$ is a clique. This proves (1).

By hypothesis, there are no nonadjacent antisimplicial twins in $G$; and in particular, vertices in $X$ are not antisimplicial. Thus $Y$ is not stable, and so we may assume that $|Y_1| \geq 2$. Let $N$ be the set of vertices in $Z$ that are complete to $Y_1$. By hypothesis, there are no adjacent simplicial
twins, and in particular the vertices in \( Y_1 \) are not simplicial; so \( N \) is not a clique. Hence there is an anticomponent \( D \) of \( N \) with at least two vertices. Let \( N' = N \setminus D \), and \( Y' = Y_2 \cup \cdots \cup Y_k \). Thus the six sets \( X, D, N', Z \setminus N, Y_1, Y' \) are pairwise disjoint and have union \( V(G) \).

(2) \( D \) is a homogeneous set, and \( D \) is stable.

Since \( G[D] \) is anticonnected, if \( D \) is not a homogeneous set then there exists \( v \in V(G) \setminus D \) and nonadjacent \( d, d' \in D \), such that \( v \) is adjacent to \( d \) and not to \( d' \). Since \( X \cup N' \cup Y_1 \) is complete to \( D \), it follows that \( v \in (Z \setminus N) \cup Y' \). Choose \( y, y' \in Y_1 \), adjacent; then \( G[\{v, d, d', y, y'\}] \) is an antifork, a contradiction. Thus \( D \) is homogeneous. By 3.2 is either either stable or a clique, and it is not a clique since \( G[D] \) is anticonnected and \(|D| \geq 2 \). This proves (2).

Let \( U \) be the set of vertices in \( Z \setminus N \) that are complete to \( D \), and let \( W \) be the set of vertices in \( Z \setminus N \) that are anticomplete to \( D \). Thus \( U \cup W = Z \setminus N \) by (2).

(3) \( W \) is a clique.

If \( w, w' \in W \) are nonadjacent, choose \( y \in Y_1 \), \( d \in D \) and \( x \in X \); then \( G[\{w_1, w_2, y_1, d, x\}] \) is a fork, a contradiction. This proves (3).

(4) \( Y' \) is anticomplete to \( D \cup W \), and \( N' \) is complete to \( W \), and \( U \) is anticomplete to \( W \).

From the minimality of the score of \( X \), it follows that the score of \( D \) is at least that of \( X \), that is, at least \( k \). For \( 2 \leq i \leq k \), \( Y_i \) is a homogeneous set by (1), and since \( D \) is also a homogeneous set by (2), it follows that \( Y_i \) is complete or anticomplete to \( D \). Let \( C \) denote the set of vertices not in \( D \) and anticomplete to \( D \). Then \( C \) consists of \( W \) and some of \( Y_2, \ldots, Y_k \), those \( Y_i \) that are anticomplete to \( D \). In particular, since \( W \) is a clique by (3), the score of \( D \) is one more (for \( W \)) than the number of \( i \in \{2, \ldots, k\} \) such that \( Y_i \) is anticomplete to \( D \cup W \). Since the score of \( D \) is at least \( k \), it follows that \( Y_2, \ldots, Y_k \) are all anticomplete to \( D \cup W \). This proves the first assertion of (4).

For the second assertion of (4), let \( y \in Y_1 \), \( d \in D \), \( x \in X \) and \( w \in W \); then for \( n \in N' \), if \( n' \) is nonadjacent to \( w \) then \( G[\{y, d, x, w, n\}] \) is an antifork. This proves that \( N' \) is complete to \( W \).

For the third assertion, let \( u \in U \), and let \( y, d, x, w \) be as before; if \( u, w \) are adjacent then \( G[\{y, d, x, u, w\}] \) is an antifork, a contradiction. This proves (4).

If \( N' = \emptyset \), then the four sets \( W, X, D \cup U, Y \) (in this order) satisfy the hypotheses of 3.4, a contradiction. Thus \( N' \neq \emptyset \). If \( Y' = \emptyset \), then the four sets \( X, Y_1, W \cup U, N' \cup D \) satisfy the hypotheses of 3.4 applied in \( \overline{G} \), a contradiction. So \( Y' \neq \emptyset \). Choose \( y \in Y_1 \), \( y' \in Y' \), \( d \in D \), \( x \in X \), \( n \in N' \) and \( w \in W \); then \( y-d-x-w \) is a four-vertex induced path, and \( n \) is a centre and \( y' \) is an anticentre, contrary to 3.3. This proves 3.5.

4 Graphs without homogeneous sets

In view of 3.1, henceforth we can restrict our attention to uncluttered graphs with no nontrivial homogeneous set. First here is a useful lemma.
4.1 Let $G$ be a graph with no nontrivial homogeneous set, and let $A \subseteq V(G)$, not a clique, with $A \neq V(G)$. Then there exist $v \in V(G) \setminus A$ and nonadjacent $a, a' \in A$, such that $v$ is adjacent to $a$ and not to $a'$.

Proof. Since $A$ is not a clique, there is an anticomponent $X$ of $G[A]$ with at least two vertices. Since $X \neq V(G)$ and $X$ is not a nontrivial homogeneous set, there is a vertex $v \in V(G) \setminus X$ with a neighbour and a nonneighbour in $X$. Since $X$ is an anticomponent of $G[A]$, it follows that $v \notin A$. Since $G[X]$ is anticonnected, and $v$ has a neighbour and a nonneighbour in $X$, there exist nonadjacent $a, a' \in X$ such that $v$ is adjacent to $a$ and not to $a'$. This proves 4.1.

If $X$ is a subgraph of $G$, or a set of vertices of $G$, we say that $X$ is dominating if every vertex not in $X$ has a neighbour in $X$. The diamond is the graph with four vertices and five edges, and a triangle is a clique of cardinality three. We begin with:

4.2 Let $G$ be an uncluttered graph with no nontrivial homogeneous set. Then every triangle in $G$ is dominating. Moreover, every triangle that is contained in a diamond of $G$ is dominating.

Proof. Suppose that there is a diamond $D$ that is not dominating. Let $v_1, v_2$ be the two vertices of $D$ that have degree three in $D$; so there is a set $A$ of $G$, not a clique, such that $v_1, v_2 \notin A$ and $A$ is complete to $\{v_1, v_2\}$, and $A \cup \{v_1, v_2\}$ is not dominating. Choose such a set $A$, maximal. Since $A$ is not a nontrivial homogeneous set, 4.1 implies that there exist $v \in V(G) \setminus A$ and nonadjacent $a, a' \in A$ such that $v$ is adjacent to $a$ and not to $a'$. Thus $v \neq v_1, v_2$. If $v$ is nonadjacent to both $v_1, v_2$, then $G[\{v, v_1, v_2, a_1, a_2\}]$ is an antifork, a contradiction; so we may assume that $v$ is adjacent to $v_1$. Since $A \cup \{v_1, v_2\}$ is not dominating, there is a vertex $w$ that is anticomplete to $A \cup \{v_1, v_2\}$. Now there are four possibilities: $w$ may or may not be adjacent to $v$, and $v$ may or may not be adjacent to $v_2$. Suppose first that $v$ is not adjacent to $v_2$, and so $v-a-v_2-a'$ is an induced four-vertex path, $P$ say. Now $v_1$ is a centre for $P$, so by 3.3, $w$ is not an anticentre, and therefore $w$ is adjacent to $v$. But then $G[\{w, v, v_1, a, v_2\}]$ is an antifork, a contradiction. So $v$ is adjacent to $v_2$. If $w, v$ are nonadjacent, then we can add $v$ to $A$, contrary to the maximality of $A$. Thus $v, w$ are adjacent; but then $G[\{w, v, v_1, v_2, a'\}]$ is an antifork, a contradiction. This proves that every diamond is dominating.

Now let $T$ be a triangle, contained in a diamond; thus some vertex $v \in V(G) \setminus T$ has exactly two neighbours in $T$. Suppose that $T$ is not dominating, and let $w$ be anticomplete to $T$. Since the diamond $G[T \cup \{v\}]$ is dominating, $w$ is adjacent to $v$; but then $G[T \cup \{v, w\}]$ is an antifork, a contradiction. This proves 4.2.

4.3 Let $G$ be an uncluttered graph with no nontrivial homogeneous set, such that $G$ is not the line graph of a bipartite graph. Then for every nondominating triangle $T$ in $G$, there is a unique maximal clique including $T$, say $C$, and it is not dominating. Moreover, for each $v \in C$, if $C_v$ is the only maximal clique containing $v$, then there is exactly one other maximal clique containing $v$, say $C_v$, and $C_v \cap C = \{v\}$, and $C_v$ is not dominating.

Proof. Let $C$ be a maximal nondominating clique including $T$; and let $S$ be the set of vertices that are anticomplete to $C$. Thus $S \neq \emptyset$. Let $A$ be the set of vertices in $V(G) \setminus C$ with exactly one neighbour in $C$, and let $B$ be the set of vertices in $V(G) \setminus C$ with at least two neighbours in $C$. Thus
$A, B, C, S$ are pairwise disjoint and have union $V(G)$.

(1) $B$ is complete to $C \cup S$.

No triangle included in $C$ is dominating, and hence by 4.2, no triangle included in $C$ is contained in a diamond. Consequently every vertex in $B$ is complete to $C$. From the maximality of $C$, $B$ is complete to $S$. This proves (1).

Every vertex in $A$ has a unique neighbour in $C$. For each $c \in C$, let $A_c$ be the set of vertices in $A$ that are adjacent to $c$. Thus $A = \bigcup_{c \in C} A_c$.

(2) There is at most one $c \in C$ with $A_c = \emptyset$.

For otherwise the set of $c \in C$ with $A_c = \emptyset$ is a nontrivial homogeneous set, which is impossible. This proves (2).

(3) $|B| \leq 1$, and $B$ is anticomplete to $A$.

Suppose first that $B$ is not a clique. By 4.1, there is a vertex $v \in V(G) \setminus B$ and nonadjacent $b, b' \in B'$ such that $v$ is adjacent to $b$ and not to $b'$. By (1), $v \in A_c$ for some $c \in C$. Choose $c_1, c_2 \in C$ different from $c$; then $G[\{v, b, b', c_1, c_2\}]$ is an antifork, a contradiction. This proves that $B$ is a clique.

Suppose that some $b \in B$ has a neighbour in $A$. Since $G$ is anticonnected (because it has no nontrivial homogeneous set), it follows that $b$ has a nonneighbour; and since $b$ is complete to $C \cup S$ and $B$ is a clique, it follows that $b$ has a nonneighbour in $A$. By (2), there are at least two vertices $c \in C$ such that $A_c \neq \emptyset$; and so there exist distinct $c_1, c_2 \in C$ and $a_i \in A_{c_i}$ for $i = 1, 2$, such that $b$ is adjacent to $a_1$ and nonadjacent to $a_2$. (To see this, choose a neighbour $a \in A$ of $b$ and a nonneighbour $a' \in A$ of $b$. If $a, a'$ both belong to $A_{c_1}$ say, choose $a'' \in A_{c_2}$, and replace one of $a, a'$ by $a''$.) Choose $c_3 \in C \setminus \{c_1, c_2\}$. Since $G[\{b, c_1, c_2, a_1\}]$ is a diamond, the triangle $G[\{b, c_1, c_3\}]$ is dominating by 4.2, and yet $a_2$ has no neighbour in this triangle, a contradiction. Thus $B$ is anticomplete to $A$, and so $B$ is a homogeneous set; and hence $|B| \leq 1$. This proves (3).

(4) For each $c \in C$, $A_c$ is a clique; and if $B \neq \emptyset$ then $S$ is a clique.

Suppose that $c \in C$ and $A_c$ is not a clique; then by 4.1, there exists $v \notin A_c$ and nonadjacent $a, a' \in A_c$ such that $v$ is adjacent to $a$ and not to $a'$. Thus $v \in A \cup S$, by (3); and so there exists $c' \in C \setminus \{c\}$ nonadjacent to $v$. But then $G[\{c, c', a, a', v\}]$ is a fork, a contradiction. Thus each $A_c$ is a clique. Now suppose that $B = \{b\}$ say, and $S$ is not a clique. Then by 4.1, there exists $v \notin S$ and nonadjacent $s, s' \in S$ such that $v$ is adjacent to $s$ and not to $s'$. Thus $v \in A_c$ for some $c \in C$. Choose $c' \in C \setminus \{c\}$. Then $G[\{b, s, s', v, c'\}]$ is a fork, a contradiction. This proves (4).

(5) $B = \emptyset$.

Suppose that $B \neq \emptyset$, and so $|B| = 1$ by (3). Let $B = \{b\}$ say. Define $A_b = S$; thus $C \cup \{b\}$ is a clique $D$ say, and every vertex not in $D$ belongs to one of the sets $A_d$ ($d \in D$), and has a unique
neighbour in \( D \); and each set \( A_d \) (\( d \in D \)) is a clique by (4). We claim:

- for all distinct \( d_1, d_2 \in D \), each vertex in \( A_{d_1} \) has at most one neighbour in \( A_{d_2} \); and
- for all distinct \( d_1, d_2, d_3 \in D \), if \( a_i \in A_{d_i} \) for \( i = 1, 2, 3 \), and \( a_1 \) is adjacent to both \( a_2, a_3 \), then \( a_2, a_3 \) are adjacent.

To see the first claim, suppose that \( v \in A_{d_1} \) has two neighbours \( a, a' \in A_{d_2} \). By (4), \( a, a' \) are adjacent. Choose \( d_3 \in D \setminus \{d_1, d_2\} \); then \( G[v, a, a', d_1, d_2, d_3] \) is an antifork, a contradiction. This proves the first claim.

For the second claim, suppose that \( d_1, d_2, d_3 \in D \) are distinct, and \( a_i \in A_{d_i} \) for \( i = 1, 2, 3 \), and \( a_1 \) is adjacent to \( a_2, a_3 \), and \( a_2, a_3 \) are not adjacent. Since \( |C| \geq 3 \) it follows that \( |D| \geq 4 \); choose \( d_4 \in D \setminus \{d_1, d_2, d_3\} \). Then \( G[a_1, a_2, a_3, d_1, d_4] \) is a fork, a contradiction. This proves the second claim.

Let \( H \) be the subgraph of \( G \) obtained by deleting the edges of the cliques \( A_d \cup \{d\} \) (\( d \in D \)). From the two bullet claims above, it follows that each component of \( H \) is a clique. (\( D \) itself is one such component.) Thus we have found two sets of cliques of \( G \); the sets \( A_d \cup \{d\} \) (\( d \in D \)), and the components of \( H \). Each vertex of \( G \) belongs to exactly one clique in the first set, and exactly one in the second; and every edge of \( G \) belongs to one of the cliques in one of the sets. Consequently \( G \) is the line graph of a bipartite graph, a contradiction. This proves that \( B = \emptyset \), and so proves (5).

It follows that \( C \) is a maximal clique of \( G \), and is nondomninating. Moreover, every edge of \( C \) is not contained in any other maximal clique; and every vertex \( c \in C \) is contained in at most two maximal cliques, namely \( C \) and \( A_c \cup \{c\} \) if \( A_c \neq \emptyset \). To complete the proof of the theorem, we only need to show that the cliques \( A_c \cup \{c\} \) are not dominating. Suppose then that \( c \in C \), and \( A_c \cup \{c\} \) is dominating. We have already shown that every clique including a nondominating triangle is itself nondominating; and consequently all triangles included in \( A_c \cup \{c\} \) are dominating. But \( A_c \) is not dominating (because there is a vertex in \( C \) anticomplete to \( A_c \)), so \( |A_c| \leq 2 \).

(6) \( |A_c| = 2 \).

Suppose that \( |A_c| \leq 1 \); and hence \( |A_c| = 1 \), since \( A_c \cup \{c\} \) is dominating and so is not anticomplete to \( S \), since \( S \neq \emptyset \). Let \( A_c = \{a\} \); then for the same reason, \( a \) is complete to \( S \). Also \( A_c \cup \{c\} \) is not anticomplete to any vertex in \( A_{c'} \) for \( c' \in C \setminus \{c\} \); and so \( a \) is complete to \( A_{c'} \) for all \( c' \in C \setminus \{c\} \). Let \( c' \in C \setminus \{c\} \), and let \( a' \in A_{c'} \), and \( s \in S \). Since \( a', s \) are both adjacent to \( a \), and there exists \( c'' \in C \setminus \{c, c'\} \), and \( G[s, a, a', c, c''] \) is not a fork, it follows that \( a', s \) are adjacent; and so \( A_{c'} \) is complete to \( S \) for all \( c' \in C \setminus \{c\} \). If \( c_1, c_2 \in C \setminus \{c\} \) are distinct, and there exist \( a_1 \in A_{c_1} \) and \( a_2 \in A_{c_2} \) nonadjacent, then \( G[a_1, a_2, a, s, c_1] \) is an antifork, a contradiction; so all the sets \( A_{c'} \) (\( c' \in C \)) are complete to each other. Hence they are all homogeneous sets, and so is \( C \); so they all have cardinality at most one. But then \( G \) is the line graph of a bipartite graph, a contradiction. This proves (6).

Let \( A_c = \{p, q\} \) say. For each \( d \in C \setminus \{c\} \), let \( P_d \) be the set of vertices in \( A_d \) adjacent to \( p \), and define \( Q_d \) similarly.

(7) For each \( d \in C \setminus \{c\} \), \( P_d \cup Q_d \) are disjoint subsets of \( A_d \) with union \( A_d \).
Since $A_c \cup \{c\}$ is dominating, it follows that $P_d \cup Q_d = A_d$ for each $d \in C \setminus \{c\}$. If some $v \in A_d$ is adjacent to both $p, q$, choose $c' \in C \setminus \{c, d\}$; then $G[\{v, p, q, c, c'\}]$ is an antifork, a contradiction. This proves (7).

Let $S_p, S_q$ be the sets of vertices in $S$ adjacent to $p, q$ respectively.

(8) $S_p, S_q$ are disjoint subsets of $S$ with union $S$; and $P_d$ is complete to $S_p$, and $Q_d$ is complete to $S_q$; for each $d \in C \setminus \{c\}$.

Certainly $S_p \cup S_q = S$ since $A_c \cup \{c\}$ is dominating. If $v \in S$ is adjacent to both $p, q$, choose $d \in C \setminus \{c\}$, and then $G[\{v, p, q, c, d\}]$ is an antifork, a contradiction. This proves the first assertion.

Let $d \in C \setminus \{c\}$, and suppose $v \in P_d$ and $s \in S_p$ are nonadjacent. Choose $c' \in C \setminus \{c, d\}$, and then $G[\{v, s, p, c, c'\}]$ is a fork, a contradiction. Thus $P_d$ is complete to $S_p$, and similarly $Q_d$ is complete to $S_q$, for each $d \in C \setminus \{c\}$. This proves (8).

Let $P = \{p\} \cup S_p \cup \bigcup_{d \in C \setminus \{c\}} P_d$, and define $Q$ similarly.

(9) $P, Q$ are cliques.

Suppose $P$ is not a clique, say, and choose $p_1, p_2 \in P$, nonadjacent. Thus $p_1, p_2 \neq p$. If $p_1 \in S$, then by (8), $p_2 \in S$, and then $G[\{p_1, p_2, p, c, d\}]$ is a fork where $d \in C \setminus \{c\}$. Thus $p_1, p_2 \notin S$, and so we may assume that $p_i \in P_i$, for $i = 1, 2$, where $c_1, c_2, c \in C$ are distinct. But then $G[\{p_1, p_2, p, q, c_1\}]$ is a fork, a contradiction. This proves (9).

(10) For each $d \in C \setminus \{c\}$, $P_d$ is anticomplete to $S_q$, and $Q_d$ is anticomplete to $S_p$.

Let $v \in P_d$, and suppose $v$ is adjacent to $u \in S_q$. Choose $d' \in C \setminus \{c, d\}$; then $G[\{u, v, p, d, d'\}]$ is a fork, a contradiction. This proves (10).

(11) For each $d \in C \setminus \{c\}$, $P_d$ is anticomplete to $Q \setminus Q_d$.

Suppose that $v \in P_d$ and $u \in Q_{d'}$ are adjacent, where $d' \in C \setminus \{c, d\}$. Choose $s \in S$; then $s$ is adjacent to exactly one of $p, q$. If $s$ is adjacent to $p$, then $s \in P$, and by (9) $s$ is adjacent to $v$; and by (10), $s$ is nonadjacent to $u$. But then $G[\{u, v, s, q, d'\}]$ is a fork, a contradiction. Thus $s$ is adjacent to $q$ and not to $p$, and hence adjacent to $u$ and not to $v$. But then $G[\{u, v, s, p, d\}]$ is a fork, a contradiction. This proves (11).

(12) Every vertex of $S_p$ has at most one neighbour in $S_q$ and vice versa.

Suppose $v \in S_p$ is adjacent to $u, w \in S_q$. Then $G[\{v, u, w, q, c\}]$ is an antifork, a contradiction. This proves (12).

From (11), since $P_d$ is complete to $Q_d$, it follows that $P_d$ is a homogeneous set, and so $|P_d| \leq 1$, and similarly $|Q_d| \leq 1$. Let $H$ be the subgraph obtained from $G$ by deleting the edges of the three cliques $C, P, Q$. By (11), every edge of $H$ either has both ends in $A_d \cup \{d\}$ for some $d \in C$, or has
one end in $S_p$ and the other in $S_q$; and hence by (12), every component of $H$ is a clique, and has at most one vertex in common with $C, P$ or $Q$. Consequently $G$ is the line graph of a bipartite graph, a contradiction. This proves 4.3.

4.4 Let $G$ be an uncluttered graph with no nontrivial homogeneous set, such that $G$ is not the line graph of a bipartite graph. Then either every triangle is dominating, or no clique is dominating.

Proof. Let $A$ be union of the vertex sets of all dominating cliques, and let $B$ be the union of the vertex sets of all nondominating triangles. By 4.3, every vertex of a nondominating triangle only belongs to nondominating cliques, so $A \cap B = \emptyset$. Suppose there is a nondominating triangle $T$, and a dominating clique $C$. Thus $C \subseteq A$, and $T \subseteq B$. Every triangle is either dominating or nondominating, and so is a subset of one of $A, B$. Consequently, every vertex of $T$ has at most one neighbour in $C$; and hence exactly one since $C$ is dominating; and since no vertex in $C$ has more than one neighbour in $T$, it follows that there are three vertices $c_1, c_2, c_3$ such that $t_i$ is adjacent to $c_i$ for $i = 1, 2, 3$, where $T = \{t_1, t_2, t_3\}$. By 4.3, $\{c_1, c_2, c_3\}$ is dominating. Since $T$ is nondominating, there is a vertex $y$ that is anticomplete to $T$; and since $\{c_1, c_2, c_3\}$ is dominating, we may assume that $y$ is adjacent to $c_1$. If $y$ is nonadjacent to $c_2$ then $G[\{y, c_1, t_1, t_3, c_2\}]$ is a fork, a contradiction. But if $y$ is adjacent to $c_2$, then since the triangle $\{y, c_1, c_2\}$ has a vertex in $A$, it follows that $y, c_1, c_2 \in A$, and in particular $\{y, c_1, c_2\}$ is dominating; and this is impossible since $t_3$ has no neighbour in $\{y, c_1, c_2\}$. This proves 4.4.

The claw is the complete bipartite graph $K_{1,3}$, and its centre is its vertex of degree three. Thus if $T$ is a nondominating triangle in $G$, and $v$ has no neighbour in $T$, then $v$ is a claw centre in $\overline{G}$, and vice versa. Let us say an anticlaw is a four-vertex graph whose complement is a claw.

![The claw and the anticlaw.](image)

Figure 3: The claw and the anticlaw.

We need one more lemma.

4.5 Let $G$ be a graph with no claw or anticlaw. Then for one of $G, \overline{G}$, say $H$, either

- each component of $H$ is a path or cycle, and hence $H$ is the line graph of a triangle-free graph, or
- $|V(H)| \leq 9$, and $H$ is a line graph of a bipartite graph.

Proof. The net is the graph on six vertices consisting of three pairwise adjacent vertices $t_1, t_2, t_3$, and three more vertices $s_1, s_2, s_3$, where for $1 \leq i \leq 3$ $s_i$ has degree one and $t_i$ is its unique neighbour. The antinet is the complement graph of the net.
We begin with:

1. If $G$ contains a net as an induced subgraph, then $G$ is a net, and so $G$ is the line graph of a bipartite graph.

Suppose that $s_1, s_2, s_3, t_1, t_2, t_3$ are distinct vertices of $G$ forming a net in the notation above. Let $W = \{s_1, s_2, s_3, t_1, t_2, t_3\}$. If $W = V(G)$ the claim holds, so we assume there exists $v \in V(G) \setminus W$. Since $G[V \setminus \{v, t_1, t_2, t_3\}]$ is not an anticlaw, $v$ is adjacent to one of $t_1, t_2, t_3$, say $t_1$. Since $G[V \setminus \{v, s_1, s_2\}]$ is not a claw, $v$ is nonadjacent to one of $s_1, s_2$, say $s_1$. Since $G[V \setminus \{v, s_1, t_1\}]$ is not an anticlaw, $v$ is nonadjacent to $s_1$. Since $G[V \setminus \{v, t_1, t_3\}]$ is not a claw, $v$ is adjacent to $t_3$. But then $G[V \setminus \{v, t_1, s_2, t_3\}]$ is an anticlaw, a contradiction. This proves (1).

From (1) we may assume $G$ contains no net as an induced subgraph, and by taking complements we may also assume that $G$ contains no antinet.

The bull is the graph with five vertices $t_1, \ldots, t_5$, where $t_1-t_2-t_3-t_4$ is an induced path and $t_5$ is adjacent to $t_2, t_3$ and nonadjacent to $t_1, t_4$.

2. We may assume that $G$ contains a bull as an induced subgraph.

Suppose first that $G$ has no triangle. Then $G$ has maximum degree at most two, since it has no claw or triangle, and the theorem holds. So we may assume that $G$ has a triangle, and (by taking complements) $G$ has a stable set of cardinality three. Choose a triangle $T$, and a set $S$ of three pairwise nonadjacent vertices, with $S \cup T$ minimal. Let $S = \{s_1, s_2, s_3\}$ and $T = \{t_1, t_2, t_3\}$. Certainly $|S \cup T| \leq 6$; suppose that equality holds. If $t_1$ is adjacent to at most one of $s_1, s_2, s_3$, say not to $s_2, s_3$, then $\{t_1, s_2, s_3\}$ is a stable set, contradicting the minimality of $S \cup T$. So each $t_i$ is adjacent to at least two of $s_1, s_2, s_3$. By the same argument in the complement, each $s_j$ is nonadjacent to at least two of $t_1, t_2, t_3$; but this is impossible. Thus $|S \cup T| \leq 5$; and so equality holds, since $|S \cap T| \leq 1$. We may assume that $s_3 = t_3$. Now each of $s_1, s_2$ is adjacent to one of $t_1, t_2$, since it is not an anticlaw centre; but each of $t_1, t_2$ is adjacent to at most one of $s_1, s_2$, since it is not a claw centre. Consequently $G[S \cup T]$ is a bull. This proves (2).
Let $W = \{t_1, \ldots, t_5\}$, and let $G[W]$ be a bull in $G$, with notation as in the definition of a bull. For each vertex $v \in V(G) \setminus W$, let $W(v)$ denote the set of neighbours of $v$ in $W$. Define $A_1, A_2, B_1, B_2$ by:

- $A_1$ is the set of $v \in V(G) \setminus W$ with $W(v) = \{t_1, t_5\}$;
- $A_2$ is the set of $v \in V(G) \setminus W$ with $W(v) = \{t_4, t_5\}$;
- $B_1$ is the set of $v \in V(G) \setminus W$ with $W(v) = \{t_1, t_2, t_4\}$;
- $B_2$ is the set of $v \in V(G) \setminus W$ with $W(v) = \{t_1, t_3, t_4\}$.

(See figure 6.)

![Figure 6: For step (3) of the proof of 4.5.](image)

We claim:

(3) $A_1 \cup A_2 \cup B_1 \cup B_2 = V(G) \setminus W$.

Let $v \in V(G) \setminus W$. Assume first that $t_5 \in W(v)$. Since $G[\{v, t_5, t_1, t_4\}]$ is not a claw, one of $t_1, t_4 \notin W(v)$, say $t_1$. Since $G[\{v, t_5, t_3, t_1\}]$ is not an anticlaw, $t_3 \notin W(v)$. Since $G[\{v, t_1, t_2, t_3\}]$ is not a claw, $t_2 \notin W(v)$. Since $G[W \cup \{v\}]$ is not a net, $t_4 \in W(v)$. But then $v \in A_2$. Now assume that $t_5 \notin W(v)$. We apply the same argument in the complement to deduce that $v \in B_1 \cup B_2$. This proves (3).

(4) Each of $A_1, A_2, B_1, B_2$ has cardinality at most one.

If there exist $u, v \in A_1$, nonadjacent, then $G[\{u, v, t_1, t_2\}]$ is a claw, a contradiction; and if there exist $u, v \in A_1$ adjacent, then $G[\{u, v, t_1, t_4\}]$ is an anticlaw, a contradiction. Similarly $|A_2| \leq 1$; and by taking complements it follows that $|B_1|, |B_2| \leq 1$. This proves (4).

(5) The pairs $(A_1, A_2), (A_1, B_2), (A_2, B_1)$ are complete, and the pairs $(B_1, B_2), (A_1, B_1), (A_2, B_2)$ are anticomplete.

If there exist $u \in A_1$ and $v \in A_2$, nonadjacent, then $G[\{u, v, t_5, t_2\}]$ is a claw; so $A_1$ is complete to $A_2$. By taking complement it follows that $B_1$ is anticomplete to $B_2$. If there exists $u \in A_1$ and $v \in B_1$, adjacent, then $G[\{u, v, t_1, t_3\}]$ is an anticlaw, a contradiction; so $A_1$ is anticomplete to $B_1$, and
and from the symmetry $A_2$ is anticomplete to $B_2$. By taking complements it follows that $A_1$ is complete to $B_2$, and $A_2$ is complete to $B_1$. This proves (5).

From (4) and (5) it follows that $|V(G)| \leq 9$, and $G$ is an induced subgraph of the line graph of $K_{3,3}$. This proves 4.5.

We use 4.5 to prove the next result:

**4.6** Let $G$ be an unchatttered graph with no nontrivial homogeneous set. Then one of $G,\overline{G}$ is the line graph of a triangle-free graph.

**Proof.** Suppose that $G$ contains a claw and an anticlaw. Then since there is a nondominating triangle, it follows from 4.4 that no clique is dominating. If $a \in V(G)$ is a claw centre, then it is in at least three maximal cliques, and so by 4.3 it is in no triangle (since all triangles are nondominating), and hence the set of neighbours of $a$ is stable. By taking complements, if $b \in V(G)$ is an anticlaw centre, then the set of vertices nonadjacent to $b$ is a clique.

(1) *There do not exist a claw centre $a$ and an anticlaw centre $b$ with $a \neq b$.***

Suppose there exist such $a, b$. By taking complements if necessary, we may assume that $a, b$ are nonadjacent. Since $b$ is an anticlaw centre, the set of vertices of $G$ nonadjacent to $b$ is a clique $C$ say. Now $|C| \geq 3$ since $b$ is an anticlaw centre; and $a \in C$. Thus $a$ belongs to a triangle, and so belongs to at most two maximal cliques by 4.3, contradicting that $a$ is a claw centre. This proves (1).

(2) *There do not exist both a claw and an anticlaw in $G$.***

Suppose there is both a claw and an anticlaw. Then by (1), there is a vertex $c$ that is the unique claw centre and the unique anticlaw centre. Let $A$ be its set of neighbours and let $B = V(G) \setminus (A \cup \{c\})$. Since $a$ is not in a triangle, $A$ is stable, and similarly $B$ is a clique. Since $a$ is a claw centre, $|A| \geq 3$, and similarly $|B| \geq 3$. Let $b \in B$. If $b$ has no neighbour in $A$, then the stable set $\{b, c\}$ is a dominating clique of $\overline{G}$, a contradiction to 4.4 applied to $\overline{G}$. If $b$ has at least two nonneighbours in $A$, say $a_1, a_2$, let $b$ be adjacent to $a_3 \in A$ and then $G[\{b, a_1, a_2, a_3, c\}]$ is a fork, a contradiction. But $b$ has at most two neighbours in $A$ since $b$ is not a claw centre; so $|A| = 3$ and $b$ has exactly two neighbours in $A$. By the same argument applied in the complement, $|B| = 3$ and every vertex in $A$ has exactly two nonneighbours in $B$, which is impossible by counting edges between $A$ and $B$. This proves (2).

From (2), and taking complements if necessary, we may assume there is no claw in $G$. By 4.5 we may assume that there is an anticlaw in $G$, that is, there is a nondominating triangle. Consequently every clique is nondominating, by 4.4. By 4.3 every vertex that is in a triangle is in at most two maximal cliques. But if a vertex $v$ is not in any triangle, then since there is no claw it follows that $v$ has degree at most two, and so $v$ is in at most two maximal cliques. This proves that every vertex is in at most two maximal cliques. Let $C_1, \ldots, C_t$ be the maximal cliques of $G$, and make a graph $H$ with vertex set $\{1, \ldots, t\}$, where distinct $i, j$ are adjacent if $C_i \cap C_j \neq \emptyset$. If $|C_i \cap C_j| \geq 2$ for some distinct $i, j$, then $C_i \cap C_j$ is a nontrivial homogeneous set of $G$, which is impossible; so $G$ is the line graph of $H$. It remains to show that $H$ is triangle-free. Suppose not; then we may assume that $C_1, C_2, C_3$ pairwise intersect. Choose $v_1 \in C_2 \cap C_3$, and define $v_2, v_3$ similarly. Then $v_1 \notin C_1$, since
every vertex belongs to at most two of $C_1, C_2, C_3$, and similarly $v_2 \notin C_2$ and $v_3 \notin C_3$. But $v_1, v_2$ are adjacent, because they both belong to $C_3$, and similarly $v_1, v_2, v_3$ are pairwise adjacent; and hence there is a maximal clique containing all three of $v_1, v_2, v_3$. It is different from $C_1, C_2, C_3$, and so $v_1$ belongs to three different maximal cliques, a contradiction. This proves that $H$ is triangle-free, and so proves 4.6.

By combining 4.6 and 3.1, we deduce our main result, which we restate:

4.7 Let $G$ be an uncluttered graph. Then either

- one of $G, \overline{G}$ is disconnected; or
- one of $G, \overline{G}$ has adjacent simplicial twins; or
- one of $G, \overline{G}$ is candled; or
- one of $G, \overline{G}$ is the line graph of a triangle-free graph.

Proof. If $G$ is a nontrivial homogeneous set, then by 3.1 either one of $G, \overline{G}$ is disconnected, or one of $G, \overline{G}$ has adjacent simplicial twins, or one of $G, \overline{G}$ is candled, and in each case the theorem holds. If $G$ has no nontrivial homogeneous set, then by 4.6, one of $G, \overline{G}$ is the line graph of a triangle-free graph, and again the theorem holds. This proves 4.7.

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References

